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April 1<sup>st</sup>, 2010  
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# Application Note

## 78K0R/Kx3-L, 78K0R/Ix3, 78K0R/Kx3-C

### 16-bit Single-Chip Microcontrollers

### Flash Memory Programming (Programmer)

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78K0R/KC3-L:  $\mu$ PD78F1000, 78F1001, 78F1002, 78F1003

78K0R/KD3-L:  $\mu$ PD78F1004, 78F1005, 78F1006

78K0R/KE3-L:  $\mu$ PD78F1007, 78F1008, 78F1009

78K0R/KF3-L:  $\mu$ PD78F1010, 78F1011, 78F1012

78K0R/KG3-L:  $\mu$ PD78F1013, 78F1014

78K0R/IB3:  $\mu$ PD78F1201, 78F1203

78K0R/IC3:  $\mu$ PD78F1211, 78F1213, 78F1214, 78F1215

78K0R/ID3:  $\mu$ PD78F1223, 78F1224, 78F1225

78K0R/IE3:  $\mu$ PD78F1233, 78F1234, 78F1235

78K0R/KF3-C:  $\mu$ PD78F1846, 78F1847

78K0R/KG3-C:  $\mu$ PD78F1848, 78F1849

[MEMO]

## NOTES FOR CMOS DEVICES

### ① VOLTAGE APPLICATION WAVEFORM AT INPUT PIN

Waveform distortion due to input noise or a reflected wave may cause malfunction. If the input of the CMOS device stays in the area between  $V_{IL}$  (MAX) and  $V_{IH}$  (MIN) due to noise, etc., the device may malfunction. Take care to prevent chattering noise from entering the device when the input level is fixed, and also in the transition period when the input level passes through the area between  $V_{IL}$  (MAX) and  $V_{IH}$  (MIN).

### ② HANDLING OF UNUSED INPUT PINS

Unconnected CMOS device inputs can be cause of malfunction. If an input pin is unconnected, it is possible that an internal input level may be generated due to noise, etc., causing malfunction. CMOS devices behave differently than Bipolar or NMOS devices. Input levels of CMOS devices must be fixed high or low by using pull-up or pull-down circuitry. Each unused pin should be connected to  $V_{DD}$  or GND via a resistor if there is a possibility that it will be an output pin. All handling related to unused pins must be judged separately for each device and according to related specifications governing the device.

### ③ PRECAUTION AGAINST ESD

A strong electric field, when exposed to a MOS device, can cause destruction of the gate oxide and ultimately degrade the device operation. Steps must be taken to stop generation of static electricity as much as possible, and quickly dissipate it when it has occurred. Environmental control must be adequate. When it is dry, a humidifier should be used. It is recommended to avoid using insulators that easily build up static electricity. Semiconductor devices must be stored and transported in an anti-static container, static shielding bag or conductive material. All test and measurement tools including work benches and floors should be grounded. The operator should be grounded using a wrist strap. Semiconductor devices must not be touched with bare hands. Similar precautions need to be taken for PW boards with mounted semiconductor devices.

### ④ STATUS BEFORE INITIALIZATION

Power-on does not necessarily define the initial status of a MOS device. Immediately after the power source is turned ON, devices with reset functions have not yet been initialized. Hence, power-on does not guarantee output pin levels, I/O settings or contents of registers. A device is not initialized until the reset signal is received. A reset operation must be executed immediately after power-on for devices with reset functions.

### ⑤ POWER ON/OFF SEQUENCE

In the case of a device that uses different power supplies for the internal operation and external interface, as a rule, switch on the external power supply after switching on the internal power supply. When switching the power supply off, as a rule, switch off the external power supply and then the internal power supply. Use of the reverse power on/off sequences may result in the application of an overvoltage to the internal elements of the device, causing malfunction and degradation of internal elements due to the passage of an abnormal current.

The correct power on/off sequence must be judged separately for each device and according to related specifications governing the device.

### ⑥ INPUT OF SIGNAL DURING POWER OFF STATE

Do not input signals or an I/O pull-up power supply while the device is not powered. The current injection that results from input of such a signal or I/O pull-up power supply may cause malfunction and the abnormal current that passes in the device at this time may cause degradation of internal elements. Input of signals during the power off state must be judged separately for each device and according to related specifications governing the device.

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## INTRODUCTION

<b>Caution</b>	Explanations in this application note assume use of the 78K0R/Kx3-L as the representative microcontroller. Users of a product other than the 78K0R/Kx3-L should read 78K0R/Kx3-L as referring to that product.												
<b>Target Readers</b>	This application note is intended for users who understand the functions of the 78K0R/Kx3-L, 78K0R/lx3, and 78K0R/Kx3-C and who will use these products to design application systems.												
<b>Purpose</b>	<p>The purpose of this application note is to help users understand how to develop dedicated flash memory programmers for rewriting the internal flash memory of the 78K0R/Kx3-L, 78K0R/lx3, and 78K0R/Kx3-C.</p> <p>The sample programs and circuit diagrams shown in this document are for reference only and are not intended for use in actual design-ins.</p> <p>Therefore, these sample programs must be used at the user's own risk. Correct operation is not guaranteed if these sample programs are used.</p>												
<b>Organization</b>	<p>This manual consists of the following main sections.</p> <ul style="list-style-type: none"><li>• Flash memory programming</li><li>• Command/data frame format</li><li>• Description of command processing</li><li>• UART communication mode</li><li>• Flash memory programming parameter characteristics</li></ul>												
<b>How to Read This Manual</b>	<p>It is assumed that the reader of this manual has general knowledge in the fields of electrical engineering, logic circuits, and microcontrollers.</p> <ul style="list-style-type: none"><li>• To gain a general understanding of functions: → Read this manual in the order of the <b>CONTENTS</b>.</li><li>• To learn more about hardware functions of the 78K0R/Kx3-L, 78K0R/lx3, and 78K0R/Kx3-C: → See the user's manual of the 78K0R/Kx3-L, 78K0R/lx3, or 78K0R/Kx3-C.</li></ul>												
<b>Conventions</b>	<table><tr><td>Data significance:</td><td>Higher digits on the left and lower digits on the right</td></tr><tr><td>Active low representation:</td><td><math>\overline{xxx}</math> (overscore over pin or signal name)</td></tr><tr><td><b>Note:</b></td><td>Footnote for item marked with <b>Note</b> in the text</td></tr><tr><td><b>Caution:</b></td><td>Information requiring particular attention</td></tr><tr><td><b>Remark:</b></td><td>Supplementary information</td></tr><tr><td>Numeral representation:</td><td>Binary.....xxxx or xxxxB Decimal .....xxxx Hexadecimal .....xxxxH</td></tr></table>	Data significance:	Higher digits on the left and lower digits on the right	Active low representation:	$\overline{xxx}$ (overscore over pin or signal name)	<b>Note:</b>	Footnote for item marked with <b>Note</b> in the text	<b>Caution:</b>	Information requiring particular attention	<b>Remark:</b>	Supplementary information	Numeral representation:	Binary.....xxxx or xxxxB Decimal .....xxxx Hexadecimal .....xxxxH
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<b>Caution:</b>	Information requiring particular attention												
<b>Remark:</b>	Supplementary information												
Numeral representation:	Binary.....xxxx or xxxxB Decimal .....xxxx Hexadecimal .....xxxxH												

**Related Documents**

The related documents indicated in this publication may include preliminary versions. However, preliminary versions are not marked as such.

**Documents related to the 78K0R/Kx3-L, 78K0R/lx3, and 78K0R/Kx3-C**

Document Name	Document No.
78K0R/Kx3-L User's Manual	U19291E
78K0R/IB3 User's Manual	U19018E
78K0R/IC3 User's Manual	U19120E
78K0R/ID3 User's Manual	U19185E
78K0R/IE3 User's Manual	U19163E
78K0R/KG3-C User's Manual	U19274E
78K0R Microcontrollers Instructions User's Manual	U17792E

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## CHAPTER 1 FLASH MEMORY PROGRAMMING

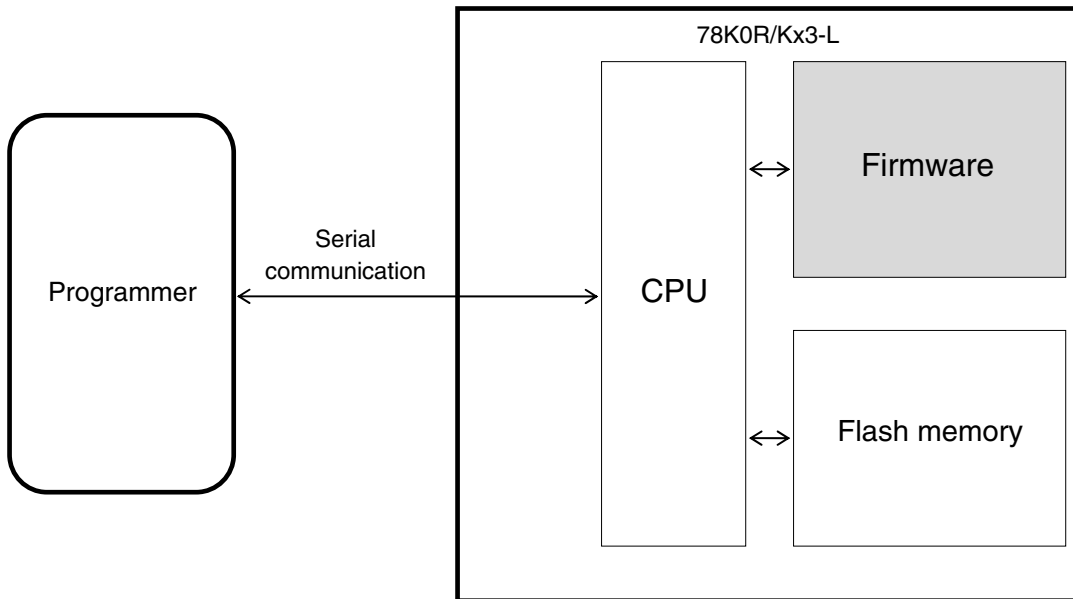
To rewrite the contents of the internal flash memory of the 78K0R/Kx3-L, a dedicated flash memory programmer (hereafter referred to as the “programmer”) is usually used.

This Application Note explains how to develop a dedicated programmer.

### 1.1 Overview

The 78K0R/Kx3-L incorporates firmware that controls flash memory programming. The programming to the internal flash memory is performed by transmitting/receiving commands between the programmer and the 78K0R/Kx3-L via serial communication.

Figure 1-1. System Outline of Flash Memory Programming in 78K0R/Kx3-L



## 1.2 System Configuration

Examples of the system configuration for programming the flash memory are illustrated in Figure 1-2.

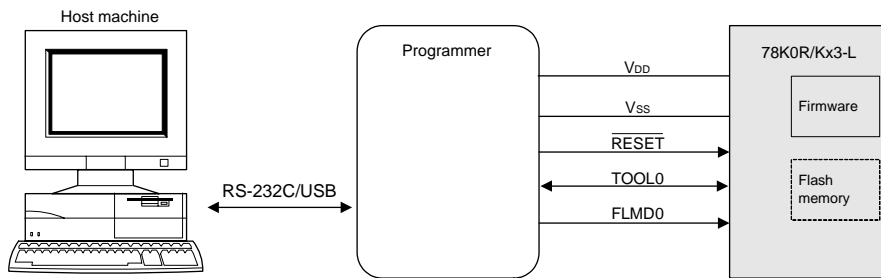
These figures illustrate how to program the flash memory with the programmer, under control of a host machine.

Depending on how the programmer is connected, the programmer can be used in a standalone mode without using the host machine, if a user program has been downloaded to the programmer in advance.

For example, NEC Electronics' flash memory programmer PG-FP5 can execute programming either by using the GUI software with a host machine connected or by itself (standalone).

**Figure 1-2. System Configuration Example**

Single-wire UART communication mode (LSB-first transfer)



- Remarks 1.** The 78K0R/Kx3-L can only communicate via the single-wire UART communication mode.
- 2.** For the pins used by flash memory programming and the recommended connections of unused pins, see the user's manual of each product.

### 1.3 Flash Memory Configuration

The 78K0R/Kx3-L must manage product-specific information (such as device name and memory information) via the programmer.

Table 1-1 shows the flash memory size of the 78K0R/Kx3-L and Figure 1-3 shows the configuration of the flash memory.

**Table 1-1. Size of Flash Memory for Each Product**

**(a) Size of Flash Memory for 78K0R/Kx3-L**

Device Name	Flash Memory Size
$\mu$ PD78F1000	16 KB
$\mu$ PD78F1001, 78F1004, 78F1007	32 KB
$\mu$ PD78F1002, 78F1005, 78F1008	48 KB
$\mu$ PD78F1003, 78F1006, 78F1009, 78F1010	64 KB
$\mu$ PD78F1011, 78F1013	96 KB
$\mu$ PD78F1012, 78F1014	128 KB

**(b) Size of Flash Memory for 78K0R/Ix3**

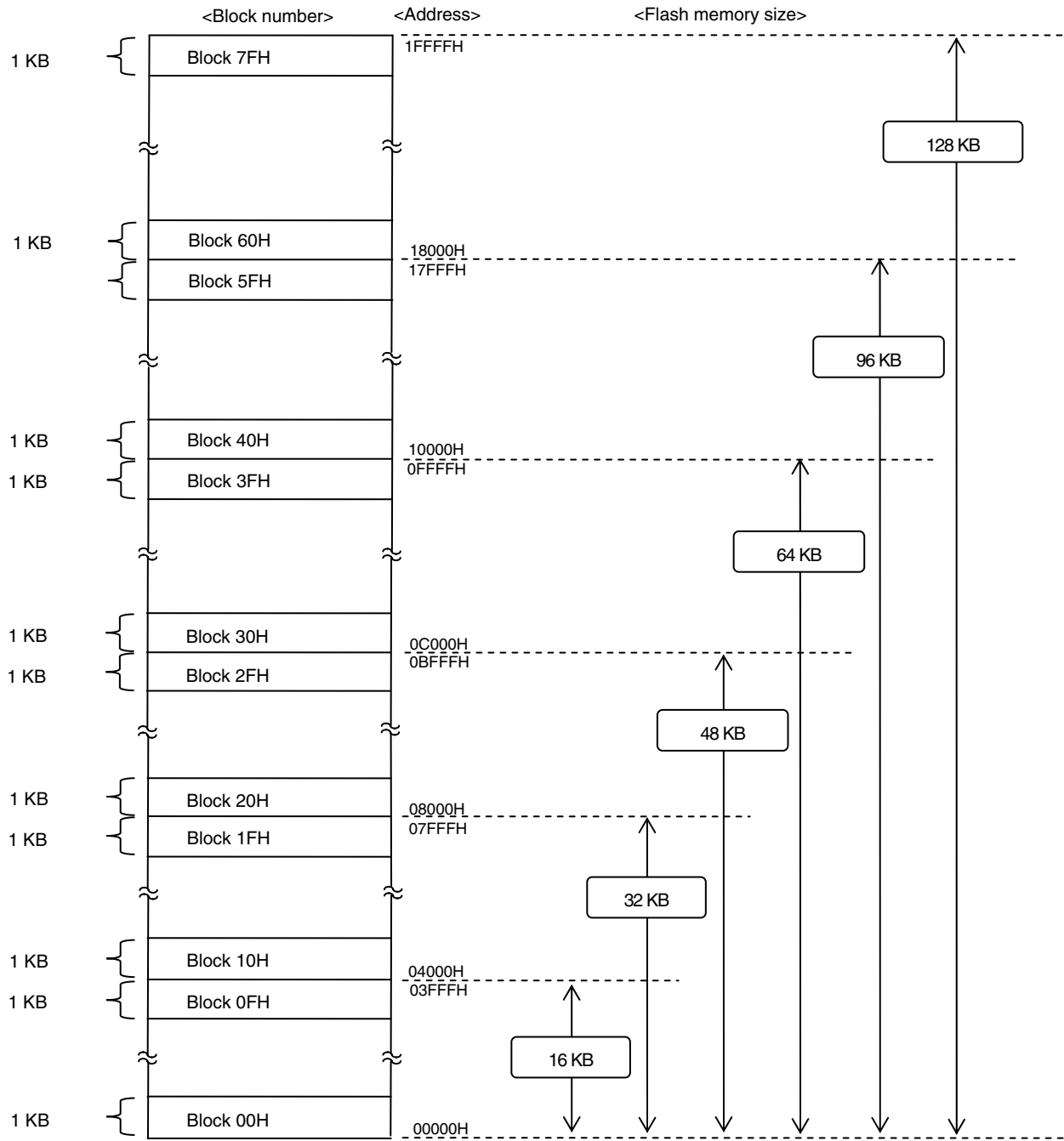
Device Name	Flash Memory Size
$\mu$ PD78F1211	16 KB
$\mu$ PD78F1213, 78F1223, 78F1233	32 KB
$\mu$ PD78F1214, 78F1224, 78F1234	48 KB
$\mu$ PD78F1215, 78F1225, 78F1235	64 KB

**(c) Size of Flash Memory for 78K0R/Kx3-C**

Device Name	Flash Memory Size
$\mu$ PD78F1846, 78F1848	96 KB
$\mu$ PD78F1847, 78F1849	128 KB

**Remark** Products under development are included in the above tables..

Figure 1-3. Flash Memory Configuration



**Remark** Each block consists of 1 KB (this figure only illustrates some blocks in the flash memory).



## 1.4 Command List and Status List

The flash memory incorporated in the 78K0R/Kx3-L can be rewritten by using the commands listed in Table 1-2. The programmer transmits commands to control these functions to the 78K0R/Kx3-L, and checks the response status sent from the 78K0R/Kx3-L, to manipulate the flash memory.

### 1.4.1 Command list

The commands used by the programmer and their functions are listed below.

**Table 1-2. List of Commands Transmitted from Programmer to 78K0R/Kx3-L**

Command Number	Command Name	Function Name	Function
20H	Chip Erase	Erase	Erases the entire flash memory area.
22H	Block Erase		Erases a specified area in the flash memory.
40H	Programming	Write	Writes data to a specified area in the flash memory.
13H	Verify	Verify	Compares the contents in a specified area in the flash memory with the data transmitted from the programmer.
32H	Block Blank Check	Blank check	Checks the erase status of a specified block in the flash memory.
C0H	Silicon Signature	Information acquisition	Acquires 78K0R/Kx3-L information (product name, flash memory configuration, etc.).
C5H	Version Get		Acquires version of the 78K0R/Kx3-L and firmware.
B0H	Checksum		Acquires checksum data of a specified area.
A0H	Security Set	Security	Sets security information.
00H	Reset	Others	Detects synchronization in communication.
9AH	Baud Rate Set		Sets the baud rate when UART communication mode is selected.

### 1.4.2 Status list

The following table lists the status codes the programmer receives from the 78K0R/Kx3-L.

**Table 1-3. Status Code List**

Status Code	Status	Description
04H	Command number error	Error returned if a command not supported is received
05H	Parameter error	Error returned if command information (parameter) is invalid
06H	Normal acknowledgment (ACK)	Normal acknowledgment
07H	Checksum error	Error returned if data in a frame transmitted from the programmer is abnormal
0FH	Verify error	Error returned if a verify error has occurred upon verifying data transmitted from the programmer
10H	Protect error	Error returned if an attempt is made to execute processing that is prohibited by the Security Set command
15H	Negative acknowledgment (NACK)	Negative acknowledgment
1AH	MRG10 error	Erase verify error
1BH	MRG11 error	Internal verify error or blank check error during data write
1CH	Write error	Write error

Reception of a checksum error or NACK is treated as an immediate abnormal end in this manual. When a dedicated programmer is developed, however, the processing may be retried without problem from the wait immediately before transmission of the command that results a checksum error or NACK. In this event, limiting the retry count is recommended for preventing infinite repetition of the retry operation.

Although not listed in the above table, if a time-out error (BUSY time-out or time-out in data frame reception during UART communication) occurs, it is recommended to shutdown the power supply to the 78K0R/Kx3-L (refer to **1.6 Shutting Down Target Power Supply**) and then connect the power supply again.

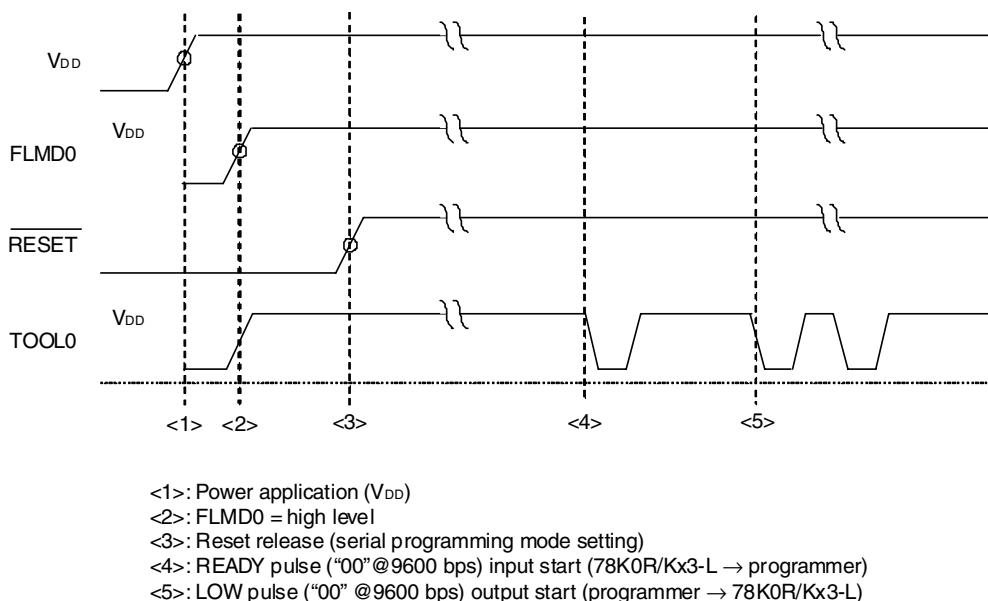
### 1.5 Power Application and Setting Flash Memory Programming Mode

To rewrite the contents of the flash memory with the programmer, the 78K0R/Kx3-L must first be set to the flash memory programming mode by supplying a specific voltage to the flash memory programming mode setting pin (FLMD0) in the 78K0R/Kx3-L, then releasing a reset.

To select a communication mode for flash memory rewriting, input pulses to the FLMD0 pin in the flash memory programming mode.

The following figure illustrates a timing chart for setting the flash memory programming mode and selecting the communication mode.

**Figure 1-4. Setting Flash Memory Programming Mode and Selecting Communication Mode**



The relationship between the setting of the FLMD0 pin after reset release and the operating mode is shown below.

**Table 1-4. Relationship Between FLMD0 Pin Setting After Reset Release and Operating Mode**

FLMD0	Operating Mode
Low (GND)	Normal operating mode
High ( $V_{DD}$ )	Flash memory programming mode

The following table shows the communication mode that can be selected with the 78K0R/Kx3-L, the number of pulses input to FLMD0 (pulse count), and the port to be used.

**Table 1-5. Relationship Between FLMD0 Pulse Count in 78K0R/Kx3-L and Communication Mode**

Communication Mode	FLMD0 Pulse Count	Port Used for Communication
Single-wire UART	0	TOOL0 (P40)

### 1.5.1 UART communication mode

The TOOL0 pin is used for UART communication. The communication conditions are as shown below.

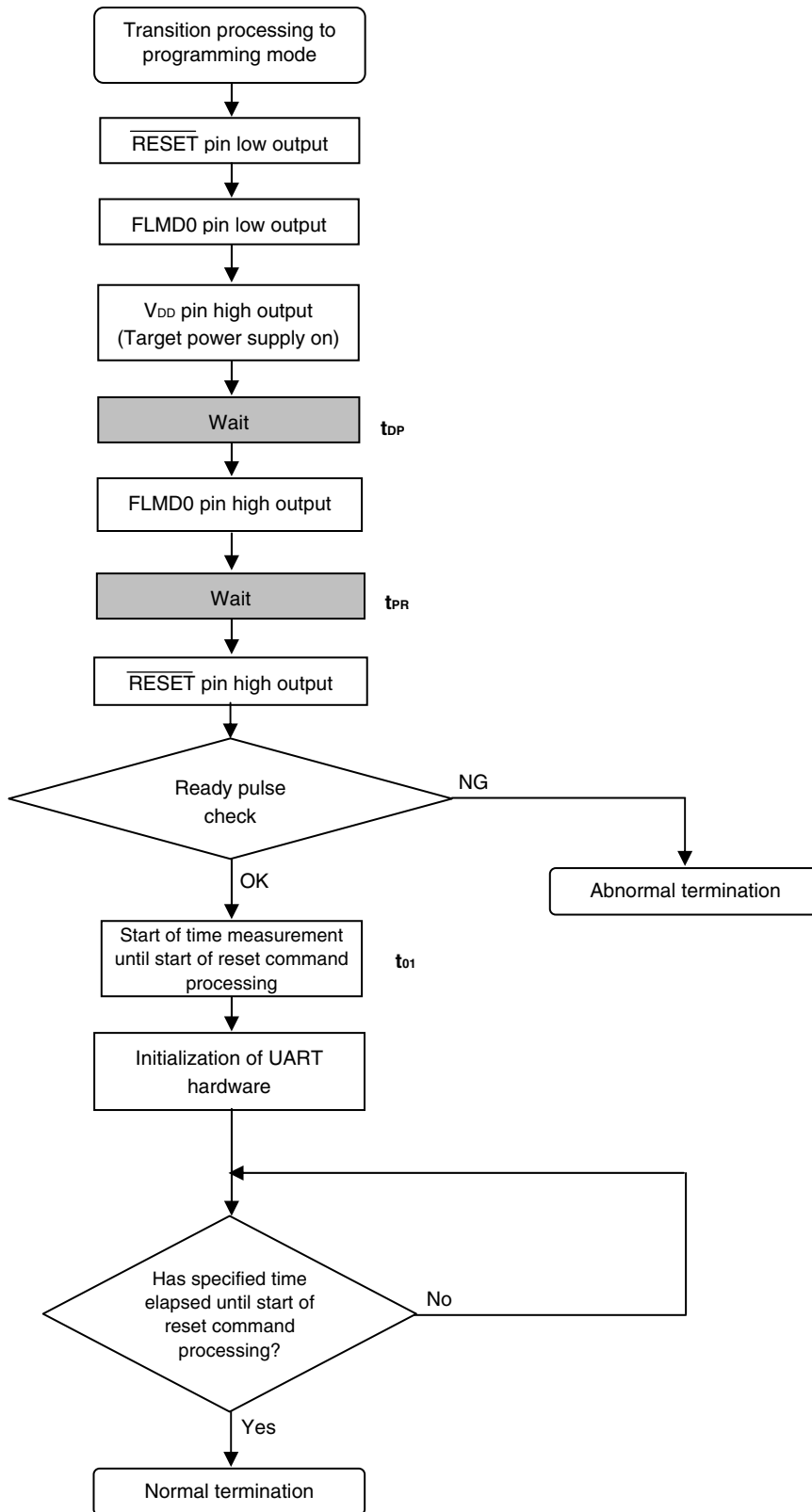
**Table 1-6. UART Communication Conditions**

Item	Description
Baud rate	Communication is performed at 9,600 bps until the Baud Rate Set command for baud rate setting command processing is transmitted. The transmission rate is changed to the baud rate set by the Baud Rate Set command from the transmission of the Reset command for baud rate command processing. For details of the settable baud rate, refer to <b>3.2 Baud Rate Set Command</b> .
Parity bit	None
Data length	8 bits (LSB first)
Stop bit	2 bits (programmer → 78K0R/Kx3-L)/1 bit (78K0R/Kx3-L → programmer)

The role of the master and slave is occasionally exchanged during UART communication, so communication at the optimum timing is possible.

**Caution** Set the same baud rate to the master and slave devices when performing UART communication.

1.5.2 Mode setting flowchart



### 1.5.3 Sample program

The following shows a sample program for mode setting processing.

```

/*****
/*
/* connect to Flash device
/*
/*
/*****
u16    fl_con_dev(void)
{
extern void  init_fl_uart(void);
extern void  init_fl_csi(void);
extern void  stop_UART0(void);

    u16    rc = NO_ERROR;

    SRMK0 = true;           // disable UART Rx INT.
    UARTE0 = false;        // disable UART H.W.
    stop_UART0();         // TxD/RxD = Hi-Z

    pFL_RES      = low;      // RESET = low
    pmFL_FLMD0   = PM_OUT;   // FLMD0 = Low output
    pFL_FLMD0    = low;
    FL_VDD_HI();           // VDD = high

    fl_wait(tDP);         // wait

    pFL_FLMD0    = hi;      // FLMD0 = high
    fl_wait(tPR);         // wait

    pFL_RES      = hi;      // RESET = high

    rc = check_ready_pulse(); // check "READY PULSE" from target device
    if (rc){
        return rc;         // pulse width/timing error
    }
    start_flto(t01);       // start "t01" wait timer

    init_fl_uart();        // Initialize UART h.w.(for Flash device control)
    UARTE0 = true;         // enable UART h.w.
    SRIF0 = false;        // clear UART Rx IRQ flag
    SRMK0 = false;        // enable UART Rx INT.

    while(!check_flto())  // timeout "t01" ?
        ;                 // no

    return rc;
    // start RESET command proc.
}

```

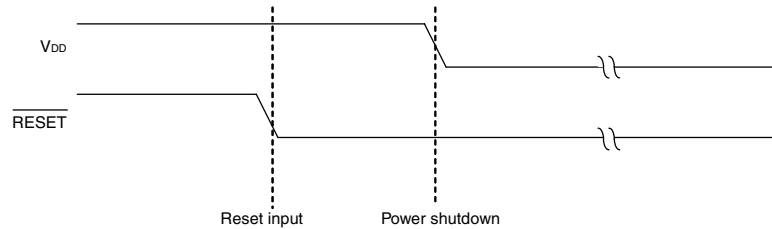
## 1.6 Shutting Down Target Power Supply

After each command execution is completed, shut down the power supply to the target after setting the  $\overline{\text{RESET}}$  pin to low level, as shown below.

Set other pins to Hi-Z when shutting down the power supply to the target.

**Caution** Shutting down the power supply and inputting a reset during command processing are prohibited.

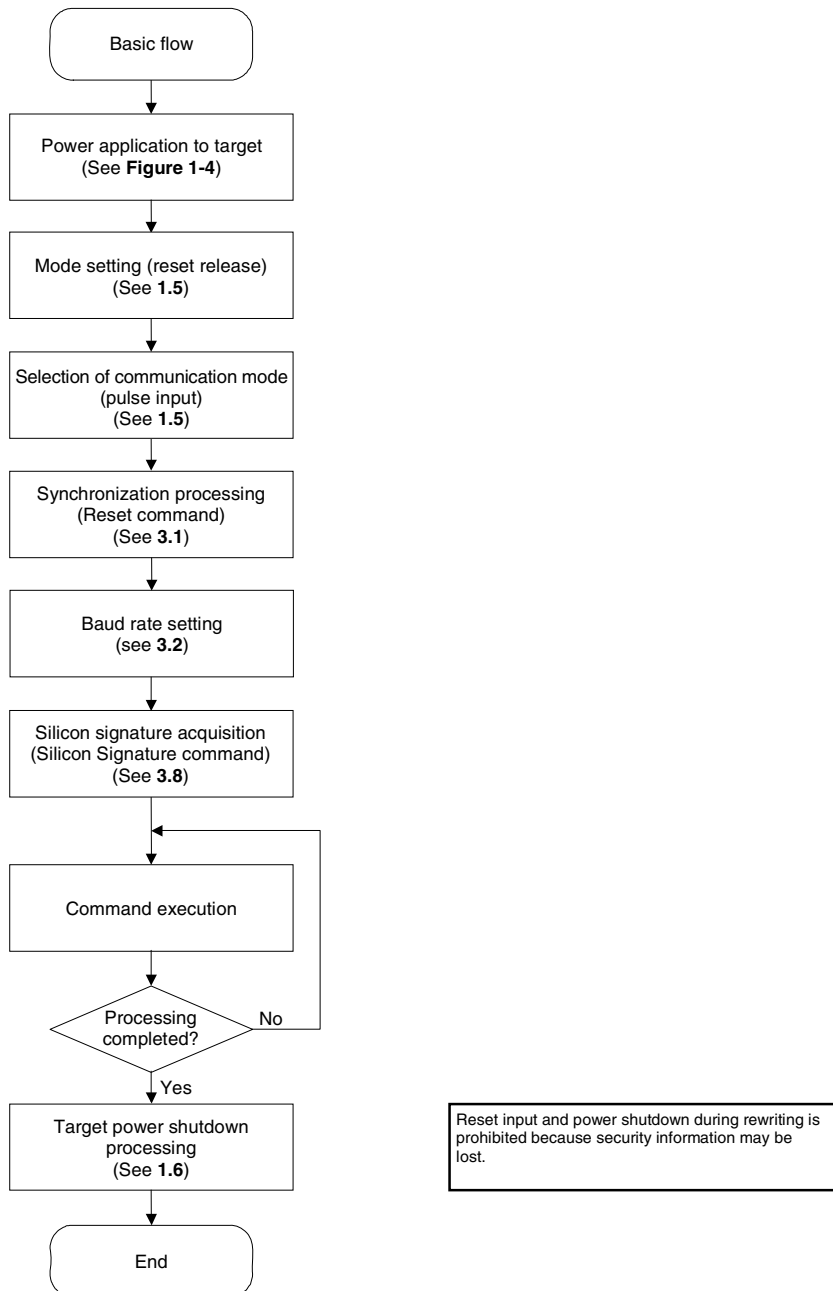
**Figure 1-5. Timing for Terminating Flash Memory Programming Mode**



## 1.7 Command Execution Flow at Flash Memory Rewriting

Figure 1-6 illustrates the basic flowchart when flash memory rewriting is performed with the programmer. Other than commands shown in Figure 1-6, the Verify command and Checksum command are also supported.

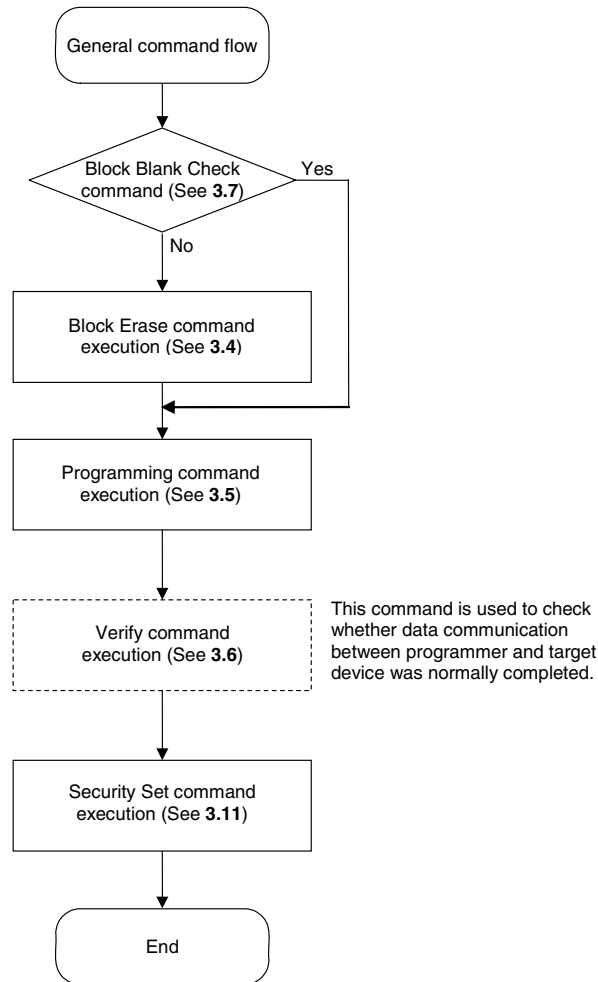
Figure 1-6. Basic Flowchart for Flash Memory Rewrite Processing



**Remark** The example of each command execution is shown in Figure 1-7.



Figure 1-7. General Command Execution Flow at Flash Memory Rewriting



## CHAPTER 2 COMMAND/DATA FRAME FORMAT

The programmer uses the command frame to transmit commands to the 78K0R/Kx3-L. The 78K0R/Kx3-L uses the data frame to transmit write data or verify data to the programmer. A header, footer, data length information, and checksum are appended to each frame to enhance the reliability of the transferred data.

The following shows the format of a command frame and data frame.

**Figure 2-1. Command Frame Format**

SOH (1 byte)	LEN (1 byte)	COM (1 byte)	Command information (variable length) (Max. 255 bytes)	SUM (1 byte)	ETX (1 byte)
-----------------	-----------------	-----------------	---	-----------------	-----------------

**Figure 2-2. Data Frame Format**

STX (1 byte)	LEN (1 byte)	Data (variable length) (Max. 256 bytes)	SUM (1 byte)	ETX or ETB (1 byte)
-----------------	-----------------	--	-----------------	------------------------

**Table 2-1. Description of Symbols in Each Frame**

Symbol	Value	Description
SOH	01H	Command frame header
STX	02H	Data frame header
LEN	–	Data length information (00H indicates 256). Command frame: COM + command information length Data frame: Data field length
COM	–	Command number
SUM	–	Checksum data for a frame Obtained by sequentially subtracting all of calculation target data from the initial value (00H) in 1-byte units (borrow is ignored). The calculation targets are as follows. Command frame: LEN + COM + all of command information Data frame: LEN + all of data
ETB	17H	Footer of data frame other than the last frame
ETX	03H	Command frame footer, or footer of last data frame

The following shows examples of calculating the checksum (SUM) for a frame.

**[Command frame]**

No command information is included in the following example of a Chip Erase command frame, so LEN and COM are targets of checksum calculation.

SOH	LEN	COM	SUM	ETX
01H	01H	20H	Checksum	03H
Checksum calculation targets				

For this command frame, checksum data is obtained as follows.

$$00\text{H (initial value)} - 01\text{H (LEN)} - 20\text{H (COM)} = \text{DFH (Borrow ignored. Lower 8 bits only.)}$$

The command frame finally transmitted is as follows.

SOH	LEN	COM	SUM	ETX
01H	01H	20H	DFH	03H

**[Data frame]**

To transmit a data frame as shown below, LEN and D1 to D4 are targets of checksum calculation.

STX	LEN	D1	D2	D3	D4	SUM	ETX
02H	04H	FFH	80H	40H	22H	Checksum	03H
Checksum calculation targets							

For this data frame, checksum data is obtained as follows.

$$00\text{H (initial value)} - 04\text{H (LEN)} - \text{FFH (D1)} - 80\text{H (D2)} - 40\text{H (D3)} - 22\text{H (D4)} \\ = 1\text{BH (Borrow ignored. Lower 8 bits only.)}$$

The data frame finally transmitted is as follows.

STX	LEN	D1	D2	D3	D4	SUM	ETX
02H	04H	FFH	80H	40H	22H	1BH	03H

When a data frame is received, the checksum data is calculated in the same manner, and the obtained value is used to detect a checksum error by judging whether the value is the same as that stored in the SUM field of the receive data. When a data frame as shown below is received, for example, a checksum error is detected.

STX	LEN	D1	D2	D3	D4	SUM	ETX
02H	04H	FFH	80H	40H	22H	1AH	03H

↑ Normally 1BH

## 2.1 Command Frame Transmission Processing

For details of the flowchart of processing to transmit command frames, read **4.1 Command Frame Transmission Processing Flowchart**.

## 2.2 Data Frame Transmission Processing

The write data frame (user program), verify data frame (user program), and security data frame (security flag) are transmitted as a data frame.

For details of the flowchart of processing to transmit data frames, read **4.2 Data Frame Transmission Processing Flowchart**.

## 2.3 Data Frame Reception Processing

The status frame, silicon signature data frame, version data frame, and checksum data frame are received as a data frame.

For details of the flowchart of processing to receive data frames, read **4.3 Data Frame Reception Processing Flowchart**.

## CHAPTER 3 DESCRIPTION OF COMMAND PROCESSING

### 3.1 Reset Command

#### 3.1.1 Description

This command is used to check the establishment of communication between the programmer and the 78K0R/Kx3-L after the communication mode is set.

The same baud rate must be set for the programmer and 78K0R/Kx3-L, however, the 78K0R/Kx3-L cannot detect its own baud rate generation clock frequency so the baud rate cannot be set. The 78K0R/Kx3-L is enabled to detect the baud rate generation clock frequency by itself, when "00H" is transmitted twice at 9,600 bps from the programmer, and the 78K0R/Kx3-L measures the low-level width of "00H" and calculates the average of the two sent signals. The baud rate can consequently be set, which enables synchronous detection in communication.

#### 3.1.2 Command frame and status frame

Figure 3-1 shows the format of a command frame for the Reset command, and Figure 3-2 shows the status frame for the command.

**Figure 3-1. Reset Command Frame (from Programmer to 78K0R/Kx3-L)**

SOH	LEN	COM	SUM	ETX
01H	01H	00H (Reset)	Checksum	03H

**Figure 3-2. Status Frame for Reset Command (from 78K0R/Kx3-L to Programmer)**

STX	LEN	Data	SUM	ETX
02H	1	ST1	Checksum	03H

**Remark** ST1: Synchronization detection result

See 4.4 **Reset Command** for details on the flowchart of the processing sequence between the programmer and the 78K0R/Kx3-L, the flowchart of command processing, and the sample program.

## 3.2 Baud Rate Set Command

### 3.2.1 Description

This command is used to change the baud rate for UART communication (9,600 bps by default).

After the Baud Rate Set command has been executed, the Reset command must be executed to check synchronization at the changed baud rate.

### 3.2.2 Command frame and status frame

Figure 3-3 shows the format of a command frame for the Baud Rate Set command, and Figure 3-4 shows the status frame for the command.

**Figure 3-3. Baud Rate Set Command Frame (from Programmer to 78K0R/Kx3-L)**

SOH	LEN	COM	Command Information <sup>Note</sup>					SUM	ETX
01H	06H	9AH	D01	D02H	D02L	D03	D04	Checksum	03H

**Note** For details of the command information setting, refer to **Table 3-1**. If data other than in Table 3-1 is set, a time-out error will occur.

If a time-out error has occurred, execute a hardware reset and re-set the flash memory programming mode.

**Remark**

D01:	Synchronization correction mode
D02H, D02L:	Baud rate setting
D03:	Noise filter setting
D04:	Flash rewriting voltage mode setting

**Table 3-1. Command Information Setting**

Synchronization Correction Mode	D01	D02H	D02L	D03	D04
Microcontroller correction mode	00H	Fixed to 00H	Fixed to 0AH (115,200 bps)	Noise filter 00H: Off	00H: When writing in full-speed mode (2.7 V ≤ V <sub>DD</sub> ≤ 5.5 V)
Programmer correction mode	01H	<b>Note 2</b>	<b>Note 2</b>	01H: On	01H: When writing in wide-voltage mode <sup>Note 1</sup> (1.8 V ≤ V <sub>DD</sub> ≤ 5.5 V)

**Notes 1.** Support of the wide-voltage mode differs depending on the products. Check if it is supported by referring to the user's manual of each product.

- 2.** To communicate in the programmer correction mode, set a value to D02H/D02L by executing the following procedures using the programmer.

<1> Measure the length of the READY pulse output from the microcontroller.

<2> Calculate errors of the READY pulse (low-level 9 bits @ 9,600 bps).

<3> Calculate the k value by using the following equation, taking errors in consideration.

$$k = (f_{IH} \times E) \div (\text{Baud rate} \times 2)$$

E: READY pulse (9,600 bps) error when the flash memory programming mode has been set

**Example:** E = 1.00 when 0% error for READY pulse (low-level 9 bits @ 9,600 bps) length

E = 1.02 when +2% error for READY pulse (low-level 9 bits @ 9,600 bps) length

E = 0.98 when -2% error for READY pulse (low-level 9 bits @ 9,600 bps) length

f<sub>IH</sub>: High-speed internal oscillator frequency

<4> Convert the k value to a hexadecimal number.

<5> Set the value obtained in step <4> to D02H/D02L, and then execute the Baud Rate Set command.

However, steps <1> to <3> can be omitted if the following condition is satisfied.

**Condition:** All the k values are the same as a result of calculating with consideration for the deviation of the device's operating frequency. (The k value is rounded to a whole number.)

**Example:** If specifying 1 Mbps with  $f_{IH} = 20 \text{ MHz} \pm 2\%$ ,  $k = 10$  is obtained by using the following equations (D02H/D02L = 00H/0AH).

- 20.4 MHz (20 MHz + 2%)  
 $k = (f_{IH} \times E) \div (\text{Baud rate} \times 2)$   
 $= (20.4 \times 10^6) \div (1 \times 10^6 \times 2)$   
 $= 10.2$   
 $\cong 10$
- 20 MHz (20 MHz  $\pm 0\%$ )  
 $k = (f_{IH} \times E) \div (\text{Baud rate} \times 2)$   
 $= (20 \times 10^6) \div (1 \times 10^6 \times 2)$   
 $= 10$
- 19.6 MHz (20 MHz - 2%)  
 $k = (f_{IH} \times E) \div (\text{Baud rate} \times 2)$   
 $= (19.6 \times 10^6) \div (1 \times 10^6 \times 2)$   
 $= 9.8$   
 $\cong 10$

Figure 3-4. Status Frame for Baud Rate Set Command (from 78K0R/Kx3-L to Programmer)

STX	LEN	Data	SUM	ETX
02H	01H	ST1	checksum	03H

**Remark** ST1: Synchronization detection result

See 4.5 **Baud Rate Set Command** for details on the flowchart of the processing sequence between the programmer and the 78K0R/Kx3-L, the flowchart of command processing, and the sample program.

### 3.3 Chip Erase Command

#### 3.3.1 Description

This command is used to erase the entire contents of the flash memory. In addition, all of the information that is set by security setting processing can be initialized by chip erase processing, as long as erasure is not prohibited by the security setting (see 3.11 Security Set Command).

#### 3.3.2 Command frame and status frame

Figure 3-5 shows the format of a command frame for the Chip Erase command, and Figure 3-6 shows the status frame for the command.

**Figure 3-5. Chip Erase Command Frame (from Programmer to 78K0R/Kx3-L)**

SOH	LEN	COM	SUM	ETX
01H	01H	20H (Chip Erase)	Checksum	03H

**Figure 3-6. Status Frame for Chip Erase Command (from 78K0R/Kx3-L to Programmer)**

STX	LEN	Data	SUM	ETX
02H	01H	ST1	Checksum	03H

**Remark** ST1: Chip erase result

See 4.6 Chip Erase Command for details on the flowchart of the processing sequence between the programmer and the 78K0R/Kx3-L, the flowchart of command processing, and the sample program.



### 3.4 Block Erase Command

#### 3.4.1 Description

This command is used to erase the content of flash memory of the block with the specified number.

A block can be specified with the first address of the block where erasing starts and the last address where erasing ends. Successive multiple blocks can be specified.

Erasing cannot be performed, however, if erasing is prohibited due to the security setting (see **3.11 Security Set Command**).

#### 3.4.2 Command frame and status frame

Figure 3-7 shows the format of a command frame for the Block Erase command, and Figure 3-8 shows the status frame for the command.

**Figure 3-7. Block Erase Command Frame (from Programmer to 78K0R/Kx3-L)**

SOH	LEN	COM	Command Information						SUM	ETX
01H	07H	22H (Block Erase)	SAH	SAM	SAL	EAH	EAM	EAL	Checksum	03H

**Remark** SAH, SAM, SAL: Block erase start address (start address of any block)

SAH: Start address, high (bits 23 to 16)

SAM: Start address, middle (bits 15 to 8)

SAL: Start address, low (bits 7 to 0)

EAH, EAM, EAL: Block erase end address (last address of any block)

EAH: End address, high (bits 23 to 16)

EAM: End address, middle (bits 15 to 8)

EAL: End address, low (bits 7 to 0)

**Figure 3-8. Status Frame for Block Erase Command (from 78K0R/Kx3-L to Programmer)**

STX	LEN	Data	SUM	ETX
02H	01H	ST1	Checksum	03H

**Remark** ST1: Block erase result

See **4.7 Block Erase Command** for details on the flowchart of the processing sequence between the programmer and the 78K0R/Kx3-L, the flowchart of command processing, and the sample program.

### 3.5 Programming Command

#### 3.5.1 Description

This command is used to write the user program to the flash memory by transmitting write data after having transmitted the write start address and the write end address. Internal verification is then executed after the last data has been transmitted and writing has been completed.

The write start/end address can be set only in the block start/end address units.

If both of the status frames (ST1 and ST2) after the last data transmission indicate ACK, the 78K0R/Kx3-L firmware automatically executes internal verify. Therefore, the Status command for this internal verify must be transmitted.

#### 3.5.2 Command frame and status frame

Figure 3-9 shows the format of a command frame for the Programming command, and Figure 3-10 shows the status frame for the command.

**Figure 3-9. Programming Command Frame (from Programmer to 78K0R/Kx3-L)**

SOH	LEN	COM	Command Information						SUM	ETX
01H	07H	40H (Programming)	SAH	SAM	SAL	EAH	EAM	EAL	Checksum	03H

**Remark** SAH, SAM, SAL: Write start addresses  
EAH, EAM, EAL: Write end addresses

**Figure 3-10. Status Frame for Programming Command (from 78K0R/Kx3-L to Programmer)**

STX	LEN	Data	SUM	ETX
02H	01H	ST1 (a)	Checksum	03H

**Remark** ST1 (a): Command reception result

#### 3.5.3 Data frame and status frame

Figure 3-11 shows the format of a frame that includes data to be written, and Figure 3-12 shows the status frame for the data.

**Figure 3-11. Data Frame to Be Written (from Programmer to 78K0R/Kx3-L)**

STX	LEN	Data	SUM	ETX/ETB
02H	00H to FFH (00H = 256)	Write Data	Checksum	03H/17H

**Remark** Write Data: User program to be written

**Figure 3-12. Status Frame for Data Frame (from 78K0R/Kx3-L to Programmer)**

STX	LEN	Data		SUM	ETX
02H	02H	ST1 (b)	ST2 (b)	Checksum	03H

**Remark** ST1 (b): Data reception check result  
ST2 (b): Write result

### 3.5.4 Completion of transferring all data and status frame

Figure 3-13 shows the status frame after transfer of all data is completed.

**Figure 3-13. Status Frame After Completion of Transferring All Data (from 78K0R/Kx3-L to Programmer)**

STX	LEN	Data	SUM	ETX
02H	01H	ST1 (c)	Checksum	03H

**Remark** ST1 (c): Internal verify result

See **4.8 Programming Command** for details on the flowchart of the processing sequence between the programmer and the 78K0R/Kx3-L, the flowchart of command processing, and the sample program.

## 3.6 Verify Command

### 3.6.1 Description

This command is used to compare the data transmitted from the programmer with the data read from the 78K0R/Kx3-L (read level) in the specified address range, and check whether they match.

The verify start/end address can be set only in the block start/end address units.

### 3.6.2 Command frame and status frame

Figure 3-14 shows the format of a command frame for the Verify command, and Figure 3-15 shows the status frame for the command.

**Figure 3-14. Verify Command Frame (from Programmer to 78K0R/Kx3-L)**

SOH	LEN	COM	Command Information						SUM	ETX
01H	07H	13H (Verify)	SAH	SAM	SAL	EAH	EAM	EAL	Checksum	03H

**Remark** SAH, SAM, SAL: Verify start addresses

EAH, EAM, EAL: Verify end addresses

**Figure 3-15. Status Frame for Verify Command (from 78K0R/Kx3-L to Programmer)**

STX	LEN	Data	SUM	ETX
02H	01H	ST1 (a)	Checksum	03H

**Remark** ST1 (a): Command reception result

### 3.6.3 Data frame and status frame

Figure 3-16 shows the format of a frame that includes data to be verified, and Figure 3-17 shows the status frame for the data.

**Figure 3-16. Data Frame of Data to Be Verified (from Programmer to 78K0R/Kx3-L)**

STX	LEN	Data	SUM	ETX/ETB
02H	00H to FFH (00H = 256)	Verify Data	Checksum	03H/17H

**Remark** Verify Data: User program to be verified

**Figure 3-17. Status Frame for Data Frame (from 78K0R/Kx3-L to Programmer)**

STX	LEN	Data		SUM	ETX
02H	02H	ST1 (b)	ST2 (b)	Checksum	03H

**Remark** ST1 (b): Data reception check result  
ST2 (b): Verify result<sup>Note</sup>

**Note** Even if a verify error occurs in the specified address range, ACK is always returned as the verify result. The status of all verify errors are reflected in the verify result for the last data. Therefore, the occurrence of verify errors can be checked only when all the verify processing for the specified address range is completed.

See **4.9 Verify Command** for details on the flowchart of the processing sequence between the programmer and the 78K0R/Kx3-L, the flowchart of command processing, and the sample program.

### 3.7 Block Blank Check Command

#### 3.7.1 Description

This command is used to check if a block in the flash memory, with a specified block number, is blank (erased state).

A block can be specified with the start address of the blank check start block and the last address of the blank check end block. Successive multiple blocks can be specified.

#### 3.7.2 Command frame and status frame

Figure 3-18 shows the format of a command frame for the Block Blank Check command, and Figure 3-19 shows the status frame for the command.

**Figure 3-18. Block Blank Check Command Frame (from Programmer to 78K0R/Kx3-L)**

SOH	LEN	COM	Command Information							SUM	ETX
01H	08H	32H (Block Blank Check)	SAH	SAM	SAL	EAH	EAM	EAL	D01	Checksum	03H

**Remark** SAH, SAM, SAL: Block blank check start address (start address of any block)

SAH: Start address, high (bits 23 to 16)

SAM: Start address, middle (bits 15 to 8)

SAL: Start address, low (bits 7 to 0)

EAH, EAM, EAL: Block blank check end address (last address of any block)

EAH: End address, high (bits 23 to 16)

EAM: End address, middle (bits 15 to 8)

EAL: End address, low (bits 7 to 0)

D01: 00H: When performing a block blank check for a single block

01H: When performing a blank check for the complete area before erasing the chip

**Figure 3-19. Status Frame for Block Blank Check Command (from 78K0R/Kx3-L to Programmer)**

STX	LEN	Data	SUM	ETX
02H	01H	ST1	Checksum	03H

**Remark** ST1: Block blank check result

See **4.10 Block Blank Check Command** for details on the flowchart of the processing sequence between the programmer and the 78K0R/Kx3-L, the flowchart of command processing, and the sample program.

### 3.8 Silicon Signature Command

#### 3.8.1 Description

This command is used to read information such as the write protocol information (silicon signature) of the device and security flag information.

If the programmer supports a programming protocol that is not supported in the 78K0R/Kx3-L, for example, execute this command to select an appropriate protocol in accordance with the values of the second and third bytes.

#### 3.8.2 Command frame and status frame

Figure 3-20 shows the format of a command frame for the Silicon Signature command, and Figure 3-21 shows the status frame for the command.

**Figure 3-20. Silicon Signature Command Frame (from Programmer to 78K0R/Kx3-L)**

SOH	LEN	COM	SUM	ETX
01H	01H	C0H (Silicon Signature)	Checksum	03H

**Figure 3-21. Status Frame for Silicon Signature Command (from 78K0R/Kx3-L to Programmer)**

STX	LEN	Data	SUM	ETX
02H	01H	ST1	Checksum	03H

**Remark** ST1: Command reception result

### 3.8.3 Silicon signature data frame

Figure 3-22 shows the format of a frame that includes silicon signature data.

**Figure 3-22. Silicon Signature Data Frame (from 78K0R/Kx3-L to Programmer)**

STX	LEN	Data							
02H	n	VEN	MET	MSC	DEC1	DEC2	DEC3	UAE(3)	DEV(10)

Data (continued)							SUM	ETX
SCF	BOT	FSWSH	FSWSL	FSWEH	FSWEL	RES(2)	Checksum	03H

- Remarks 1.**
- n (LEN): Data length
  - VEN: Vendor code (NEC: 10H)
  - MET: Macro extension code
  - MSC: Macro function code
  - DEC1: Device extension code 1
  - DEC2: Device extension code 2
  - DEC3: Device extension code 3
  - UAE: User flash ROM last address (3 bytes)
  - DEV: Device name (10 bytes)
  - SCF: Security flag information
  - BOT: Boot block number
  - FSWSH: Higher 8-bit side of flash shield window (FSW) start block
  - FSWSL: Lower 8-bit side of flash shield window (FSW) start block
  - FSWEH: Higher 8-bit side of flash shield window (FSW) end block
  - FSWEL: Lower 8-bit side of flash shield window (FSW) end block
  - RES: Reserved (2 bytes)
- 2.** For the vendor code (VEN), extension code (MET), function code (MSC), device extension code 1 (DEC1), device extension code 2 (DEC2), and device extension code 3 (DEC3), the lower 7 bits are used as data entity, and the highest bit is used as an odd parity. The following shows an example.



Table 3-2. Example of Silicon Signature Data ( $\mu$ PD78F1000 (78K0R/KC3-L))

Field Name	Content	Length (Byte)	Example of Silicon Signature Data	Actual Value	Parity
VEN	Vendor code (NEC)	1	10H (00010000B)	10H	Added
MET	Macro extension code	1	EFH (11101111B)	EFH	Added
MSC	Macro function code	1	04H (01000000B)	04H	Added
DEC1	Device extension code 1	1	DCH (11011100B)	DCH	Added
DEC2	Device extension code 2	1	FDH (11111101B)	FDH	Added
DEC3	Device extension code 3	1	FDH (11111101B)	FDH	Added
UAE	User flash ROM last address	3	FFH (11111111B)	003FFFH	Not added
			3FH (00111111B)		
			00H (00000000B)		
DEV	Device name	10	44H (01000100B) = 'D'	'D'	Not added
			37H (00110111B) = '7'	'7'	
			38H (00111000B) = '8'	'8'	
			46H (01000110B) = 'F'	'F'	
			31H (00110001B) = '1'	'1'	
			30H (00110000B) = '0'	'0'	
			30H (00110000B) = '0'	'0'	
			30H (00110000B) = '0'	'0'	
			20H (00100000B) = ''	''	
			20H (00100000B) = ''	''	
SCF	Security flag information	1	Any	Same as left column	Not added
BOT	Boot block number (fixed)	1	03H (00000011B)	03H	Not added
FSWS(H)	Higher 8-bit side of flash shield window start block	1	Any	Same as left column	Not added
FSWS(L)	Lower 8-bit side of flash shield window start block	1	Any	Same as left column	Not added
FSWE(H)	Higher 8-bit side of flash shield window end block	1	Any	Same as left column	Not added
FSWE(L)	Lower 8-bit side of flash shield window end block	1	Any	Same as left column	Not added
RES	Reserved	2	FFFFH (1111111111111111B)	FFFFH	Not added

See 4.11 **Silicon Signature Command** for details on the flowchart of the processing sequence between the programmer and the 78K0R/Kx3-L, the flowchart of command processing, and the sample program.

## 3.8.4 Silicon signature list

## (1) 78K0R/Kx3-L silicon signature list

Table 3-3. 78K0R/Kx3-L Silicon Signature Data List

Item	Description	Length (Byte)	Data (Hex)
Vendor code	NEC	1	10
Extension code	Extension code	1	7F
Function code	Function information	1	04
Device information	Device information	3	DC
			FD
			FD
Internal flash ROM last address	Transmitted from lower bytes of address	3	<b>Note 1</b>
Device name ( $\mu$ PD)	78F1000/78F1001/78F1002/78F1003 78F1004/78F1005/78F1006 78F1007/78F1008/78F1009 78F1010/78F1011/78F1012 78F1013/78F1014	10	<b>Note 2</b>
Security information	Security information	1	Any
Boot block number	The last block number of the boot cluster that is currently selected	1	03
FSW block number	FSW information	4	Any
Reserved	Reserved	2	FFFF

**Note 1.** The list of internal-flash-ROM last addresses is as follows.

Item	Description	Length (Byte)	Data (Hex)
Internal flash ROM last address	16 KB (003FFFH)	3	FF3F00
	32 KB (007FFFH)		FF7F00
	48 KB (00BFFFH)		FFBF00
	64 KB (00FFFFH)		FFFF00
	96 KB (017FFFH)		FF7F01
	128 KB (01FFFFH)		FFFF01

(Note 2 is on the next page.)

**Note 2.** The device names are listed below.

**Device name list**

Generic Name	Part Number	Length (Bytes)	Actual Value										
			( Upper Row: Signature Data Lower Row: Character Code )										
78K0R/KC3-L	D78F1000	10	44	37	38	46	31	30	30	30	20	20	
			D	7	8	F	1	0	0	0	-	-	
	D78F1001		44	37	38	46	31	30	30	31	20	20	
			D	7	8	F	1	0	0	1	-	-	
	D78F1002		44	37	38	46	31	30	30	32	20	20	
			D	7	8	F	1	0	0	2	-	-	
	D78F1003		44	37	38	46	31	30	30	33	20	20	
			D	7	8	F	1	0	0	3	-	-	
	78K0R/KD3-L		D78F1004	44	37	38	46	31	30	30	34	20	20
				D	7	8	F	1	0	0	4	-	-
			D78F1005	44	37	38	46	31	30	30	35	20	20
				D	7	8	F	1	0	0	5	-	-
D78F1006		44	37	38	46	31	30	30	36	20	20		
		D	7	8	F	1	0	0	6	-	-		
78K0R/KE3-L	D78F1007	44	37	38	46	31	30	30	37	20	20		
		D	7	8	F	1	0	0	7	-	-		
	D78F1008	44	37	38	46	31	30	30	38	20	20		
		D	7	8	F	1	0	0	8	-	-		
	D78F1009	44	37	38	46	31	30	30	39	20	20		
		D	7	8	F	1	0	0	9	-	-		
78K0R/KF3-L	D78F1010	44	37	38	46	31	30	31	30	20	20		
		D	7	8	F	1	0	1	0	-	-		
	D78F1011	44	37	38	46	31	30	31	31	20	20		
		D	7	8	F	1	0	1	1	-	-		
	D78F1012	44	37	38	46	31	30	31	32	20	20		
		D	7	8	F	1	0	1	2	-	-		
78K0R/KG3-L	D78F1013	44	37	38	46	31	30	31	33	20	20		
		D	7	8	F	1	0	1	3	-	-		
	D78F1014	44	37	38	46	31	30	31	34	20	20		
		D	7	8	F	1	0	1	4	-	-		

## (2) 78K0R/lx3 silicon signature list

Table 3-4. 78K0R/lx3 Silicon Signature Data List

Item	Description	Length (Byte)	Data (Hex)
Vendor code	NEC	1	10
Extension code	Extension code	1	7F
Function code	Function information	1	04
Device information	Device information	3	DC
			FD
			FD
Internal flash ROM last address	Transmitted from lower bytes of address	3	<b>Note 1</b>
Device name ( $\mu$ PD)	78F1201/78F1203 78F1211/78F1213/78F1214/78F1215 78F1223/78F1224/78F1225 78F1233/78F1234/78F1235	10	<b>Note 2</b>
Security information	Security information	1	Any
Boot block number	The last block number of the boot cluster that is currently selected	1	03
FSW block number	FSW information	4	Any
Reserved	Reserved	2	FFFF

**Note 1.** The list of internal-flash-ROM last addresses is as follows.

Item	Description	Length (Byte)	Data (Hex)
Internal flash ROM last address	16 KB (003FFFH)	3	FF3F00
	32 KB (007FFFH)		FF7F00
	48 KB (00BFFFH)		FFBF00
	64 KB (00FFFFH)		FFFF00

(Note 2 is on the next page.)

**Note 2.** The device names are listed below.

**Device name list**

Generic Name	Part Number	Length (Bytes)	Actual Value									
			( Upper Row: Signature Data Lower Row: Character Code )									
78K0R/IB3	D78F1221	10	44	37	38	46	31	32	32	31	20	20
			D	7	8	F	1	2	2	1	-	-
D78F1222	44		37	38	46	31	32	32	32	20	20	
			D	7	8	F	1	2	2	2	-	-
78K0R/IC3	D78F1211		44	37	38	46	31	32	31	31	20	20
			D	7	8	F	1	2	1	1	-	-
	D78F1213		44	37	38	46	31	32	31	33	20	20
				D	7	8	F	1	2	1	3	-
	D78F1214		44	37	38	46	31	32	31	34	20	20
				D	7	8	F	1	2	1	4	-
D78F1215	44	37	38	46	31	32	31	35	20	20		
		D	7	8	F	1	2	1	5	-	-	
78K0R/ID3	D78F1223	44	37	38	46	31	32	32	33	20	20	
			D	7	8	F	1	2	2	3	-	-
	D78F1224	44	37	38	46	31	32	32	34	20	20	
			D	7	8	F	1	2	2	4	-	-
	D78F1225	44	37	38	46	31	32	32	35	20	20	
	D	7	8	F	1	2	2	5	-	-		
78K0R/IE3	D78F1233	44	37	38	46	31	32	33	33	20	20	
			D	7	8	F	1	2	3	3	-	-
	D78F1234	44	37	38	46	31	32	33	34	20	20	
			D	7	8	F	1	2	3	4	-	-
	D78F1235	44	37	38	46	31	32	33	35	20	20	
	D	7	8	F	1	2	3	5	-	-		

(3) 78K0R/Kx3-C silicon signature list

Table 3-5. 78K0R/Kx3-C Silicon Signature Data List

Item	Description	Length (Byte)	Data (Hex)
Vendor code	NEC	1	10
Extension code	Extension code	1	7F
Function code	Function information	1	04
Device information	Device information	3	DC
			FD
			FD
Internal flash ROM last address	Transmitted from lower bytes of address	3	<b>Note 1</b>
Device name (μPD)	78F1846/78F1847 78F1848/78F1849	10	<b>Note 2</b>
Security information	Security information	1	Any
Boot block number	The last block number of the boot cluster that is currently selected	1	03
FSW block number	FSW information	4	Any
Reserved	Reserved	2	FFFF

**Notes 1.** The list of internal-flash-ROM last addresses is as follows.

Item	Description	Length (Byte)	Data (Hex)
Internal flash ROM last address	96 KB (017FFFH)	3	FF7F01
	128 KB (01FFFFH)		FFFF01

**2.** The device names are listed below.

**Device name list**

Generic Name	Part Number	Length (Bytes)	Actual Value									
			( Upper Row: Signature Data Lower Row: Character Code )									
78K0R/KF3-C	D78F1846	10	44	37	38	46	31	38	34	36	20	20
			D	7	8	F	1	8	4	6	-	-
	D78F1847		44	37	38	46	31	38	34	37	20	20
			D	7	8	F	1	8	4	7	-	-
78K0R/KG3-C	D78F1848	10	44	37	38	46	31	38	34	38	20	20
			D	7	8	F	1	8	4	8	-	-
	D78F1849		44	37	38	46	31	38	34	39	20	20
			D	7	8	F	1	8	4	9	-	-

### 3.9 Version Get Command

#### 3.9.1 Description

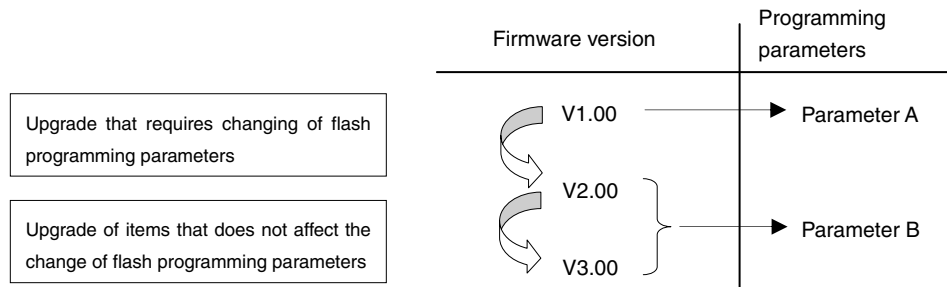
This command is used to acquire information on the 78K0R/Kx3-L device version and firmware version.

The device version value is fixed to 00H.

Use this command when the programming parameters must be changed in accordance with the 78K0R/Kx3-L firmware version.

**Caution** The firmware version may be updated during firmware update that does not affect the change of flash programming parameters (at this time, update of the firmware version is not reported).

**Example** Firmware version and reprogramming parameters



#### 3.9.2 Command frame and status frame

Figure 3-23 shows the format of a command frame for the Version Get command, and Figure 3-24 shows the status frame for the command.

**Figure 3-23. Version Get Command Frame (from Programmer to 78K0R/Kx3-L)**

SOH	LEN	COM	SUM	ETX
01H	01H	C5H (Version Get)	Checksum	03H

**Figure 3-24. Status Frame for Version Get Command (from 78K0R/Kx3-L to Programmer)**

STX	LEN	Data	SUM	ETX
02H	01H	ST1	Checksum	03H

**Remark** ST1: Command reception result

### 3.9.3 Version data frame

Figure 3-25 shows the data frame of version data.

**Figure 3-25. Version Data Frame (from 78K0R/Kx3-L to Programmer)**

STX	LEN	Data						SUM	ETX
02H	06H	DV1	DV2	DV3	FV1	FV2	FV3	Checksum	03H

**Remark** DV1: Integer of device version (fixed to 00H)  
DV2: First decimal place of device version (fixed to 00H)  
DV3: Second decimal place of device version (fixed to 00H)  
FV1: Integer of firmware version  
FV2: First decimal place of firmware version  
FV3: Second decimal place of firmware version

See **4.12 Version Get Command** for details on the flowchart of the processing sequence between the programmer and the 78K0R/Kx3-L, the flowchart of command processing, and the sample program.



### 3.10 Checksum Command

#### 3.10.1 Description

This command is used to acquire the checksum data in the specified area.

For the checksum calculation start/end address, specify a fixed address in block units (1 KB) starting from the top of the flash memory.

Checksum data is obtained by sequentially subtracting data in the specified address range from the initial value (0000H) in 1-byte units.

#### 3.10.2 Command frame and status frame

Figure 3-26 shows the format of a command frame for the Checksum command, and Figure 3-27 shows the status frame for the command.

**Figure 3-26. Checksum Command Frame (from Programmer to 78K0R/Kx3-L)**

SOH	LEN	COM	Command Information						SUM	ETX
01H	07H	B0H (Checksum)	SAH	SAM	SAL	EAH	EAM	EAL	Checksum	03H

**Remark** SAH, SAM, SAL: Checksum calculation start addresses  
EAH, EAM, EAL: Checksum calculation end addresses

**Figure 3-27. Status Frame for Checksum Command (from 78K0R/Kx3-L to Programmer)**

STX	LEN	Data	SUM	ETX
02H	01H	ST1	Checksum	03H

**Remark** ST1: Command reception result

#### 3.10.3 Checksum data frame

Figure 3-28 shows the format of a frame that includes checksum data.

**Figure 3-28. Checksum Data Frame (from 78K0R/Kx3-L to Programmer)**

STX	LEN	Data		SUM	ETX
02H	02H	CK1	CK2	Checksum	03H

**Remark** CK1: Higher 8 bits of checksum data  
CK2: Lower 8 bits of checksum data

See 4.13 **Checksum Command** for details on the flowchart of the processing sequence between the programmer and the 78K0R/Kx3-L, the flowchart of command processing, and the sample program.

### 3.11 Security Set Command

#### 3.11.1 Description

This command is used to perform security settings (enabling/disabling of write, block erase, chip erase, and boot block rewriting, and setting of flash shield window start/end block number). By performing these settings with this command, rewriting of the flash memory by an unauthorized party can be restricted and the rewrite area for self programming can be specified.

**Caution** Even after the security setting, additional setting of changing from enable to disable can be performed; however, changing from disable to enable is not possible. If an attempt is made to perform such a setting, a protect error (10H) will occur. If such setting is required, all of the security flags must first be initialized by executing the Chip Erase command (the Block Erase command cannot be used to initialize the security flags).

If chip erase or boot block rewrite has been disabled, however, chip erase itself will be impossible, so the settings cannot be erased from the programmer. Re-confirmation of security setting execution is therefore recommended before disabling chip erase, due to this programmer specification.

#### 3.11.2 Command frame and status frame

Figure 3-29 shows the format of a command frame for the Security Set command, and Figure 3-30 shows the status frame for the command.

**Figure 3-29. Security Set Command Frame (from Programmer to 78K0R/Kx3-L)**

SOH	LEN	COM	Command Information		SUM	ETX
01H	03H	A0H (Security Set)	00H (fixed)	00H (fixed)	Checksum	03H

**Figure 3-30. Status Frame for Security Set Command (from 78K0R/Kx3-L to Programmer)**

STX	LEN	Data	SUM	ETX
02H	01H	ST1 (a)	Checksum	03H

**Remark** ST1 (a): Command reception result

3.11.3 Data frame and status frame

Figure 3-31 shows the format of a security data frame, and Figure 3-32 shows the status frame for the data.

Figure 3-31. Security Data Frame (from Programmer to 78K0R/Kx3-L)

STX	LEN	Data								SUM	ETX
02H	06H	FLG	BOT	FSWSH	FSWSL	FSWEH	FSWEL	FFH (fixed)	FFH (fixed)	Checksum	03H

- Remarks 1.**
- FLG: Security flag
  - BOT: Boot cluster last block number (fixed to 03H)
  - FSWSH: Higher 8 bits of flash shield window start block number
  - FSWSL: Lower 8 bits of flash shield window start block number
  - FSWEH: Higher 8 bits of flash shield window end block number
  - FSWEL: Lower 8 bits of flash shield window end block number
- 2.** If the flash shield window is not to be set, set FSWS to 0000H and the end block to the target device end block number.

Figure 3-32. Status Frame for Security Data Writing (from 78K0R/Kx3-L to Programmer)

STX	LEN	Data	SUM	ETX
02H	01H	ST1 (b)	Checksum	03H

**Remark** ST1 (b): Security data write result

3.11.4 Internal verify check and status frame

Figure 3-33 shows the status frame for internal verify check.

Figure 3-33. Status Frame for Internal Verify Check (from 78K0R/Kx3-L to Programmer)

STX	LEN	Data	SUM	ETX
02H	01H	ST1 (c)	Checksum	03H

**Remark** ST1 (c): Internal verify result

The following table shows the contents in the security flag field.

Table 3-6. Contents of Security Flag Field

Item	Contents
Bit 7	Fixed to "1"
Bit 6	
Bit 5	
Bit 4	Boot block rewrite disable flag (1: Enables boot block rewrite, 0: Disable boot block rewrite)
Bit 3	Fixed to "1"
Bit 2	Programming disable flag (1: Enables programming, 0: Disable programming)
Bit 1	Block erase disable flag (1: Enables block erase, 0: Disable block erase)
Bit 0	Chip erase disable flag (1: Enables chip erase, 0: Disable chip erase)

The following table shows the relationship between the security flag field settings and the enable/disable status of each operation.

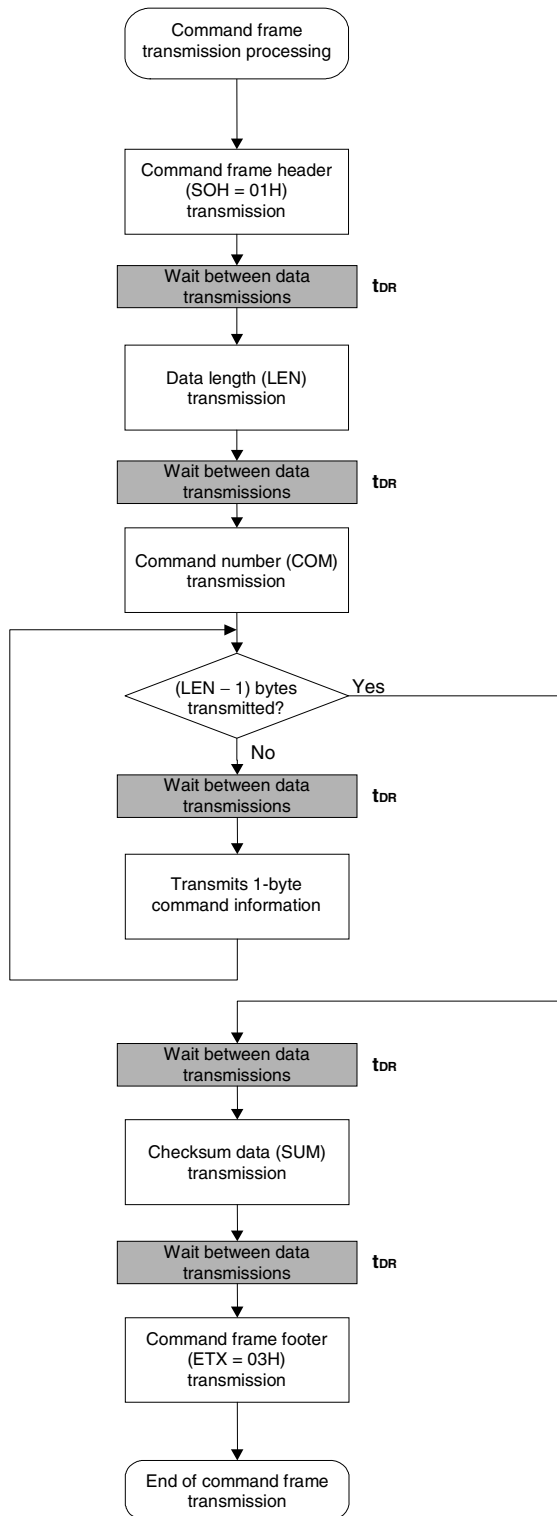
**Table 3-7. Security Flag Field and Enable/Disable Status of Each Operation**

Operating Mode	Flash Memory Programming Mode			Self-Programming Mode
Security Setting Item	Command Operation After Security Setting √: Execution possible, ×: Execution impossible △: Writing and block erase in boot area are impossible			<ul style="list-style-type: none"> <li>All commands can be executed regardless of the security setting values</li> <li>Only retention of security setting values is possible</li> </ul>
	Programming	Chip Erase	Block Erase	
Disable programming	×	√	×	
Disable chip erase	√	×	×	
Disable block erase	√	√	×	
Boot block rewrite disable flag	△	×	△	Same condition as that in flash memory programming mode (on-board/off-board programming)

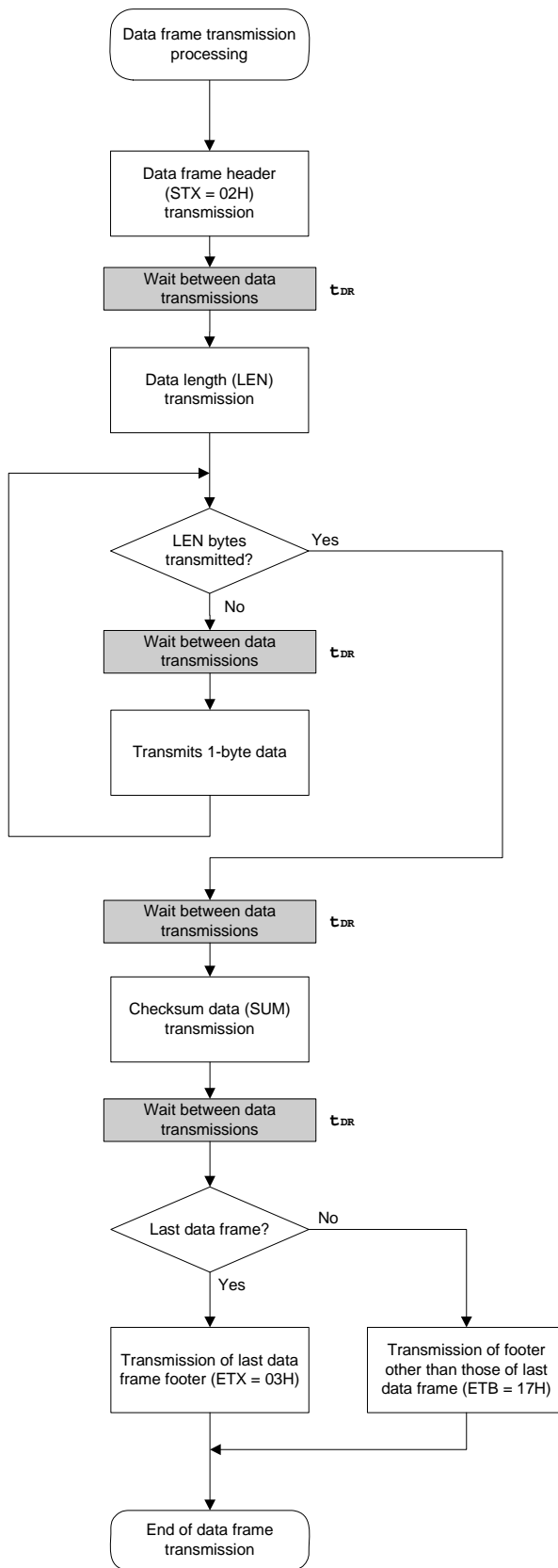
See **4.14 Security Set Command** for details on the flowchart of the processing sequence between the programmer and the 78K0R/Kx3-L, the flowchart of command processing, and the sample program.

## CHAPTER 4 UART COMMUNICATION MODE

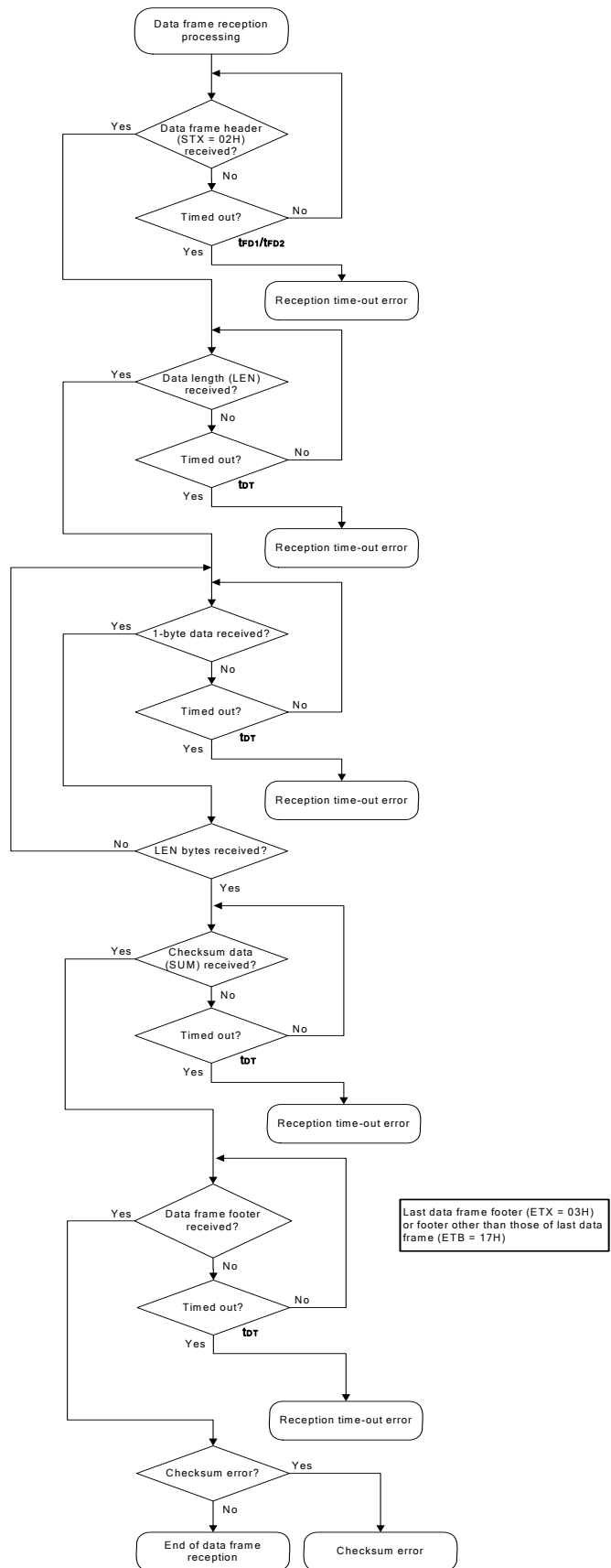
### 4.1 Command Frame Transmission Processing Flowchart



4.2 Data Frame Transmission Processing Flowchart



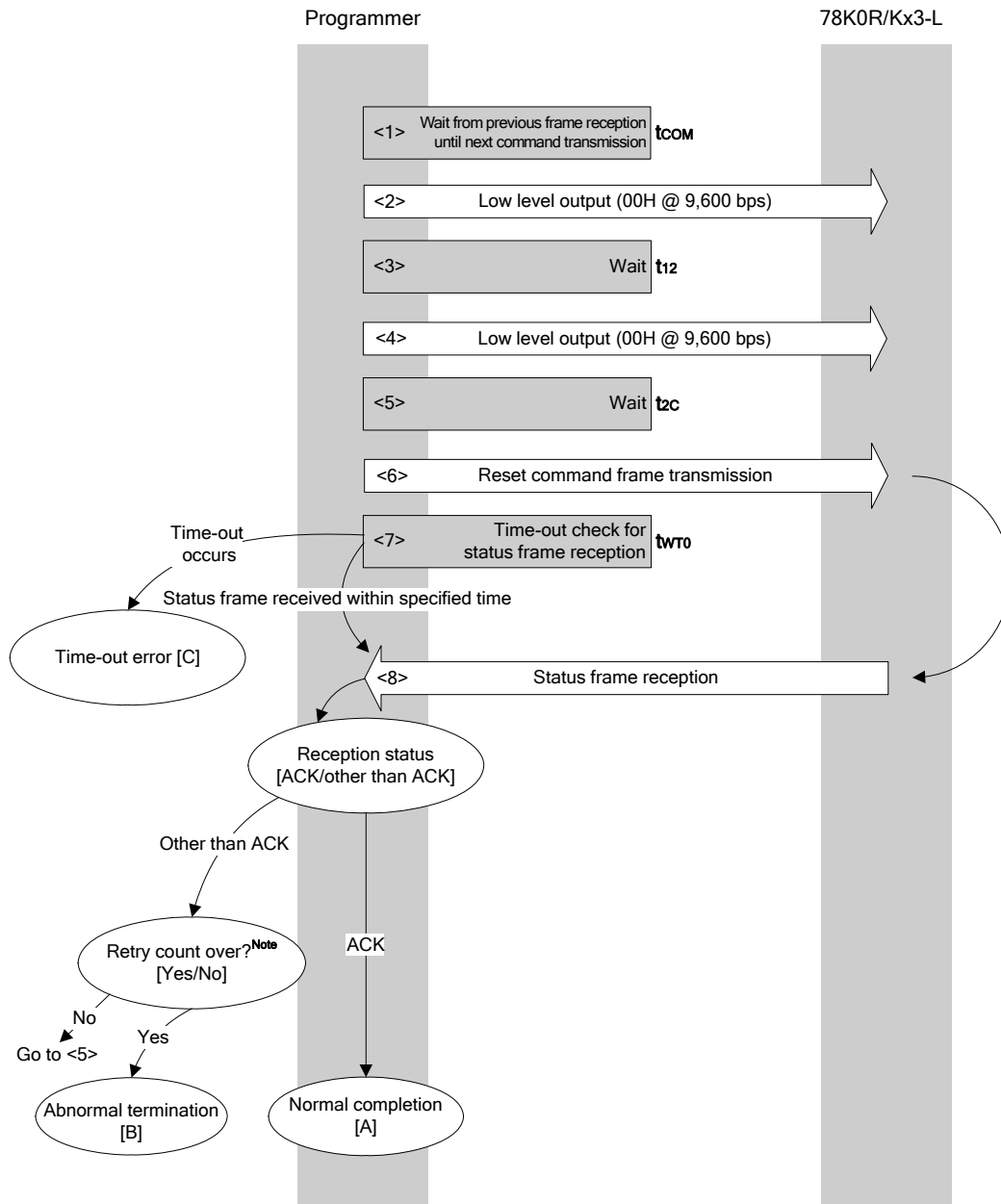
### 4.3 Data Frame Reception Processing Flowchart



4.4 Reset Command

4.4.1 Processing sequence chart

Reset command processing sequence



**Note** Do not exceed the retry count for the reset command transmission (up to 16 times).



#### 4.4.2 Description of processing sequence

- <1> Waits from the previous frame reception until the next command processing starts (wait time  $t_{COM}$ ).
- <2> The low level is output (data 00H is transmitted at 9,600 bps).
- <3> Wait state (wait time  $t_{12}$ ).
- <4> The low level is output (data 00H is transmitted at 9,600 bps).
- <5> Wait state (wait time  $t_{2C}$ ).
- <6> The Reset command is transmitted by command frame transmission processing.
- <7> A time-out check is performed from command transmission until status frame reception.  
If a time-out occurs, a time-out error [C] is returned (time-out time  $t_{WTO}$ ).
- <8> The status code is checked.

When ST1 = ACK: Normal completion [A]

When ST1  $\neq$  ACK: The retry count ( $t_{RS}$ ) is checked.

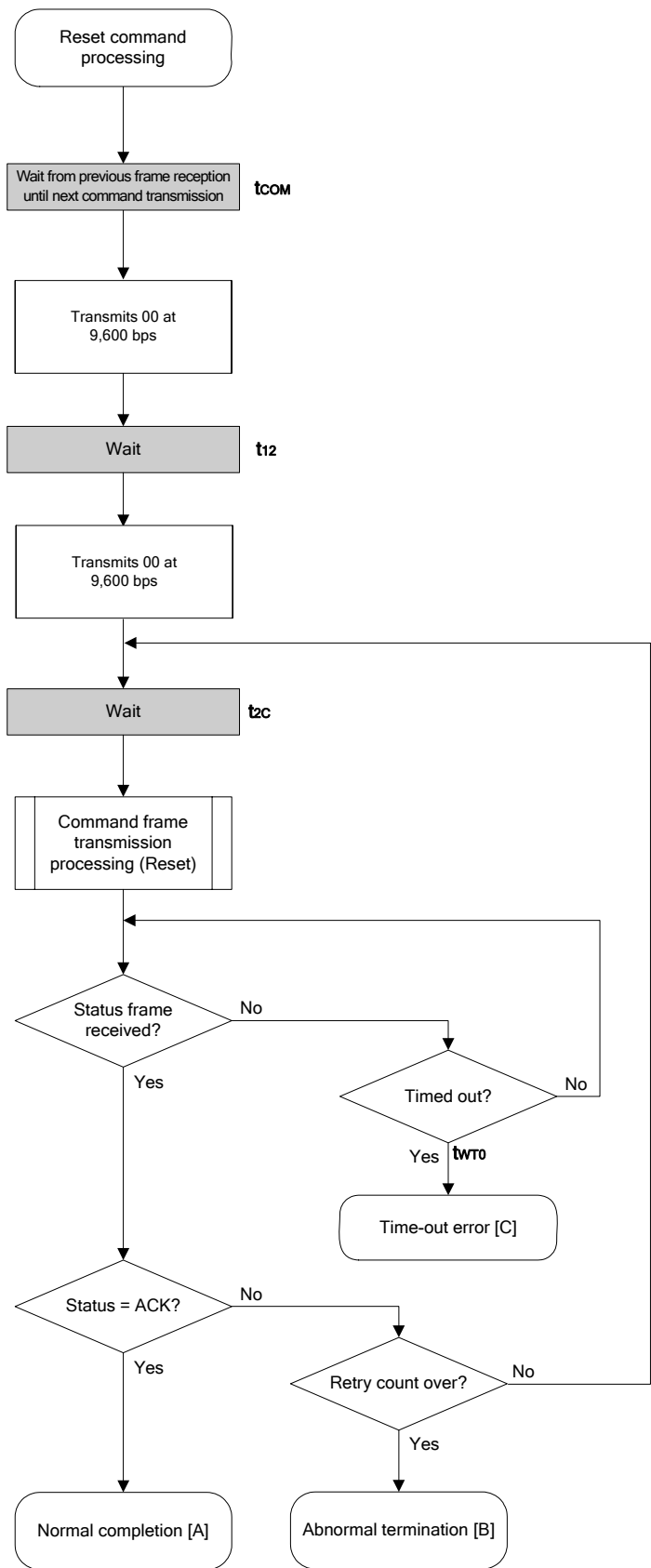
The sequence is re-executed from <5> if the retry count is not over.

If the retry count is over, the processing ends abnormally [B].

#### 4.4.3 Status at processing completion

Status at Processing Completion		Status Code	Description
Normal completion [A]	Normal acknowledgment (ACK)	06H	The command was executed normally and synchronization between the programmer and the 78K0R/Kx3-L has been established.
Abnormal termination [B]	Checksum error	07H	The checksum of the transmitted command frame does not match.
	Negative acknowledgment (NACK)	15H	Command frame data is abnormal (such as invalid data length (LEN) or no ETX).
Time-out error [C]		–	The status frame was not received within the specified time.

4.4.4 Flowchart



#### 4.4.5 Sample program

The following shows a sample program for Reset command processing.

```

/*****
/*
/*  Reset command
/*
/*****
/*  [r] u16          ... error code
/*****
u16 fl_ua_reset(void)
{
    u16  rc;
    u32  retry;

    set_uart0_br(BR_9600);    // change to 9600bps

    fl_wait(tCOM);           // wait

    set_ua_dir_tx();         // Change Mono-wire UART transmit mode
    putc_ua(0x00);           // send 0x00 @ 9600bps

    fl_wait(t12);           // wait

    putc_ua(0x00);           // send 0x00 @ 9600bps
    set_ua_dir_rx();         // Change Mono-wire UART receive mode

    for (retry = 0; retry < tRS; retry++){

        fl_wait(t2C);        // wait

        put_cmd_ua(FL_COM_RESET, 1, fl_cmd_prm); // send RESET command

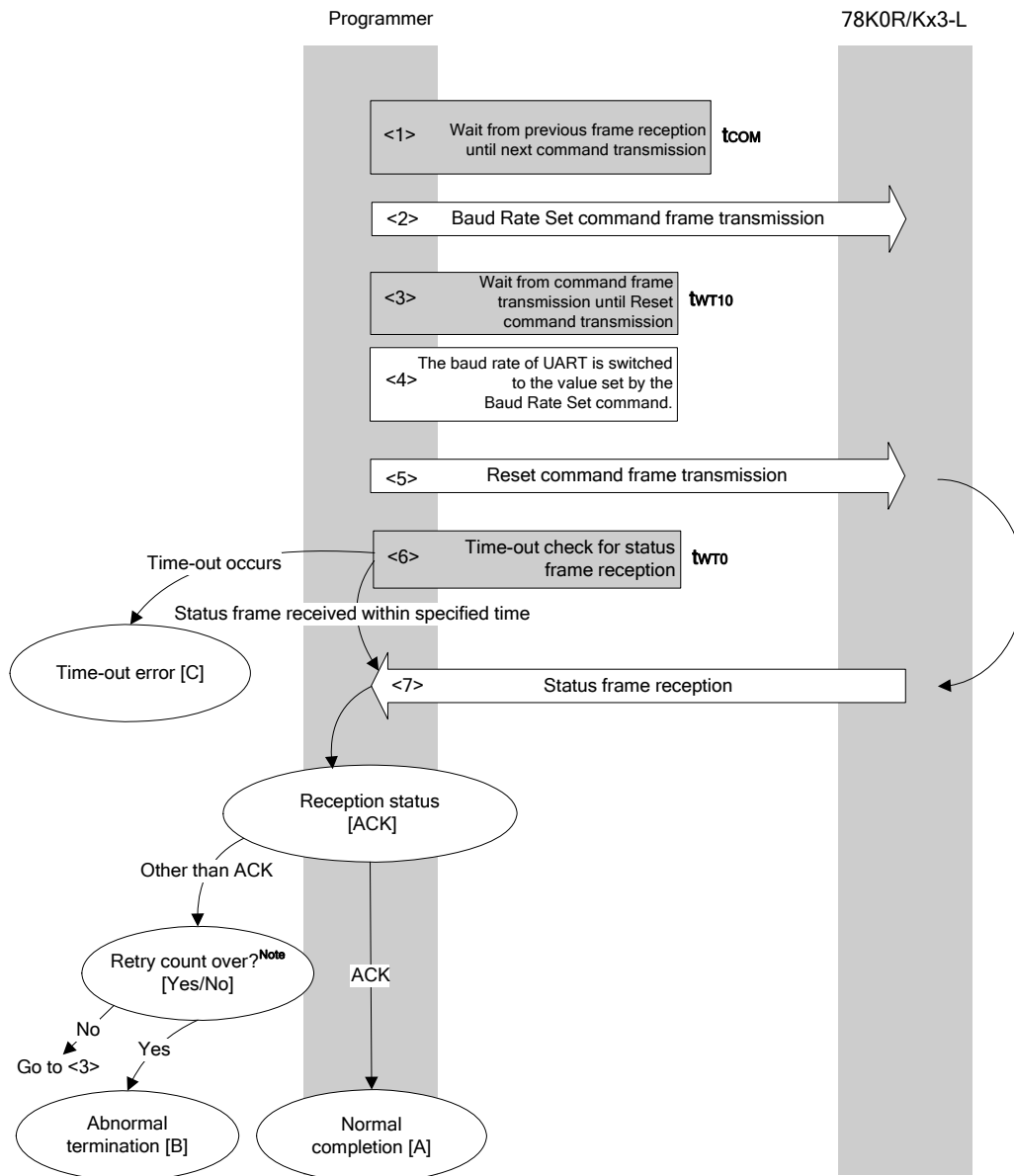
        rc = get_sfrm_ua(fl_ua_sfrm, tWT0_MAX);
        if (rc == FLC_DFTO_ERR)           // t.o. ?
            break;                       // yes // case [C]
        if (rc == FLC_ACK){               // ACK ?
            break;                       // yes // case [A]
        }
        else{
            NOP();
        }
        //continue;                    // case [B] (if exit from loop)
    }
    // switch(rc) {
    //
    //     case  FLC_NO_ERR:    return rc;    break; // case [A]
    //     case  FLC_DFTO_ERR: return rc;    break; // case [C]
    //     default:            return rc;    break; // case [B]
    // }
    return rc;
}

```

4.5 Baud Rate Set Command

4.5.1 Processing sequence chart

Baud Rate Set command processing sequence



**Note** Do not exceed the retry count for the reset command transmission (up to 16 times).

4.5.2 Description of processing sequence

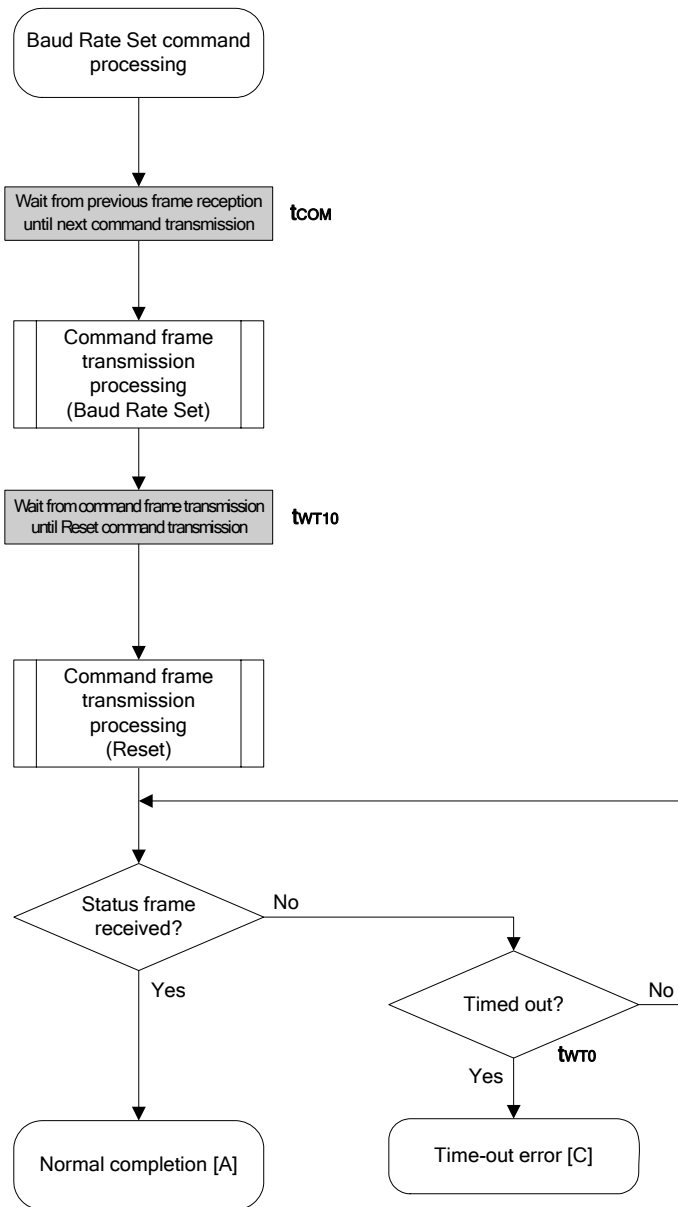
- <1> Waits from the previous frame reception until the next command transmission (wait time  $t_{COM}$ ).
- <2> The Baud Rate Set command is transmitted by command frame transmission processing.
- <3> Waits from command transmission until Reset command transmission (wait time  $t_{WT10}$ ).
- <4> The baud rate of UART communication is switched to the value set by the Baud Rate Set command.
- <5> The Reset command is transmitted by command frame transmission processing.
- <6> A time-out check is performed from command transmission until status frame reception.  
If a time-out occurs, a time-out error [C] is returned (time-out time  $t_{WTO}$ ).
- <7> Since the status code should be ACK, the processing ends normally [A].

4.5.3 Status at processing completion

Status at Processing Completion		Status Code	Description
Normal completion [A]	Normal acknowledgment (ACK)	06H	The command was executed normally and the synchronization of the UART communication speed has been established between the programmer and the 78K0R/Kx3-L.
Abnormal termination [B]	Checksum error	07H	The checksum of the transmitted command frame does not match.
	Negative acknowledgment (NACK)	15H	Command frame data is abnormal (such as invalid data length (LEN) or no ETX).
Time-out error [C] <sup>Note</sup>		–	Data frame reception was timed out. With the 78K0R/Kx3-L, this command also results in errors in the following cases. <ul style="list-style-type: none"> <li>• Command information (D01, D02H, D02L, D03) is invalid</li> <li>• The command frame includes the checksum error</li> <li>• The data length of the command frame (LEN) is invalid</li> <li>• The footer of the command frame (ETX) is missing</li> <li>• The Reset command was not detected after setting the baud rate and receiving command frame data for 16 times.</li> </ul>

**Note** If a time-out error has occurred, execute a hardware reset and re-set to the flash memory programming mode.

4.5.4 Flowchart



## 4.5.5 Sample program

The following shows a sample program for Baud Rate Set command processing.

```

/*****/
/*                                                                 */
/*Set baudrate command                                          */
/*                                                                 */
/*****/
/* [i] u8 brid ... baudrate ID                                  */
/* [i] u8 fpv ... Flash programming voltage                    */
/* [r] u16 ... error code                                       */
/*****/
u16    fl_ua_setbaud(u8 brid, u8 fpv)
{
    u16    rc;
    u8     br;
    u32    retry;

    fl_cmd_prm[0] = 0x00;    // "D01" : adjust by target device (115200bps)
    fl_cmd_prm[1] = 0x00;    // "D02H" : adjust by target device (115200bps)
    fl_cmd_prm[2] = 0x0a;    // "D02L" : (fixed value)
    fl_cmd_prm[3] = 0x01;    // "D03" : noise filter on
    fl_cmd_prm[4] = fpv;     // "D04" : Flash programming voltage

    fl_wait(tCOM);          // wait before sending command
    put_cmd_ua(FL_COM_SET_BAUDRATE, 1+5, fl_cmd_prm);    // send "Baudrate Set"
command
    set_flbaud(brid);       // change baud-rate
    set_uart0_br(brid);     // change baud-rate (h.w.)

    retry = trs;
    while(1){
        fl_wait(tWT10);

        put_cmd_ua(FL_COM_RESET, 1, fl_cmd_prm);    // send RESET command
        rc = get_sfrm_ua(fl_ua_sfrm, tWT0_MAX);    // get status frame
        if (rc){
            if (retry--
                continue;
            else
                return rc;
        }
        break;    // got ACK !!
    }

    // switch(rc) {
    //     case FLC_NO_ERR: return rc; break; // case [A]
    //     case FLC_DFTO_ERR: return rc; break; // case [C]
    //     default: return rc; break; // case [B]
    // }

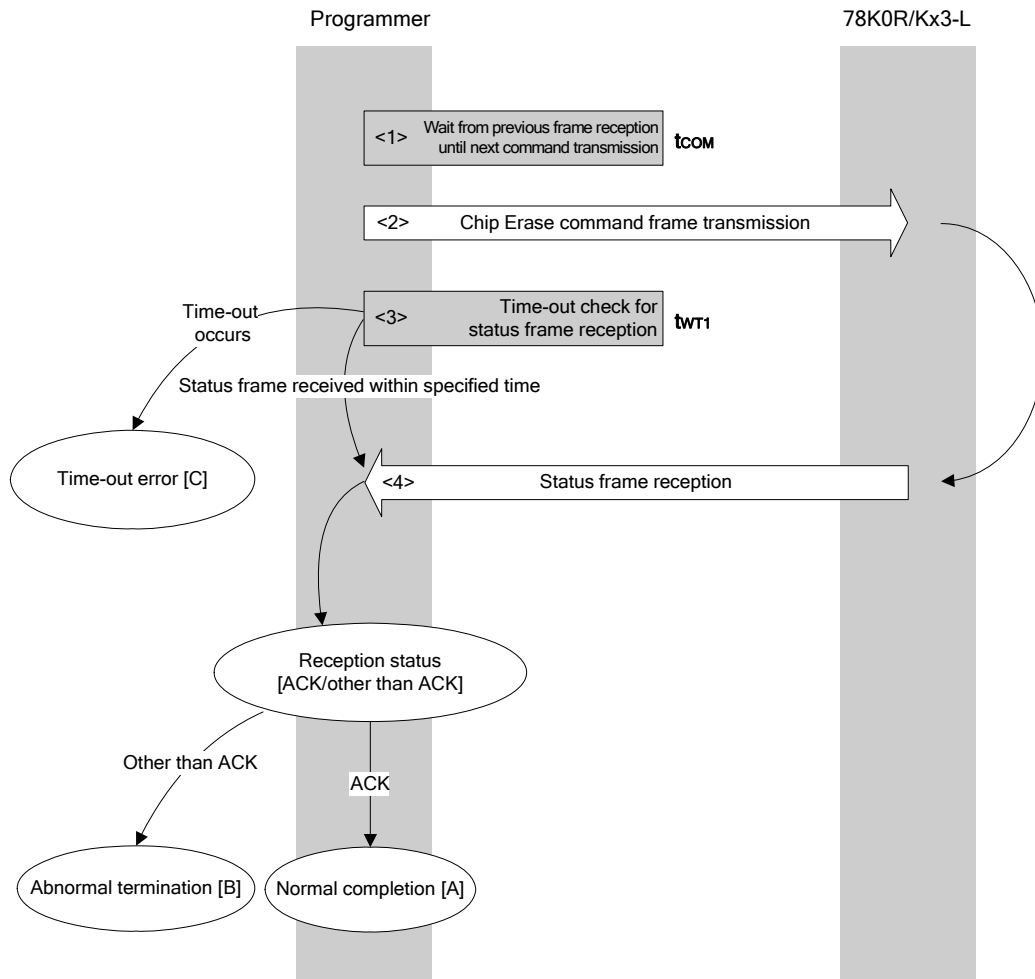
    return rc;
}

```

4.6 Chip Erase Command

4.6.1 Processing sequence chart

Chip Erase command processing sequence





#### 4.6.2 Description of processing sequence

- <1> Waits from the previous frame reception until the next command transmission (wait time  $t_{COM}$ ).
- <2> The Chip Erase command is transmitted by command frame transmission processing.
- <3> A time-out check is performed from command transmission until status frame reception.  
If a time-out occurs, a time-out error [C] is returned (time-out time  $t_{WTF1}$ ).
- <4> The status code is checked.

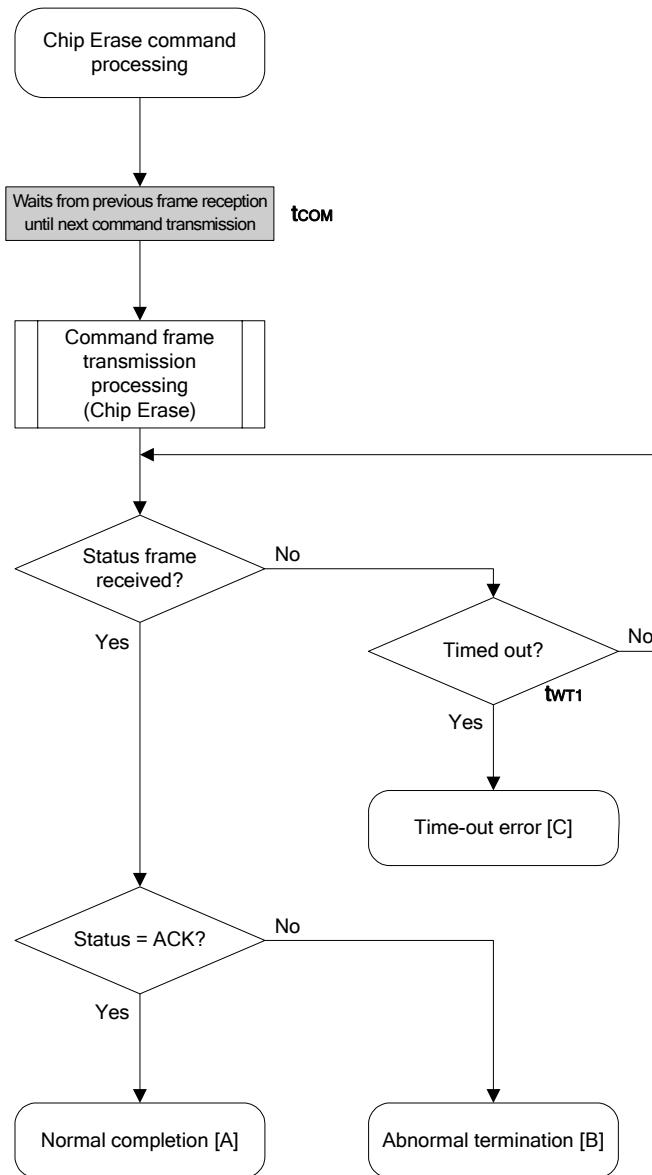
When ST1 = ACK: Normal completion [A]

When ST1 ≠ ACK: Abnormal termination [B]

#### 4.6.3 Status at processing completion

Status at Processing Completion		Status Code	Description
Normal completion [A]	Normal acknowledgment (ACK)	06H	The command was executed normally and chip erase was performed normally.
Abnormal termination [B]	Checksum error	07H	The checksum of the transmitted command frame does not match.
	Protect error	10H	Chip erase or boot block rewrite is prohibited in the security setting.
	Negative acknowledgment (NACK)	15H	Command frame data is abnormal (such as invalid data length (LEN) or no ETX).
	MRG10 error	1AH	An erase error has occurred.
	MRG11 error	1BH	
	Write error	1CH	
Time-out error [C]	–	The status frame was not received within the specified time.	

4.6.4 Flowchart



#### 4.6.5 Sample program

The following shows a sample program for Chip Erase command processing.

```

/*****
/*
/*  Erase all(chip) command
/*
/*****
/*  [r] u16          ... error code
/*****
u16      fl_ua_erase_all(void)
{
    u16    rc;

    fl_wait(tCOM);          // wait before sending command

    put_cmd_ua(FL_COM_ERASE_CHIP, 1, fl_cmd_prm); // send ERASE CHIP command

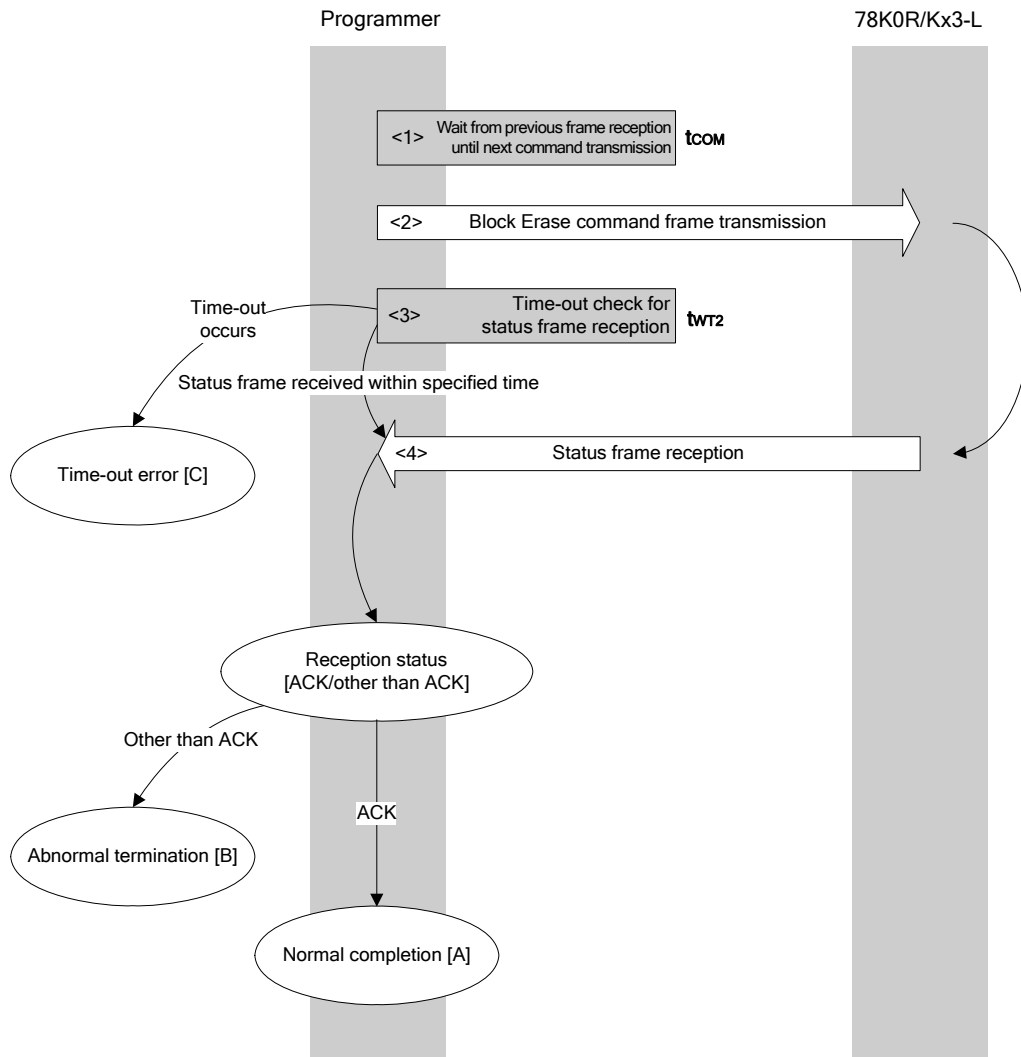
    rc = get_sfrm_ua(fl_ua_sfrm, tWT1_MAX); // get status frame
    // switch(rc) {
    //
    //     case  FLC_NO_ERR:    return rc;          break; // case [A]
    //     case  FLC_DFEO_ERR: return rc;          break; // case [C]
    //     default:            return rc;          break; // case [B]
    // }
    return rc;
}

```

4.7 Block Erase Command

4.7.1 Processing sequence chart

Block Erase command processing sequence



#### 4.7.2 Description of processing sequence

- <1> Waits from the previous frame reception until the next command transmission (wait time  $t_{COM}$ ).
- <2> The Block Erase command is transmitted by command frame transmission processing.
- <3> A time-out check is performed from command transmission until status frame reception.  
If a time-out occurs, a time-out error [C] is returned (time-out time  $t_{WT2}$ ).
- <4> The status code is checked.

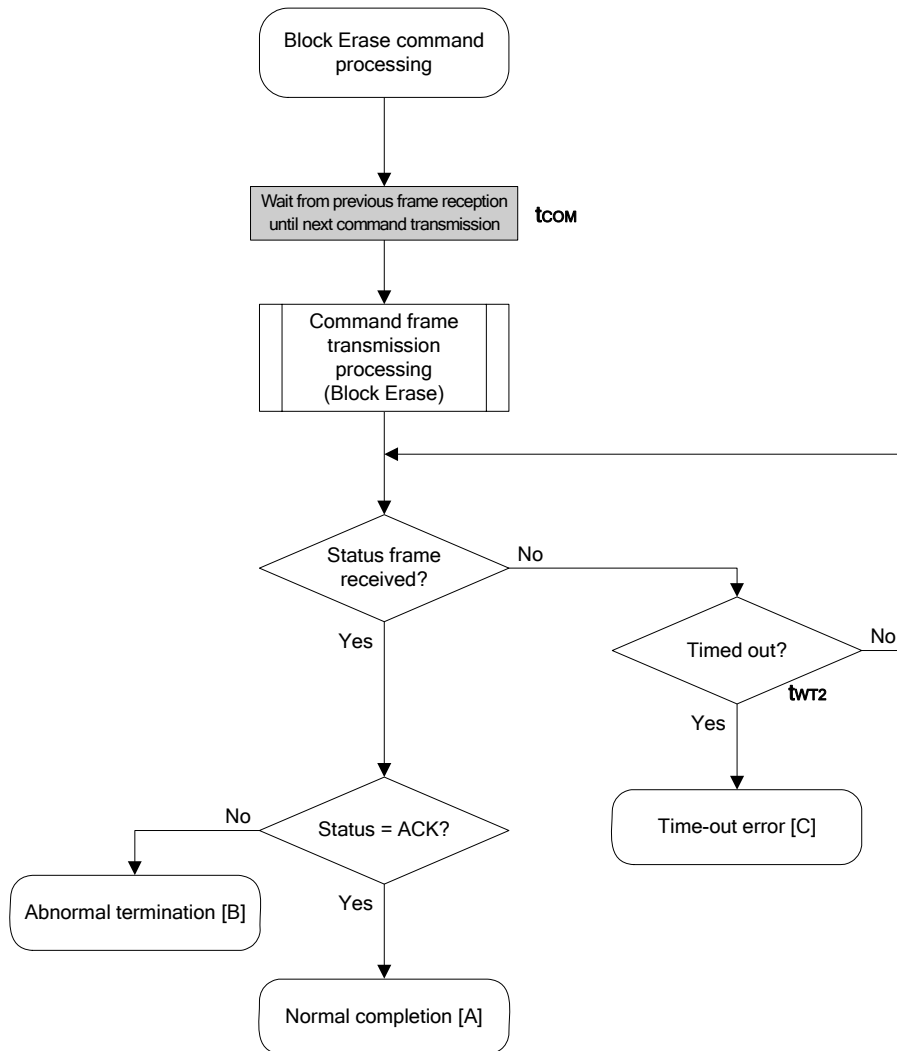
When ST1 = ACK: Normal completion [A]

When ST1 ≠ ACK: Abnormal termination [B]

#### 4.7.3 Status at processing completion

Status at Processing Completion		Status Code	Description
Normal completion [A]	Normal acknowledgment (ACK)	06H	The command was executed normally and block erase was performed normally.
Abnormal termination [B]	Parameter error	05H	The specified end address is out of the flash memory range, or the specified start/end address is not the first/end address of the block.
	Checksum error	07H	The checksum of the transmitted command frame does not match.
	Protect error	10H	Write, block erase, or chip erase is prohibited in the security setting. A boot block is included in the specified range and boot block rewrite is prohibited.
	Negative acknowledgment (NACK)	15H	Command frame data is abnormal (such as invalid data length (LEN) or no ETX).
	MRG10 error	1AH	An erase error has occurred.
Time-out error [C]		–	The status frame was not received within the specified time.

## 4.7.4 Flowchart



### 4.7.5 Sample program

The following shows a sample program for Block Erase command processing.

```

/*****
/*
/*  Erase block command
/*
/*****
/*  [i] u8 block      ... block number
/*  [r] u16           ... error code
/*****
u16      fl_ua_erase_blk(u16 sblk, u16 eblk)
{

    u16    rc;
    u32    wt2_max;
    u32    top, bottom;

    top = get_top_addr(sblk);           // get start address of start block
    bottom = get_bottom_addr(eblk);     // get end address of end block

    set_range_prm(fl_cmd_prm, top, bottom); // set SAH/SAM/SAL, EAH/EAM/EAL

    wt2_max = make_wt2_max(sblk, eblk);

    fl_wait(tCOM);                      // wait before sending command

    put_cmd_ua(FL_COM_ERASE_BLOCK, 1+6, fl_cmd_prm); // send ERASE CHIP command

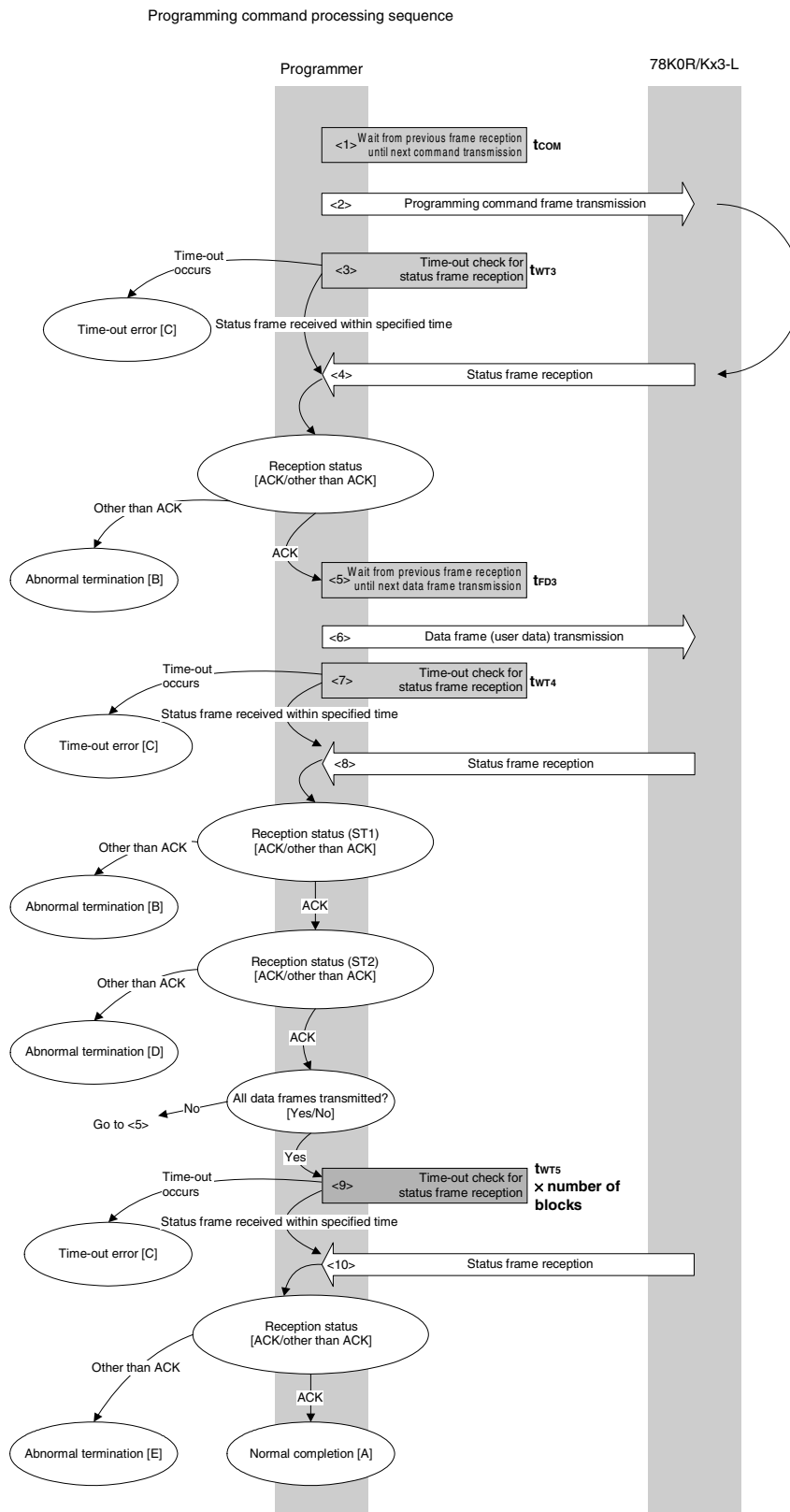
    rc = get_sfrm_ua(fl_ua_sfrm, wt2_max); // get status frame
    // switch(rc) {
    //
    //     case  FLC_NO_ERR:    return rc;        break; // case [A]
    //     case  FLC_DFTO_ERR:  return rc;        break; // case [C]
    //     default:             return rc;        break; // case [B]
    // }

    return rc;
}

```

### 4.8 Programming Command

#### 4.8.1 Processing sequence chart





#### 4.8.2 Description of processing sequence

- <1> Waits from the previous frame reception until the next command transmission (wait time  $t_{COM}$ ).
- <2> The Programming command is transmitted by command frame transmission processing.
- <3> A time-out check is performed from command transmission until status frame reception.  
If a time-out occurs, a time-out error [C] is returned (time-out time  $t_{WT3}$ ).
- <4> The status code is checked.

When ST1 = ACK: Proceeds to <5>.

When ST1  $\neq$  ACK: Abnormal termination [B]

- <5> Waits from the previous frame reception until the next data frame transmission (wait time  $t_{FD3}$ ).
- <6> User data is transmitted by data frame transmission processing.
- <7> A time-out check is performed from user data transmission until data frame reception.  
If a time-out occurs, a time-out error [C] is returned (time-out time  $t_{WT4}$ ).
- <8> The status code (ST1/ST2) is checked (also refer to the processing sequence chart and flowchart).

When ST1  $\neq$  ACK: Abnormal termination [B]

When ST1 = ACK: The following processing is performed according to the ST2 value.

- When ST2 = ACK: Proceeds to <9> when transmission of all data frames is completed.  
If there still remain data frames to be transmitted, the processing re-executes the sequence from <5>.
- When ST2  $\neq$  ACK: Abnormal termination [D]

- <9> A time-out check is performed until status frame reception.  
If a time-out occurs, a time-out error [C] is returned (time-out time  $t_{WT5} \times$  number of blocks).
- <10> The status code is checked.

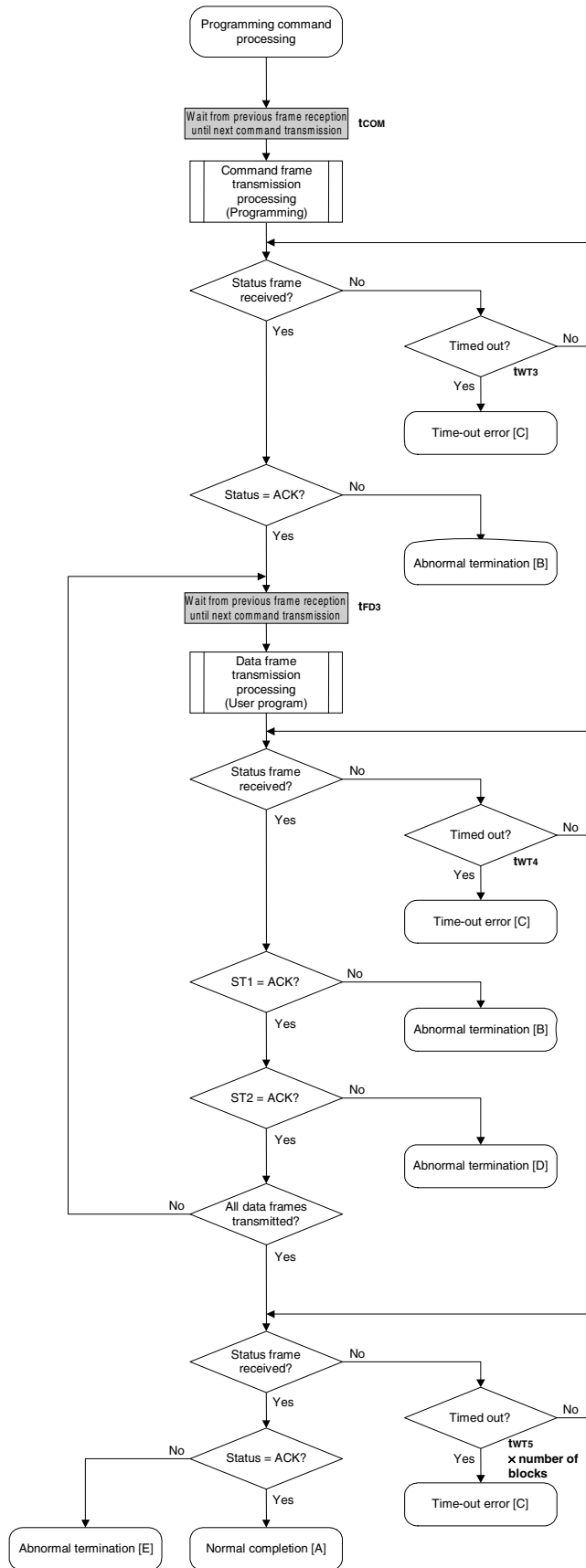
When ST1 = ACK: Normal completion [A]

When ST1  $\neq$  ACK: Abnormal termination [E]

## 4.8.3 Status at processing completion

Status at Processing Completion		Status Code	Description
Normal completion [A]	Normal acknowledgment (ACK)	06H	The command was executed normally and the user data was written normally.
Abnormal termination [B]	Parameter error	05H	The start/end address is out of the flash memory range, the specified start/end address is not the first/end address of the block, or the write start address is larger than the end address.
	Checksum error	07H	The checksum of the transmitted command frame or data frame does not match.
	Protect error	10H	Write is prohibited in the security setting. A boot block is included in the specified range and boot block rewrite is prohibited.
	Negative acknowledgment (NACK)	15H	Command frame data or data frame data is abnormal (such as invalid data length (LEN) or no ETX).
Time-out error [C]		–	The status frame was not received within the specified time.
Abnormal termination [D], [E]	MRG10 error	1AH	A write error has occurred.
	MRG11 error	1BH	
	Write error	1CH	

4.8.4 Flowchart



## 4.8.5 Sample program

The following shows a sample program for Programming command processing.

```

/*****
/*
/* Write command
/*
/*****
/* [i] u32 top      ... start address
/* [i] u32 bottom  ... end address
/* [r] u16         ... error code
/*****

#define fl_st2_ua      (fl_ua_sfrm[OFS_STA_PLD+1])

u16      fl_ua_write(u32 top, u32 bottom)
{
    u16    rc;
    u32    send_head, send_size;
    bool   is_end;
    u16    block_num;

    block_num = get_block_num(top, bottom);          // get block num

    /*****
    /*      set params
    /*
    /*****
    set_range_prm(fl_cmd_prm, top, bottom);          // set SAH/SAM/SAL, EAH/EAM/EAL

    /*****
    /*      send command & check status
    /*
    /*****
    fl_wait(tCOM);                                  // wait before sending command

    put_cmd_ua(FL_COM_WRITE, 7, fl_cmd_prm);         // send "Programming" command

    rc = get_sfrm_ua(fl_ua_sfrm, tWT3_MAX);          // get status frame
    switch(rc) {
        case      FLC_NO_ERR:                        break; // continue
    //   case      FLC_DFTO_ERR: return rc;          break; // case [C]
        default:                                     return rc;   break; // case [B]
    }

    /*****
    /*      send user data
    /*
    /*****
    send_head = top;

    while(1){
        // make send data frame
        if ((bottom - send_head) > 256){            // rest size > 256 ?
            is_end = false;                          // yes, not is_end frame

```

```

        send_size = 256;          // transmit size = 256 byte
    }
    else{
        is_end = true;
        send_size = bottom - send_head + 1;    // transmit size = (bottom
- send_head)+1 byte

    }
    memcpy(fl_txdata_frm, rom_buf+send_head, send_size); // set data frame
payload
    send_head += send_size;

    fl_wait(tFD3);                // wait before sending data
frame

    put_dfrm_ua(send_size, fl_txdata_frm, is_end); // send user data

    rc = get_sfrm_ua(fl_ua_sfrm, tWT4_MAX);      // get status frame
    switch(rc) {
        case FLC_NO_ERR:                break; // continue
        case FLC_DFTO_ERR: return rc;    break; // case [C]
        default:                return rc; break; // case [B]
    }
    if (fl_st2_ua != FLST_ACK){          // ST2 = ACK ?
        rc = decode_status(fl_st2_ua);    // No
        return rc;                       // case [D]
    }
    if (is_end)
        break;

}
/*****
/*      Check internally verify          */
/*****/
// get status frame again
if (top == 0){/* include Block0 */
    rc = get_sfrm_ua(fl_ua_sfrm, tWT5_B0_MAX + tWT5_MAX*(block_num-1));
}
else{
    rc = get_sfrm_ua(fl_ua_sfrm, tWT5_MAX*block_num);
}
switch(rc) {
// case FLC_NO_ERR: return rc; break; // case [A]
case FLC_DFTO_ERR: return rc; break; // case [C]
default:                return rc; break; // case [E]
}

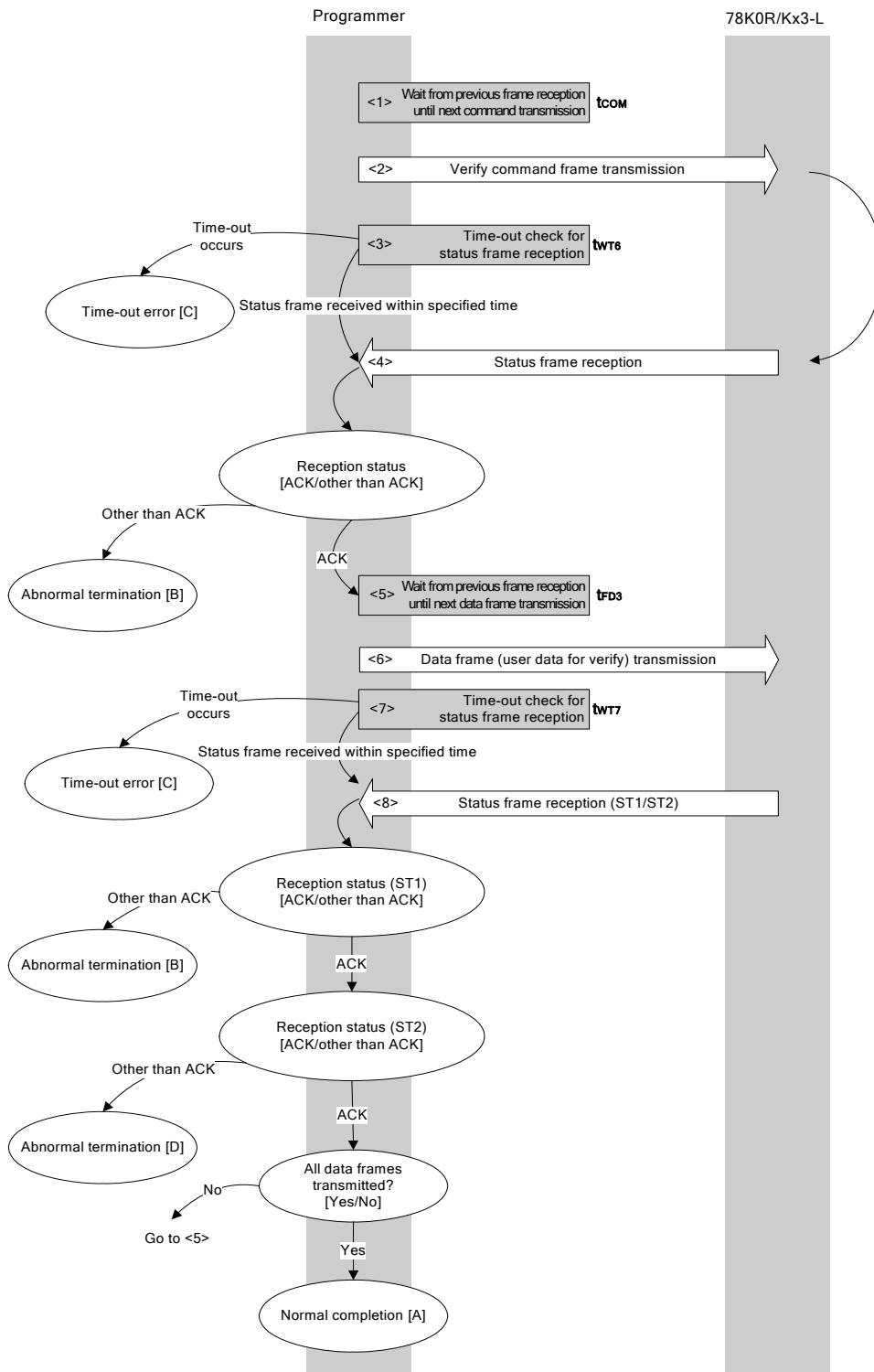
return rc;
}

```

## 4.9 Verify Command

### 4.9.1 Processing sequence chart

Verify command processing sequence



### 4.9.2 Description of processing sequence

- <1> Waits from the previous frame reception until the next command transmission (wait time  $t_{COM}$ ).
- <2> The Verify command is transmitted by command frame transmission processing.
- <3> A time-out check is performed from command transmission until status frame reception.  
If a time-out occurs, a time-out error [C] is returned (time-out time  $t_{WT6}$ ).
- <4> The status code is checked.

When ST1 = ACK: Proceeds to <5>.

When ST1 ≠ ACK: Abnormal termination [B]

- <5> Waits from the previous frame reception until the next data frame transmission (wait time  $t_{FD3}$ ).
- <6> User data for verifying is transmitted by data frame transmission processing.
- <7> A time-out check is performed from user data transmission until status frame reception.  
If a time-out occurs, a time-out error [C] is returned (time-out time  $t_{WT7}$ ).
- <8> The status code (ST1/ST2) is checked (also refer to the processing sequence chart and flowchart).

When ST1 ≠ ACK: Abnormal termination [B]

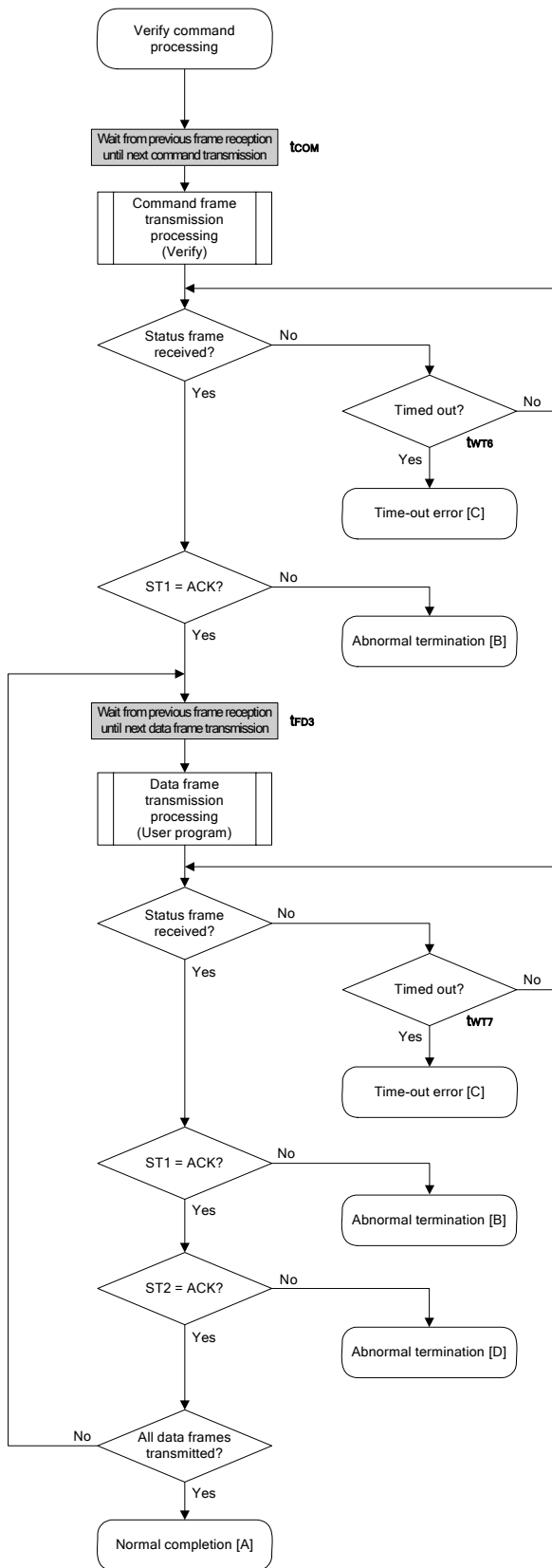
When ST1 = ACK: The following processing is performed according to the ST2 value.

- When ST2 = ACK: If transmission of all data frames is completed, the processing ends normally [A].  
If there still remain data frames to be transmitted, the processing re-executes the sequence from <5>.
- When ST2 ≠ ACK: Abnormal termination [D]

### 4.9.3 Status at processing completion

Status at Processing Completion		Status Code	Description
Normal completion [A]	Normal acknowledgment (ACK)	06H	The command was executed normally and the verify was completed normally.
Abnormal termination [B]	Parameter error	05H	The start/end address is out of the flash memory range, the start/end address is not the start/end address of the block, or the write start address is larger than the end address.
	Checksum error	07H	The checksum of the transmitted command frame or data frame does not match.
	Negative acknowledgment (NACK)	15H	Command frame data is abnormal (such as invalid data length (LEN) or no ETX).
Time-out error [C]		–	The status frame was not received within the specified time.
Abnormal termination [D]	Verify error	0FH (ST2)	A verify error has occurred.

4.9.4 Flowchart





## 4.9.5 Sample program

The following shows a sample program for Verify command processing.

```

/*****
/*
/* Verify command
/*
/*****
/* [i] u32 top          ... start address
/* [i] u32 bottom      ... end address
/* [r] u16             ... error code
/*****
u16      fl_ua_verify(u32 top, u32 bottom, u8 *buf)
{
    u16    rc;
    u32    send_head, send_size;
    bool   is_end;

    /*****
    /* set params
    /*****
    set_range_prm(fl_cmd_prm, top, bottom); // set SAH/SAM/SAL, EAH/EAM/EAL

    /*****
    /* send command & check status
    /*****

    fl_wait(tCOM); // wait before sending command

    put_cmd_ua(FL_COM_VERIFY, 7, fl_cmd_prm); // send VERIFY command

    rc = get_sfrm_ua(fl_ua_sfrm, tWT6_MAX); // get status frame
    switch(rc) {
        case FLC_NO_ERR: break; // continue
        // case FLC_DFTO_ERR: return rc; break; // case [C]
        default: return rc; break; // case [B]
    }

    /*****
    /* send user data
    /*****
    send_head = top;

    while(1){

        // make send data frame
        if ((bottom - send_head) > 256){ // rest size > 256 ?
            is_end = false; // yes, not is_end frame
            send_size = 256; // transmit size = 256 byte
        }
        else{
            is_end = true;
            send_size = bottom - send_head + 1; // transmit size = (bottom -
send_head)+1 byte

```

```
    }
    memcpy(fl_txdata_frm, buf+send_head, send_size); // set data frame payload
    send_head += send_size;

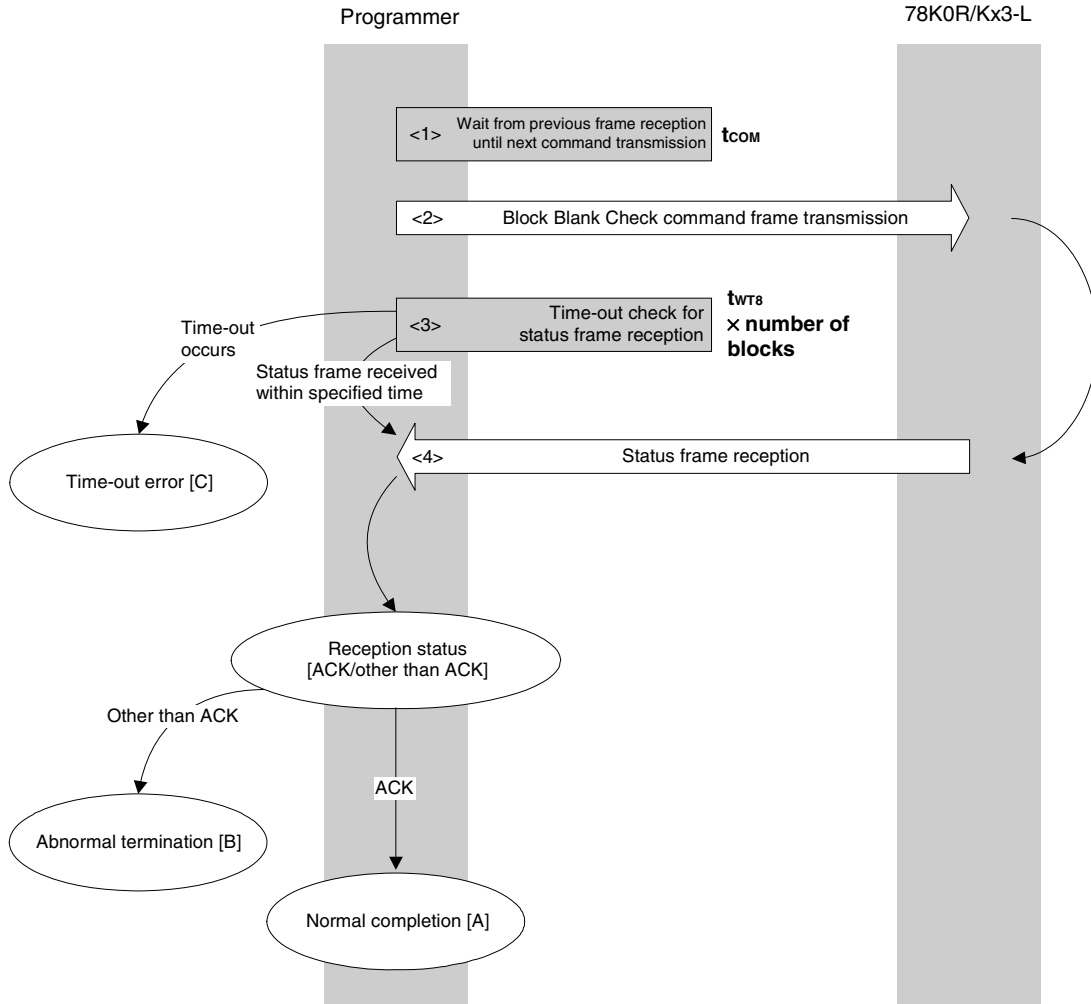
    fl_wait(tFD3);
    put_dfrm_ua(send_size, fl_txdata_frm, is_end); // send user data

    rc = get_sfrm_ua(fl_ua_sfrm, tWT7_MAX); // get status frame
    switch(rc) {
        case FLC_NO_ERR: break; // continue
        // case FLC_DFTO_ERR: return rc; break; // case [C]
        default: return rc; break; // case [B]
    }
    if (fl_st2_ua != FLST_ACK){ // ST2 = ACK ?
        rc = decode_status(fl_st2_ua); // No
        return rc; // case [D]
    }
    if (is_end) // send all user data ?
        break; // yes
    //continue;
}
return FLC_NO_ERR; // case [A]
}
```

## 4.10 Block Blank Check Command

### 4.10.1 Processing sequence chart

Block Blank Check command processing sequence



### 4.10.2 Description of processing sequence

- <1> Waits from the previous frame reception until the next command transmission (wait time  $t_{COM}$ ).
- <2> The Block Blank Check command is transmitted by command frame transmission processing.
- <3> A time-out check is performed from command transmission until status frame reception.  
If a time-out occurs, a time-out error [C] is returned (time-out time  $t_{WT8} \times$  number of blocks).
- <4> The status code is checked.

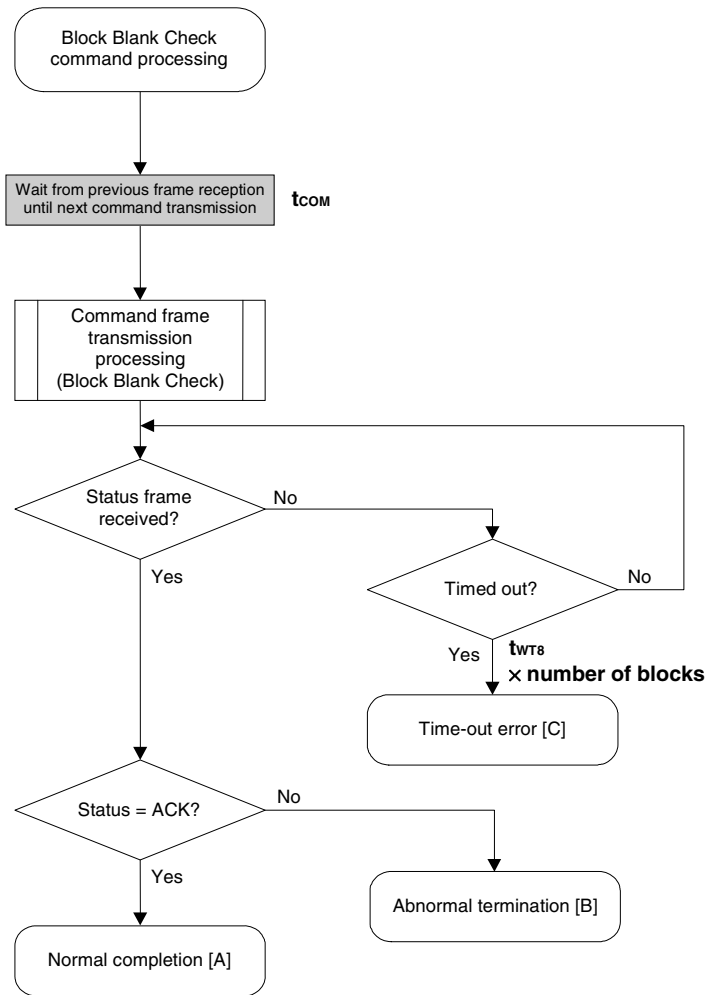
When ST1 = ACK: Normal completion [A]

When ST1  $\neq$  ACK: Abnormal termination [B]

### 4.10.3 Status at processing completion

Status at Processing Completion		Status Code	Description
Normal completion [A]	Normal acknowledgment (ACK)	06H	The command was executed normally and block blank check was executed normally.
Abnormal termination [B]	Parameter error	05H	The end address is out of the flash memory range, the start/end address is not the first/end address of the block, or the value of parameter D01 is other than 00H or 01H.
	Checksum error	07H	The checksum of the transmitted command frame does not match.
	Negative acknowledgment (NACK)	15H	Command frame data is abnormal (such as invalid data length (LEN) or no ETX).
	MRG11 error	1BH	The flash memory of the specified block is not blank.
Time-out error [C]		–	The status frame was not received within the specified time.

4.10.4 Flowchart



## 4.10.5 Sample program

The following shows a sample program for Block Blank Check command processing.

```

/*****
/*
/* Block blank check command
/*
/*****
/* [i] u32 top      ... top address of blank check
/* [i] u32 bottom  ... bottom address of blank check
/* [i] u8 whole    ... <1>check w/NON user flash
/*                <0>chek only user flash
/* [r] u16         ... error code
/*****
u16      fl_ua_blk_blank_chk(u32 top, u32 bottom, u8 whole)
{
    u16    rc;
    u16    block_num;

    set_range_prm(fl_cmd_prm, top, bottom); // set SAH/SAM/SAL, EAH/EAM/EAL
    block_num = get_block_num(top, bottom); // get block num
    fl_cmd_prm[6] = whole;                  // check only user area or not

    fl_wait(tCOM);                          // wait before sending command

    put_cmd_ua(FL_COM_BLOCK_BLANK_CHK, 7+1, fl_cmd_prm);

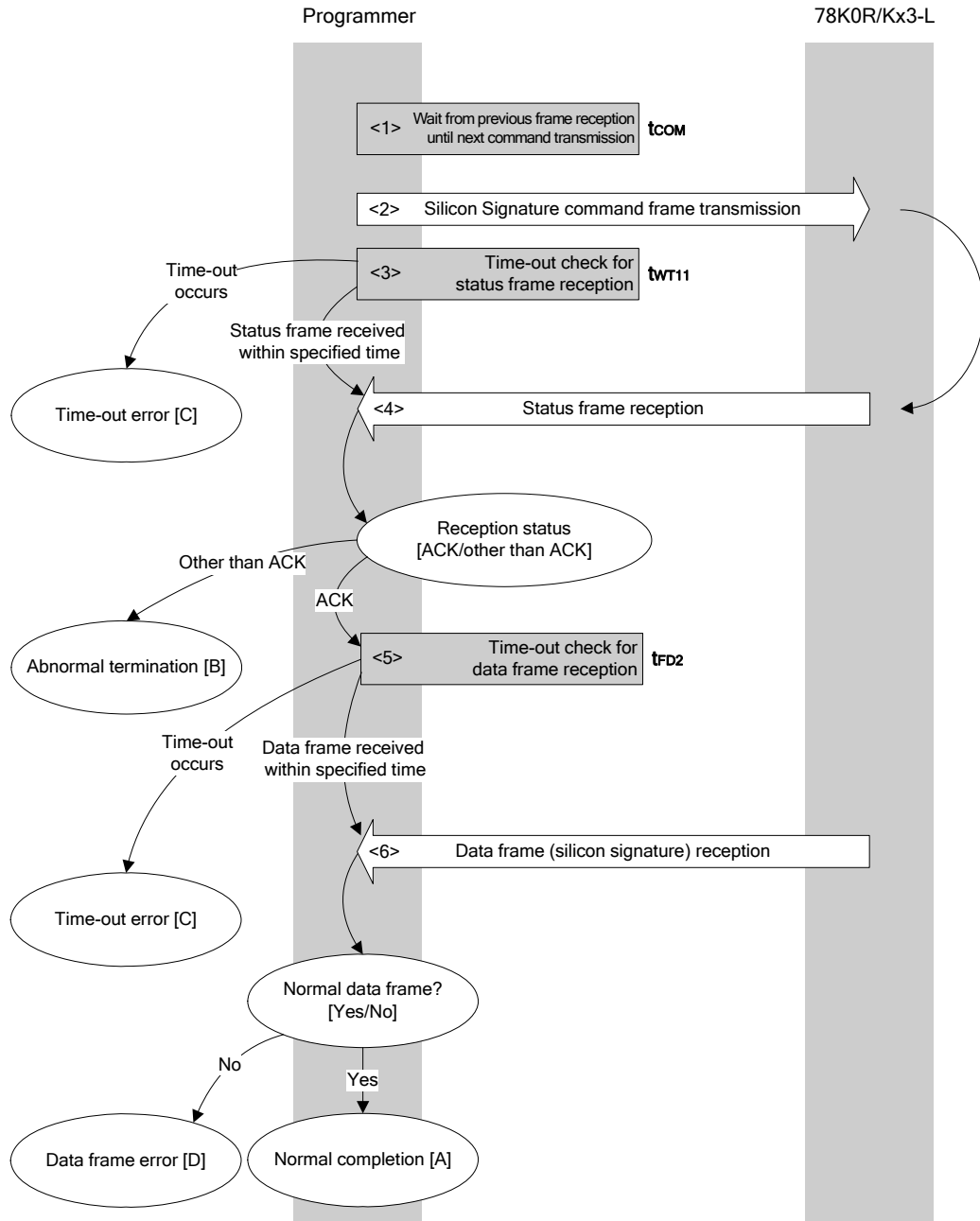
    rc = get_sfrm_ua(fl_ua_sfrm, tWT8_MAX * block_num); // get status frame
    // switch(rc) {
    //
    //     case    FLC_NO_ERR:    return rc;        break; // case [A]
    //     case    FLC_DFTO_ERR: return rc;        break; // case [C]
    //     default:                return rc;      break; // case [B]
    // }
    return rc;
}

```

### 4.11 Silicon Signature Command

#### 4.11.1 Processing sequence chart

Silicon Signature command processing sequence



### 4.11.2 Description of processing sequence

- <1> Waits from the previous frame reception until the next command transmission (wait time  $t_{COM}$ ).
- <2> The Silicon Signature command is transmitted by command frame transmission processing.
- <3> A time-out check is performed from command transmission until status frame reception.  
If a time-out occurs, a time-out error [C] is returned (time-out time  $t_{WT11}$ ).
- <4> The status code is checked.

When ST1 = ACK: Proceeds to <5>.

When ST1  $\neq$  ACK: Abnormal termination [B]

- <5> A time-out check is performed until data frame (silicon signature data) reception.  
If a time-out occurs, a time-out error [C] is returned (time-out time  $t_{FD2}$ ).
- <6> The received data frame (silicon signature data) is checked.

If data frame is normal: Normal completion [A]

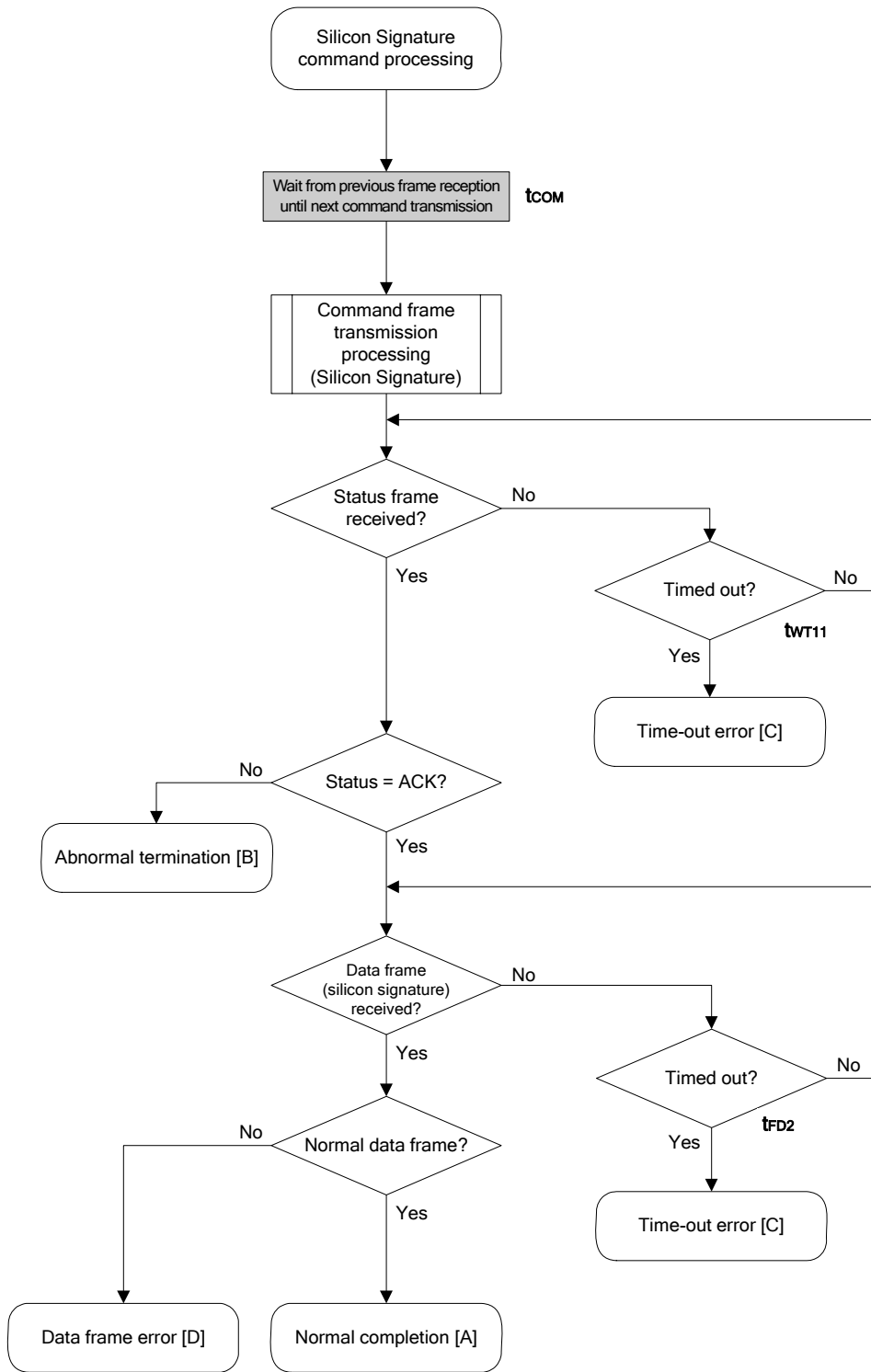
If data frame is abnormal: Data frame error [D]

### 4.11.3 Status at processing completion

Status at Processing Completion		Status Code	Description
Normal completion [A]	Normal acknowledgment (ACK)	06H	The command was executed normally and silicon signature data was acquired normally.
Abnormal termination [B]	Checksum error	07H	The checksum of the transmitted command frame does not match.
	Negative acknowledgment (NACK)	15H	Command frame data is abnormal (such as invalid data length (LEN) or no ETX).
Time-out error [C]		–	The status frame or data frame was not received within the specified time.
Data frame error [D]		–	The checksum of the data frame received as silicon signature data does not match.



4.11.4 Flowchart



## 4.11.5 Sample program

The following shows a sample program for Silicon Signature command processing.

```

/*****
/*
/* Get silicon signature command
/*
/*****
/* [i] u8 *sig ... pointer to signature save area
/* [r] u16 ... error code
/*****
u16 fl_ua_getsig(u8 *sig)
{
    u16 rc;

    fl_wait(tCOM); // wait before sending command

    put_cmd_ua(FL_COM_GET_SIGNATURE, 1, fl_cmd_prm); // send GET SIGNATURE command

    rc = get_sfrm_ua(fl_ua_sfrm, tWT11_MAX); // get status frame
    switch(rc) {
        case FLC_NO_ERR: break; // continue
        // case FLC_DF2O_ERR: return rc; break; // case [C]
        default: return rc; break; // case [B]
    }

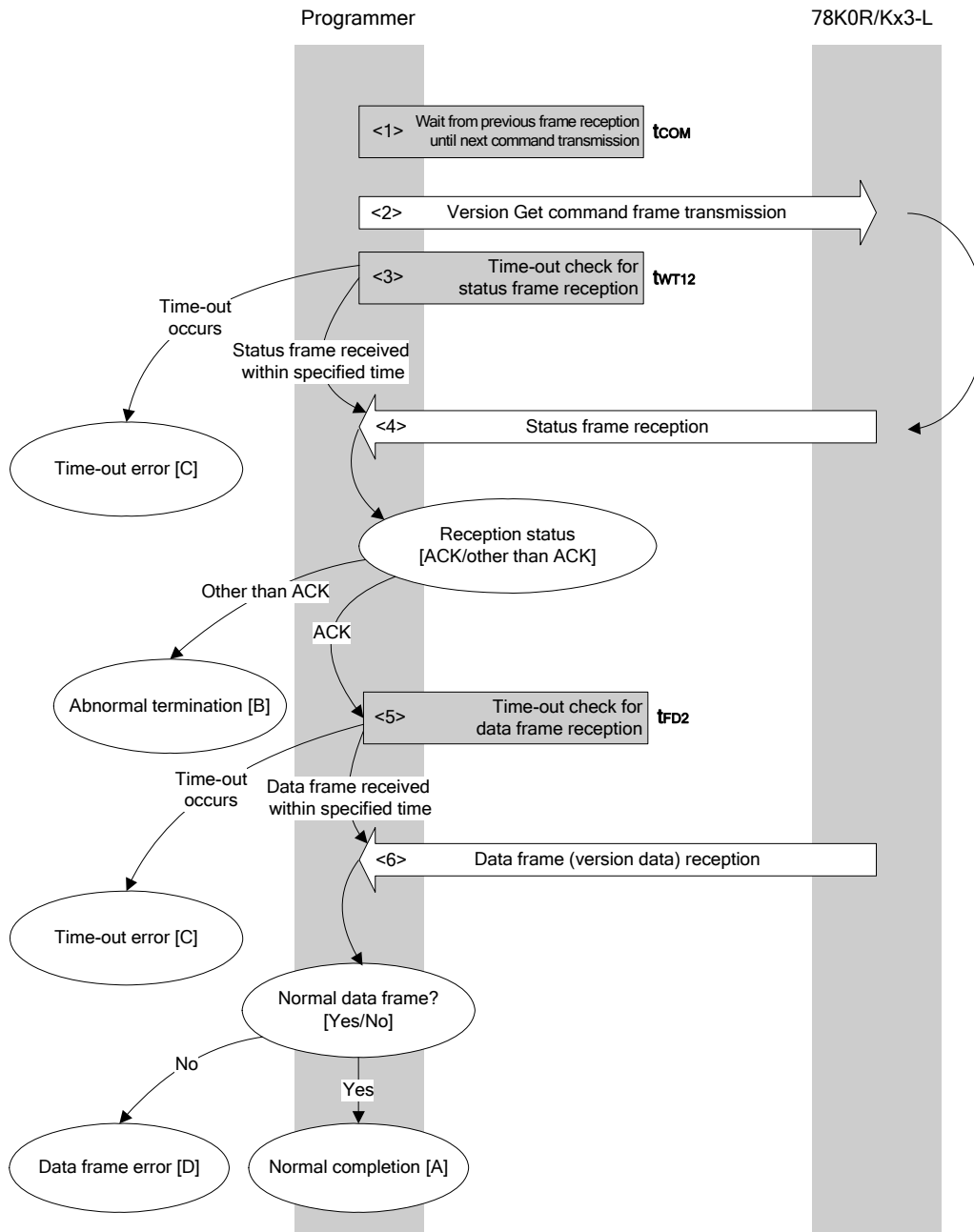
    rc = get_dfrm_ua(fl_rxddata_frm, tFD2_MAX); // get status frame
    if (rc){ // if error
        return rc; // case [D]
    }
    memcpy(sig, fl_rxddata_frm+OFS_STA_PLD, fl_rxddata_frm[OFS_LEN]); // copy Signature
data
    return rc; // case [A]
}

```

### 4.12 Version Get Command

#### 4.12.1 Processing sequence chart

Version Get command processing sequence



### 4.12.2 Description of processing sequence

- <1> Waits from the previous frame reception until the next command transmission (wait time  $t_{COM}$ ).
- <2> The Version Get command is transmitted by command frame transmission processing.
- <3> A time-out check is performed from command transmission until status frame reception.  
If a time-out occurs, a time-out error [C] is returned (time-out time  $t_{WT12}$ ).
- <4> The status code is checked.

When ST1 = ACK: Proceeds to <5>.

When ST1  $\neq$  ACK: Abnormal termination [B]

- <5> A time-out check is performed until data frame (version data) reception.  
If a time-out occurs, a time-out error [C] is returned (time-out time  $t_{FD2}$ ).
- <6> The received data frame (version data) is checked.

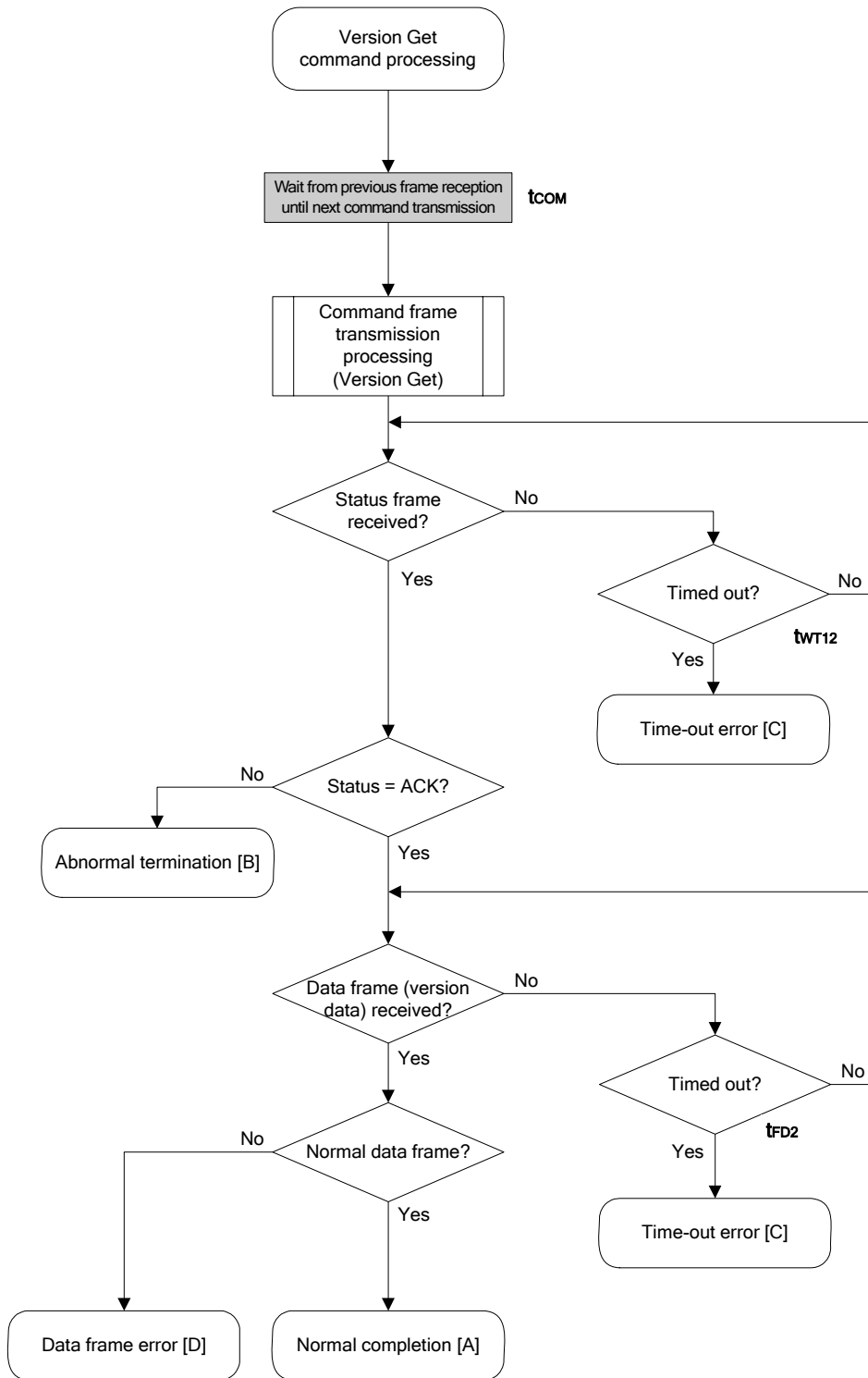
If data frame is normal: Normal completion [A]

If data frame is abnormal: Data frame error [D]

### 4.12.3 Status at processing completion

Status at Processing Completion		Status Code	Description
Normal completion [A]	Normal acknowledgment (ACK)	06H	The command was executed normally and version data was acquired normally.
Abnormal termination [B]	Checksum error	07H	The checksum of the transmitted command frame does not match.
	Negative acknowledgment (NACK)	15H	Command frame data is abnormal (such as invalid data length (LEN) or no ETX).
Time-out error [C]		–	The status frame or data frame was not received within the specified time.
Data frame error [D]		–	The checksum of the data frame received as version data does not match.

4.12.4 Flowchart



## 4.12.5 Sample program

The following shows a sample program for Version Get command processing.

```

/*****
/*
/* Get device/firmware version command
/*
/*****
/* [i] u8 *buf    ... pointer to version data save area
/* [r] u16        ... error code
/*****
u16      fl_ua_getver(u8 *buf)
{
    u16    rc;

    fl_wait(tCOM);                // wait before sending command

    put_cmd_ua(FL_COM_GET_VERSION, 1, fl_cmd_prm); // send GET VERSION command

    rc = get_sfrm_ua(fl_ua_sfrm, tWT12_MAX);        // get status frame
    switch(rc) {
        case      FLC_NO_ERR:                break; // continue
//     case      FLC_DFTO_ERR: return rc;    break; // case [C]
        default:                return rc;    break; // case [B]
    }

    rc = get_dfrm_ua(fl_rxddata_frm, tFD2_MAX);    // get data frame
    if (rc){
        return rc;                            // case [D]
    }

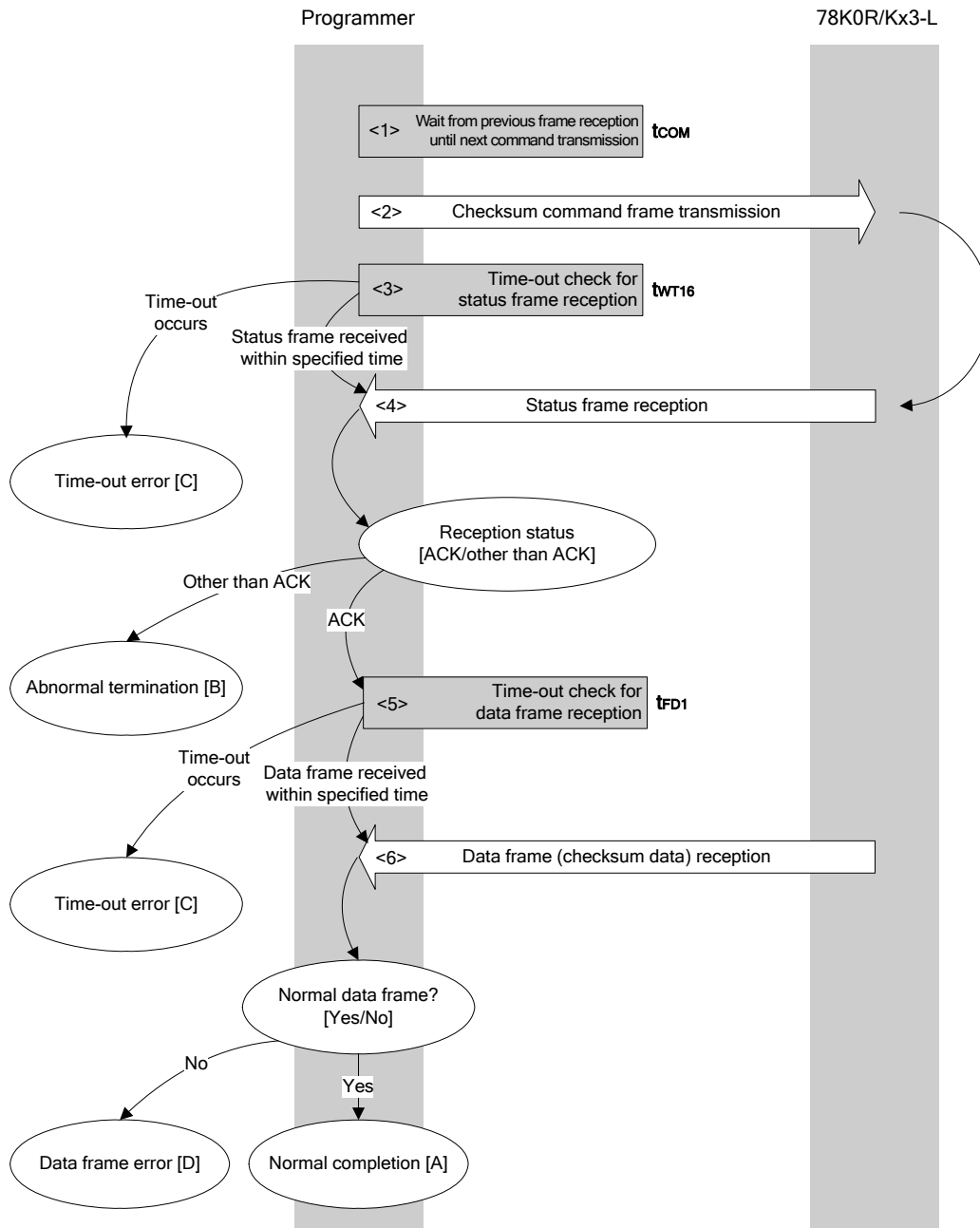
    memcpy(buf, fl_rxddata_frm+OFS_STA_PLD, DFV_LEN); // copy version data
    return rc;                                // case [A]
}

```

4.13 Checksum Command

4.13.1 Processing sequence chart

Checksum command processing sequence



### 4.13.2 Description of processing sequence

- <1> Waits from the previous frame reception until the next command transmission (wait time  $t_{COM}$ ).
- <2> The Checksum command is transmitted by command frame transmission processing.
- <3> A time-out check is performed from command transmission until status frame reception.  
If a time-out occurs, a time-out error [C] is returned (time-out time  $t_{WT16}$ ).
- <4> The status code is checked.

When ST1 = ACK: Proceeds to <5>.

When ST1  $\neq$  ACK: Abnormal termination [B]

- <5> A time-out check is performed until data frame (checksum data) reception.  
If a time-out occurs, a time-out error [C] is returned (time-out time  $t_{FD1}$ ).
- <6> The received data frame (checksum data) is checked.

If data frame is normal: Normal completion [A]

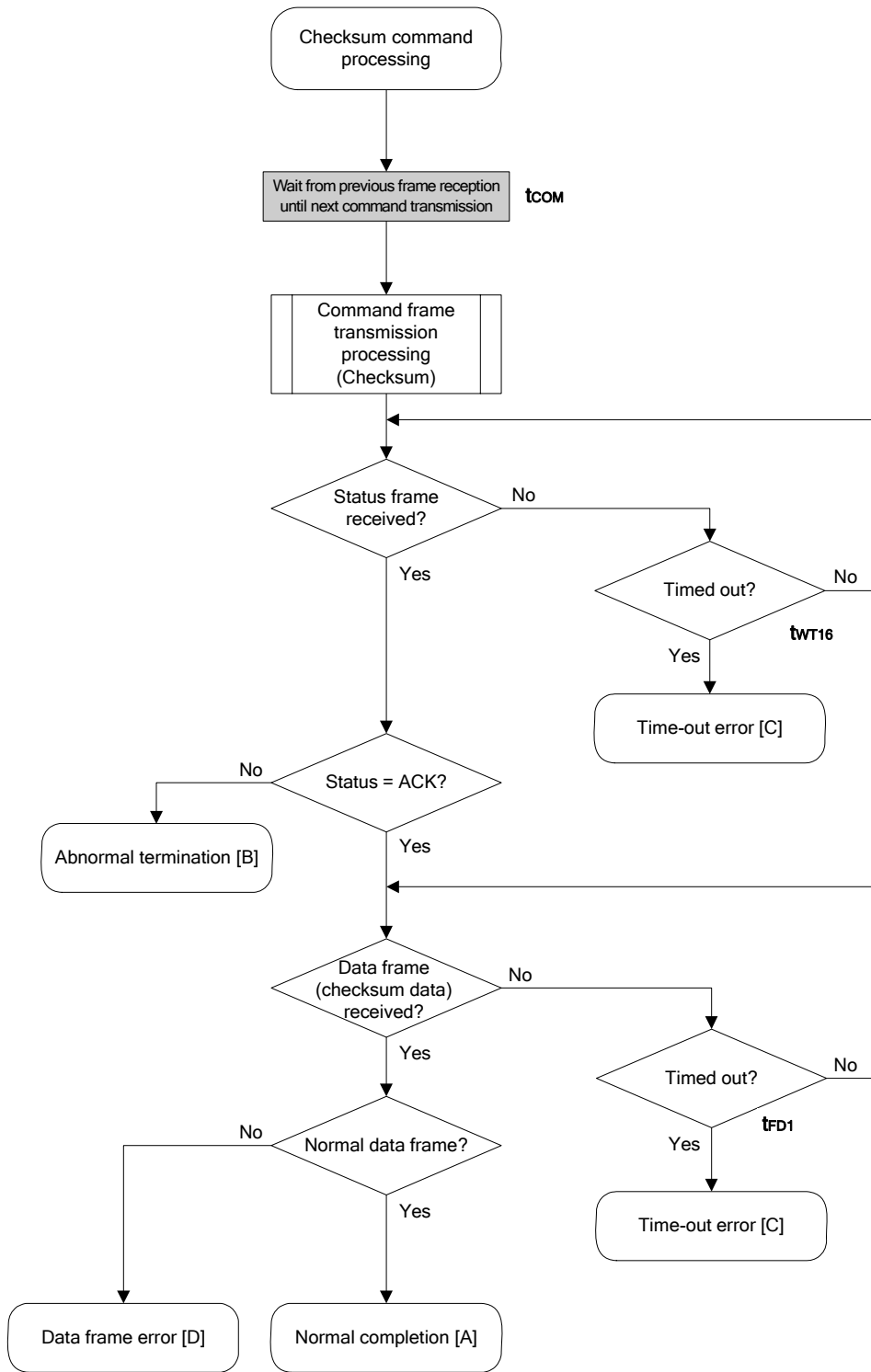
If data frame is abnormal: Data frame error [D]

### 4.13.3 Status at processing completion

Status at Processing Completion		Status Code	Description
Normal completion [A]	Normal acknowledgment (ACK)	06H	The command was executed normally and checksum data was acquired normally.
Abnormal termination [B]	Parameter error	05H	The specified start/end address is out of the flash memory range, or the start/end address is not the start/end address of the block.
	Checksum error	07H	The checksum of the transmitted command frame does not match.
	Negative acknowledgment (NACK)	15H	Command frame data is abnormal (such as invalid data length (LEN) or no ETX).
Time-out error [C]		–	The status frame or data frame was not received within the specified time.
Data frame error [D]		–	The checksum of the data frame received as checksum data does not match.



4.13.4 Flowchart



## 4.13.5 Sample program

The following shows a sample program for Checksum command processing.

```

/*****
/*
/*  Get checksum command
/*
/*****
/*  [i] u16 *sum    ... pointer to checksum save area
/*  [i] u32 top    ... start address
/*  [i] u32 bottom ... end address
/*  [r] u16        ... error code
/*****
u16      fl_ua_getsum(u16 *sum, u32 top, u32 bottom)
{
    u16    rc;

    /*****
    /*      set params
    /*****
    // set params
    set_range_prm(fl_cmd_prm, top, bottom); // set SAH/SAM/SAL, EAH/EAM/EAL

    /*****
    /*      send command
    /*****

    fl_wait(tCOM); // wait before sending command

    put_cmd_ua(FL_COM_GET_CHECK_SUM, 7, fl_cmd_prm); // send GET VERSION command

    rc = get_sfrm_ua(fl_ua_sfrm, tWT16_MAX); // get status frame
    switch(rc) {
        case      FLC_NO_ERR:          break; // continue
    // case      FLC_DFTO_ERR: return rc; break; // case [C]
        default:          return rc;   break; // case [B]
    }

    /*****
    /*      get data frame (Checksum data)
    /*****
    rc = get_dfrm_ua(fl_rxddata_frm, tFD1_MAX); // get status frame
    if (rc){ // if no error,
        return rc; // case [D]
    }

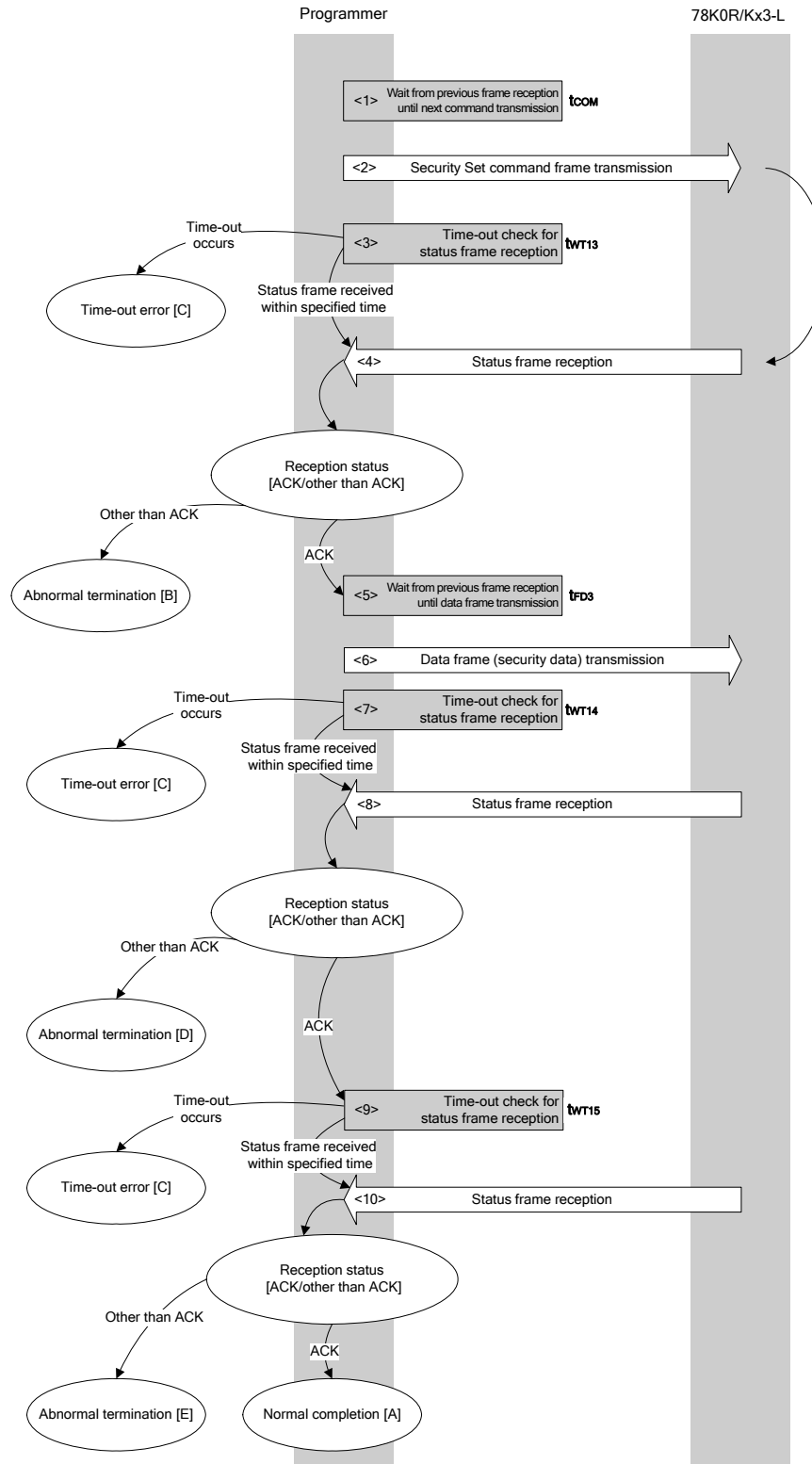
    *sum = (fl_rxddata_frm[OFS_STA_PLD] << 8) + fl_rxddata_frm[OFS_STA_PLD+1]; // set
SUM data
    return rc; // case [A]
}

```

### 4.14 Security Set Command

#### 4.14.1 Processing sequence chart

Security Set command processing sequence



#### 4.14.2 Description of processing sequence

- <1> Waits from the previous frame reception until the next command transmission (wait time  $t_{COM}$ ).
- <2> The Security Set command is transmitted by command frame transmission processing.
- <3> A time-out check is performed from command transmission until status frame reception.  
If a time-out occurs, a time-out error [C] is returned (time-out time  $t_{WT13}$ ).
- <4> The status code is checked.

When ST1 = ACK: Proceeds to <5>.

When ST1 ≠ ACK: Abnormal termination [B]

- <5> Waits from the previous frame reception until the next data frame transmission (wait time  $t_{FD3}$ ).
- <6> The data frame (security setting data) is transmitted by data frame transmission processing.
- <7> A time-out check is performed until status frame reception.  
If a time-out occurs, a time-out error [C] is returned (time-out time  $t_{WT14}$ ).
- <8> The status code is checked.

When ST1 = ACK: Proceeds to <9>.

When ST1 ≠ ACK: Abnormal termination [D]

- <9> A time-out check is performed until status frame reception.  
If a time-out occurs, a time-out error [C] is returned (time-out time  $t_{WT15}$ ).
- <10> The status code is checked.

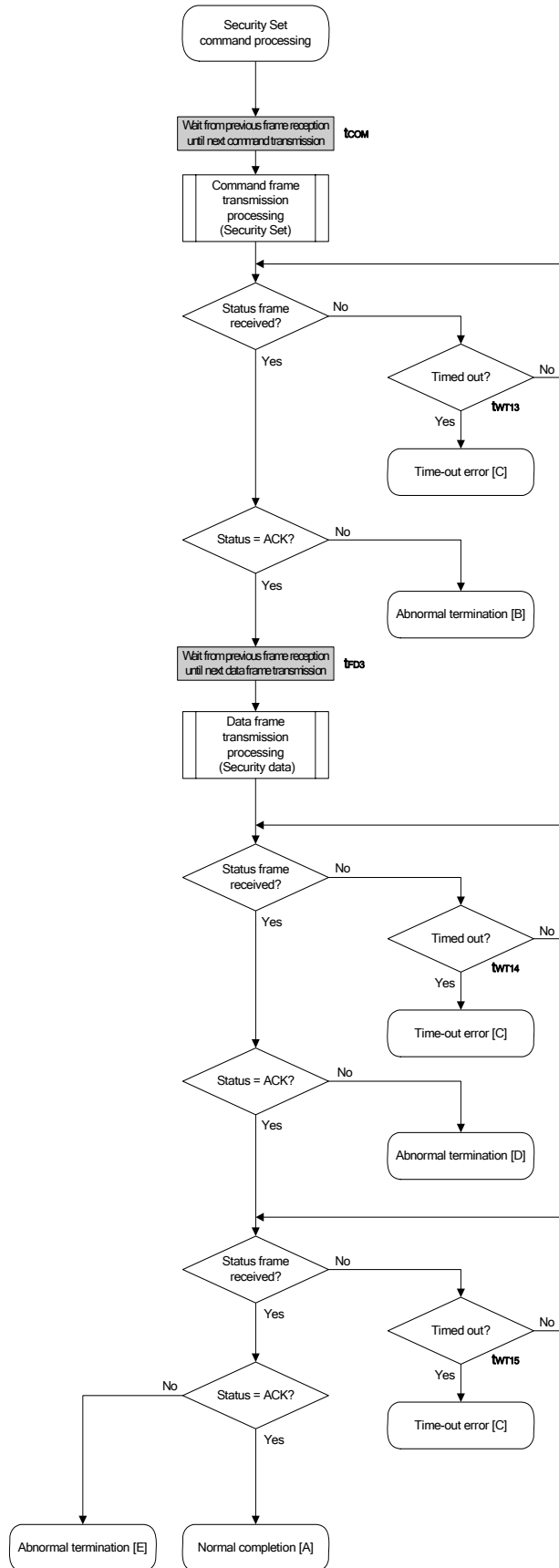
When ST1 = ACK: Normal completion [A]

When ST1 ≠ ACK: Abnormal termination [E]

#### 4.14.3 Status at processing completion

Status at Processing Completion		Status Code	Description
Normal completion [A]	Normal acknowledgment (ACK)	06H	The command was executed normally and security setting data was set normally.
Abnormal termination [B]	Parameter error	05H	Parameter BOT is not 01H, the FSW setting block is not set so that the start block number is larger than the end block number, or the FSW end block number is larger than the last block number.
	Checksum error	07H	The checksum of the transmitted command frame or data frame does not match.
	Protect error	10H	An already prohibited flag is to be enabled.
	Negative acknowledgment (NACK)	15H	Command frame data is abnormal (such as invalid data length (LEN) or no ETX).
Time-out error [C]		–	The status frame or data frame was not received within the specified time.
Abnormal termination [D], [E]	MRG10 error	1AH	Writing security data has failed.
	MRG11 error	1BH	
	Write error	1CH	

4.14.4 Flowchart



## 4.14.5 Sample program

The following shows a sample program for Security Set command processing.

```

/*****/
/*
/*      Set security flag command
/*
/*****/
/* [i] u8 scf      ... Security flag data
/* [r] u16        ... error code
/*****/
u16      fl_ua_setscf(u8 scf, u8 bot, u8 fsws, u8 fswe)
{
    u16    rc;

/*****/
/*      set params
/*
/*****/
fl_cmd_prm[0] = 0x00;           // "BLK" (must be 0x00)
fl_cmd_prm[1] = 0x00;           // "PAG" (must be 0x00)

fl_txdata_frm[0] = scf|= 0b11101000; // "FLG" (bit 7,6,5,3 must be '1')
fl_txdata_frm[1] = bot;           // "BOT"
fl_txdata_frm[2] = 0x00;           // "FSWS High"
fl_txdata_frm[3] = fsws;           // "FSWS Low"
fl_txdata_frm[4] = 0x00;           // "FSWE High"
fl_txdata_frm[5] = fswe;           // "FSWE Low"
fl_txdata_frm[6] = 0xff;           // (must be 0xff)
fl_txdata_frm[7] = 0xff;           // (must be 0xff)

/*****/
/*      send command
/*
/*****/
fl_wait(tCOM);                 // wait before sending
command

put_cmd_ua(FL_COM_SET_SECURITY, 3, fl_cmd_prm);

rc = get_sfrm_ua(fl_ua_sfrm, tWT13_MAX); // get status frame
switch(rc) {
    case      FLC_NO_ERR:           break; // continue
//   case      FLC_DFTO_ERR: return rc; break; // case [C]
    default:           return rc;   break; // case [B]
}

/*****/
/*      send data frame (security setting data)
/*
/*****/

fl_wait(tFD4);
put_dfrm_ua(6, fl_txdata_frm, true); // send securithi setting
data

//rc = get_sfrm_ua(fl_ua_sfrm, tWT14_MAX); // get status frame
rc = get_sfrm_ua(fl_ua_sfrm, tWT14_MAX+100); // get status frame (+100us)

```

```
is overhead)
switch(rc) {
    case      FLC_NO_ERR:          break; // continue
//    case      FLC_DFTO_ERR: return rc;  break; // case [C]
    default:          return rc;  break; // case [B]
}

/*****
/*      Check internally verify      */
*****/
rc = get_sfrm_ua(fl_ua_sfrm, tWT15_MAX);          // get status frame
//switch(rc) {
//
//    case      FLC_NO_ERR:  return rc;  break; // case [A]
//    case      FLC_DFTO_ERR: return rc;  break; // case [C]
//    default:          return rc;  break; // case [B]
//}
return rc;
}
```

## CHAPTER 5 FLASH MEMORY PROGRAMMING PARAMETER CHARACTERISTICS

This chapter describes the characteristics of parameter transmitted between the programmer and the devices (78K0R/Kx3-L, 78K0R/lx3, and 78K0R/Kx3-C) in the flash memory programming mode.

Refer to the user's manual of the 78K0R/Kx3-L, 78K0R/lx3, or 78K0R/Kx3-C for electrical specifications when designing a programmer.

### 5.1 Flash Memory Parameter Characteristics of 78K0R/Kx3-L

#### 5.5.1 Flash memory parameter characteristics in full-speed mode

##### (1) Flash memory programming mode setting time

Parameter	Symbol	MIN.	TYP.	MAX.
$V_{DD}\uparrow$ to FLMD0 $\uparrow$	$t_{DP}$	0		
FLMD0 $\uparrow$ to RESET $\uparrow$	$t_{PR}$	2 ms		
Ready start time from RESET $\uparrow$	$t_{R0}$	4.4 ms		100 ms
Low level data0 (Ready) width <sup>Note</sup>	$t_{L0}$	892 $\mu$ s	937.5 $\mu$ s	987 $\mu$ s
Wait for low level data1	$t_{01}$	110.6 $\mu$ s		
Wait for low level data2	$t_{02}$	4.5 $\mu$ s		
Wait for Read command	$t_{2C}$	608.1 $\mu$ s		
Low level data1/data2 width <sup>Note</sup>	$t_{L1}, t_{L2}$		937.5 $\mu$ s	

**Note** The low-level width is the same as the 00H data width at 9,600 bps. (It includes the start bit and is therefore "0" data of 9 bits.)

$t_{L0}$  is the low-level width of the data transmitted from the 78K0R/Kx3-L firmware.  $t_{L1}$  and  $t_{L2}$  are the low-level widths of the data transmitted from the flash memory programmer.



(2) Programming characteristics

Wait	Condition	Symbol	MIN.	MAX.
Between data frame transmissions	Data frame reception <sup>Note 1</sup>	$t_{DR}$	9.3 $\mu$ S	
			1.9 $\mu$ S	
	Data frame transmission	$t_{DT}$	0 <sup>Note 2</sup>	
From status frame transmission until data frame transmission	—	$t_{FD1}$	0 <sup>Note 2</sup>	
From status frame transmission until data frame reception (1)	Programming command	$t_{FD2}$	2.3 $\mu$ S	
From status frame transmission until data frame reception (2)	Verify command	$t_{FD3}$	83.8 $\mu$ S	
From status frame transmission until data frame reception (3)	Security setting command	$t_{FD4}$	236.2 $\mu$ S	
From status frame transmission until command frame reception	<b>Note 3</b>	$t_{COM}$	13.2 $\mu$ S	
			2.6 $\mu$ S	

- Notes 1.** Use the value in the upper row after the flash memory programming mode is set and until the transmission of the Baud Rate Set command of the programmer finishes. Use the value in the lower row from the time of the transmission of the Reset command after the Baud Rate Set command is transmitted.
- 2.** Enable successive reception for the programmer. Also, set the time-out time for the programmer to at least 3 seconds.
- 3.** Use the value in the upper row after the flash memory programming mode is set and until the transmission of the Reset command finishes after the Baud Rate Set command of the programmer is transmitted. Use the value in the lower row from the time of the transmission of the Reset command after the Baud Rate Set command is transmitted.

**Remark** The waits are defined as follows.

< $t_{DR}$ ,  $t_{FD2}$ ,  $t_{FD3}$ ,  $t_{FD4}$ ,  $t_{COM}$ >

The 78K0R/Kx3-L can execute the next communication after the MIN. time has elapsed after completion of the current communication.

The programmer needs to transmit the following data within the period from the MIN. time to the MAX. time after completion of the current communication.

The MAX. time is not specified, but use approximately 3 seconds.

< $t_{DT}$ ,  $t_{FD1}$ >

The 78K0R/Kx3-L can execute the next communication after the MIN. time has elapsed after completion of the current communication.

The programmer needs to be ready for reception of the following data within the MIN. time after completion of the current communication.

The MAX. time is not specified, but set the timeout to approximately 3 seconds.

## (3) Command characteristics

Command	Symbol	Condition	MIN.	MAX.
Reset	t <sub>WT0</sub>	—	0 <sup>Note 1</sup>	
Chip Erase	t <sub>WT1</sub>	—	(34.8 + 1.8 × total number of blocks) ms	(877.8 + 56.3 × total number of blocks) ms
Block Erase	t <sub>WT2</sub> <sup>Note 2</sup>	—	10.6 ms	(0.8 + 251.9 × execution count of simultaneous selection and erasure + 55.0 × number of blocks to be erased) ms
Programming	t <sub>WT3</sub>	—	0 <sup>Note 1</sup>	
	t <sub>WT4</sub> <sup>Note 3</sup>	—	1.6 ms	41.9 ms
	t <sub>WT5</sub> <sup>Note 4</sup>	Block 0	30.7 ms	633.5 ms
		Other than block 0	4.3 ms	6.7 ms
Verify	t <sub>WT6</sub>	—	0 <sup>Note 1</sup>	
	t <sub>WT7</sub> <sup>Note 3</sup>	—	0 <sup>Note 1</sup>	
Block Blank Check	t <sub>WT8</sub> <sup>Note 4</sup>	—	2.0 ms	3.7 ms
Baud Rate Set	t <sub>WT10</sub>	—	205.3 μs	
Silicon Signature	t <sub>WT11</sub>	—	0 <sup>Note 1</sup>	
Version Get	t <sub>WT12</sub>	—	0 <sup>Note 1</sup>	
Security Set	t <sub>WT13</sub>	—	0 <sup>Note 1</sup>	
	t <sub>WT14</sub>	—	7.5 μs	14.1 μs
	t <sub>WT15</sub>	—	2.6 μs	626.8 ms
Checksum	t <sub>WT16</sub>	—	0 <sup>Note 1</sup>	

- Notes 1.** Reception must be enabled for the programmer before command frame transmission.
- 2.** See **5.4 Simultaneous Selection and Erasure Performed by Block Erase Command** for the calculation method of the execution count of simultaneous selection and erasure.
- 3.** Time for 256-byte data transmission
- 4.** Time for one block transmission

**Remark** The waits are defined as follows.

<t<sub>WT0</sub> to t<sub>WT8</sub>, t<sub>WT11</sub> to t<sub>WT16</sub>>

The 78K0R/Kx3-L completes command processing between the MIN. and MAX. times and transmits a status frame.

For commands with a specified MAX. time, the programmer must wait for the start bit of the reception frame until the MAX. time has elapsed and then perform time-out processing.

For commands without a specified MAX. time, set the timeout to approximately 3 seconds.

<t<sub>WT10</sub>>

The 78K0R/Kx3-L can execute the next communication after the MIN. time has elapsed after completion of the current communication.

The programmer needs to transmit the following data within the period from the MIN. time to the MAX. time after completion of the current communication.

The MAX. time is not specified, but use approximately 3 seconds.

## 5.1.2 Flash memory parameter characteristics in wide-voltage mode

## (1) Flash memory programming mode setting time

Parameter	Symbol	MIN.	TYP.	MAX.
$V_{DD}\uparrow$ to FLMD0 $\uparrow$	$t_{DP}$	0		
FLMD0 $\uparrow$ to $\overline{RESET}\uparrow$	$t_{PR}$	2 ms		
Ready start time from $\overline{RESET}\uparrow$	$t_{R0}$	4.4 ms		100 ms
Low level data0 (Ready) width <sup>Note</sup>	$t_{L0}$	892 $\mu$ s	937.5 $\mu$ s	987 $\mu$ s
Wait for low level data1	$t_{01}$	110.6 $\mu$ s		
Wait for low level data2	$t_{02}$	4.5 $\mu$ s		
Wait for Read command	$t_{2C}$	608.1 $\mu$ s		
Low level data1/data2 width <sup>Note</sup>	$t_{L1}, t_{L2}$		937.5 $\mu$ s	

**Note** The low-level width is the same as the 00H data width at 9,600 bps. (It includes the start bit and is therefore "0" data of 9 bits.)

$t_{L0}$  is the low-level width of the data transmitted from the 78K0R/Kx3-L firmware.  $t_{L1}$  and  $t_{L2}$  are the low-level widths of the data transmitted from the flash programmer.

## (2) Programming characteristics

Wait	Condition	Symbol	MIN.	MAX.
Between data frame transmissions	Data frame reception	$t_{DR}$	9.3 $\mu s$	
	Data frame transmission	$t_{DT}$	0 <sup>Note</sup>	
From status frame transmission until data frame transmission	—	$t_{FD1}$	0 <sup>Note</sup>	
From status frame transmission until data frame reception (1)	Programming command	$t_{FD2}$	11.4 $\mu s$	
From status frame transmission until data frame reception (2)	Verify command	$t_{FD3}$	416.4 $\mu s$	
From status frame transmission until data frame reception (3)	Security setting command	$t_{FD4}$	985.8 $\mu s$	
From status frame transmission until command frame reception	—	$t_{COM}$	13.2 $\mu s$	

**Note** Enable successive reception for the programmer.

**Remark** The waits are defined as follows.

< $t_{DR}$ ,  $t_{FD2}$ ,  $t_{FD3}$ ,  $t_{FD4}$ ,  $t_{COM}$ >

The 78K0R/Kx3-L can execute the next communication after the MIN. time has elapsed after completion of the current communication.

The programmer needs to transmit the following data within the period from the MIN. time to the MAX. time after completion of the current communication.

The MAX. time is not specified, but use approximately 3 seconds.

< $t_{DT}$ ,  $t_{FD1}$ >

The 78K0R/Kx3-L can execute the next communication after the MIN. time has elapsed after completion of the current communication.

The programmer needs to be ready for reception of the following data within the MIN. time after completion of the current communication.

The MAX. time is not specified, but set the timeout to approximately 3 seconds.

## (3) Command characteristics

Command	Symbol	Condition	MIN.	MAX.
Reset	t <sub>WT0</sub>	–	0 <sup>Note 1</sup>	
Chip Erase	t <sub>WT1</sub>	–	(76.0 + 9.3 × total number of blocks) ms	(1420.1 + 281.1 × total number of blocks) ms
Block Erase	t <sub>WT2</sub> <sup>Note 2</sup>	–	20.3 ms	(3.3 + 271.6 × execution count of simultaneous selection and erasure + 275.0 × number of blocks to be erased) ms
Programming	t <sub>WT3</sub>	–	0 <sup>Note 1</sup>	
	t <sub>WT4</sub> <sup>Note 3</sup>	–	6.6 ms	149.9 ms
	t <sub>WT5</sub> <sup>Note 4</sup>	Block 0	89.8 ms	1187.5 ms
Other than block 0		23.1 ms	34.9 ms	
Verify	t <sub>WT6</sub>	–	0 <sup>Note 1</sup>	
	t <sub>WT7</sub> <sup>Note 3</sup>	–	0 <sup>Note 1</sup>	
Block Blank Check	t <sub>WT8</sub> <sup>Note 4</sup>	–	9.9 ms	18.0 ms
Baud Rate Set	t <sub>WT10</sub>	–	379.2 μs	
Silicon Signature	t <sub>WT11</sub>	–	0 <sup>Note 1</sup>	
Version Get	t <sub>WT12</sub>	–	0 <sup>Note 1</sup>	
Security Set	t <sub>WT13</sub>	–	0 <sup>Note 1</sup>	
	t <sub>WT14</sub>	–	37.6 μs	70.2 μs
	t <sub>WT15</sub>	–	13.4 μs	1152.3 ms
Checksum	t <sub>WT16</sub>	–	0 <sup>Note 1</sup>	

**Notes 1.** Reception must be enabled for the programmer before command frame transmission.

**2.** See 5.4 **Simultaneous Selection and Erasure Performed by Block Erase Command** for the calculation method of the execution count of simultaneous selection and erasure.

**3.** Time for 256-byte data transmission

**4.** Time for one block transmission

**Remark** The waits are defined as follows.

<t<sub>WT0</sub> to t<sub>WT8</sub>, t<sub>WT11</sub> to t<sub>WT16</sub>>

The 78K0R/Kx3-L completes command processing between the MIN. and MAX. times and transmits a status frame.

For commands with a specified MAX. time, the programmer must wait for the start bit of the reception frame until the MAX. time has elapsed and then perform time-out processing.

For commands without a specified MAX. time, set the timeout to approximately 3 seconds.

<t<sub>WT10</sub>>

The 78K0R/Kx3-L can execute the next communication after the MIN. time has elapsed after completion of the current communication.

The programmer needs to transmit the following data within the period from the MIN. time to the MAX. time after completion of the current communication.

The MAX. time is not specified, but use approximately 3 seconds.

## 5.2 Flash Memory Parameter Characteristics of 78K0R/Ix3

### (1) Flash memory programming mode setting time

Parameter	Symbol	MIN.	TYP.	MAX.
$V_{DD}\uparrow$ to FLMD0 $\uparrow$	$t_{DP}$	0		
FLMD0 $\uparrow$ to $\overline{\text{RESET}}\uparrow$	$t_{PR}$	2 ms		
Ready start time from $\overline{\text{RESET}}\uparrow$	$t_{R0}$	4.4 ms		100 ms
Low level data0 (Ready) width <sup>Note</sup>	$t_{L0}$	892 $\mu\text{s}$	937.5 $\mu\text{s}$	987 $\mu\text{s}$
Wait for low level data1	$t_{01}$	110.6 $\mu\text{s}$		
Wait for low level data2	$t_{02}$	4.5 $\mu\text{s}$		
Wait for Read command	$t_{2C}$	608.1 $\mu\text{s}$		
Low level data1/data2 width <sup>Note</sup>	$t_{L1}, t_{L2}$		937.5 $\mu\text{s}$	

**Note** The low-level width is the same as the 00H data width at 9,600 bps. (It includes the start bit and is therefore "0" data of 9 bits.)

$t_{L0}$  is the low-level width of the data transmitted from the 78K0R/Ix3 firmware.  $t_{L1}$  and  $t_{L2}$  are the low-level widths of the data transmitted from the flash programmer.

## (2) Programming characteristics

Wait	Condition	Symbol	MIN.	MAX.
Between data frame transmissions	Data frame reception <sup>Note 1</sup>	$t_{DR}$	9.3 $\mu$ S	
			1.9 $\mu$ S	
	Data frame transmission	$t_{DT}$	0 <sup>Note 2</sup>	
From status frame transmission until data frame transmission	—	$t_{FD1}$	0 <sup>Note 2</sup>	
From status frame transmission until data frame reception (1)	Programming command	$t_{FD2}$	2.3 $\mu$ S	
From status frame transmission until data frame reception (2)	Verify command	$t_{FD3}$	83.8 $\mu$ S	
From status frame transmission until data frame reception (3)	Security setting command	$t_{FD4}$	236.2 $\mu$ S	
From status frame transmission until command frame reception	<b>Note 3</b>	$t_{COM}$	13.2 $\mu$ S	
			2.6 $\mu$ S	

**Notes 1.** Use the value in the upper row after the flash memory programming mode is set and until the transmission of the Baud Rate Set command of the programmer finishes. Use the value in the lower row from the time of the transmission of the Reset command after the Baud Rate Set command is transmitted.

2. Enable successive reception for the programmer. Also, set the time-out time for the programmer to at least 3 seconds.
3. Use the value in the upper row after the flash memory programming mode is set and until the transmission of the Reset command finishes after the Baud Rate Set command of the programmer is transmitted. Use the value in the lower row from the time of the transmission of the Reset command after the Baud Rate Set command is transmitted.

**Remark** The waits are defined as follows.

< $t_{DR}$ ,  $t_{FD2}$ ,  $t_{FD3}$ ,  $t_{FD4}$ ,  $t_{COM}$ >

The 78K0R/lx3 can execute the next communication after the MIN. time has elapsed after completion of the current communication.

The programmer needs to transmit the following data within the period from the MIN. time to the MAX. time after completion of the current communication.

The MAX. time is not specified, but use approximately 3 seconds.

< $t_{DT}$ ,  $t_{FD1}$ >

The 78K0R/lx3 can execute the next communication after the MIN. time has elapsed after completion of the current communication.

The programmer needs to be ready for reception of the following data within the MIN. time after completion of the current communication.

The MAX. time is not specified, but set the timeout to approximately 3 seconds.

## (3) Command characteristics

Command	Symbol	Condition	MIN.	MAX.
Reset	t <sub>WT0</sub>	–	0 <sup>Note 1</sup>	
Chip Erase	t <sub>WT1</sub>	–	(34.8 + 1.8 × total number of blocks) ms	(877.8 + 56.3 × total number of blocks) ms
Block Erase	t <sub>WT2</sub> <sup>Note 2</sup>	–	10.6 ms	(0.8 + 251.9 × execution count of simultaneous selection and erasure + 55.0 × number of blocks to be erased) ms
Programming	t <sub>WT3</sub>	–	0 <sup>Note 1</sup>	
	t <sub>WT4</sub> <sup>Note 3</sup>	–	1.6 ms	41.9 ms
	t <sub>WT5</sub> <sup>Note 4</sup>	Block 0	30.7 ms	633.5 ms
Other than block 0		4.3 ms	6.7 ms	
Verify	t <sub>WT6</sub>	–	0 <sup>Note 1</sup>	
	t <sub>WT7</sub> <sup>Note 3</sup>	–	0 <sup>Note 1</sup>	
Block Blank Check	t <sub>WT8</sub> <sup>Note 4</sup>	–	2.0 ms	3.7 ms
Baud Rate Set	t <sub>WT10</sub>	–	205.3 μs	
Silicon Signature	t <sub>WT11</sub>	–	0 <sup>Note 1</sup>	
Version Get	t <sub>WT12</sub>	–	0 <sup>Note 1</sup>	
Security Set	t <sub>WT13</sub>	–	0 <sup>Note 1</sup>	
	t <sub>WT14</sub>	–	7.5 μs	14.1 μs
	t <sub>WT15</sub>	–	2.6 μs	626.8 ms
Checksum	t <sub>WT16</sub>	–	0 <sup>Note 1</sup>	

**Notes 1.** Reception must be enabled for the programmer before command frame transmission.

**2.** See 5.4 **Simultaneous Selection and Erasure Performed by Block Erase Command** for the calculation method of the execution count of simultaneous selection and erasure.

**3.** Time for 256-byte data transmission

**4.** Time for one block transmission

**Remark** The waits are defined as follows.

<t<sub>WT0</sub> to t<sub>WT8</sub>, t<sub>WT11</sub> to t<sub>WT16</sub>>

The 78K0R/Ix3 completes command processing between the MIN. and MAX. times and transmits a status frame.

For commands with a specified MAX. time, the programmer must wait for the start bit of the reception frame until the MAX. time has elapsed and then perform time-out processing.

For commands without a specified MAX. time, set the timeout to approximately 3 seconds.

<t<sub>WT10</sub>>

The 78K0R/Ix3 can execute the next communication after the MIN. time has elapsed after completion of the current communication.

The programmer needs to transmit the following data within the period from the MIN. time to the MAX. time after completion of the current communication.

The MAX. time is not specified, but use approximately 3 seconds.



### 5.3 Flash Memory Parameter Characteristics of 78K0R/Kx3-C

#### (1) Flash memory programming mode setting time

Parameter	Symbol	MIN.	TYP.	MAX.
$V_{DD}\uparrow$ to $FLMD0\uparrow$	$t_{DP}$	0		
$FLMD0\uparrow$ to $\overline{RESET}\uparrow$	$t_{PR}$	2 ms		
Ready start time from $\overline{RESET}\uparrow$	$t_{R0}$	4.4 ms		100 ms
Low level data0 (Ready) width <sup>Note</sup>	$t_{L0}$	892 $\mu$ s	937.5 $\mu$ s	987 $\mu$ s
Wait for low level data1	$t_{01}$	110.6 $\mu$ s		
Wait for low level data2	$t_{02}$	4.5 $\mu$ s		
Wait for Read command	$t_{2C}$	608.1 $\mu$ s		
Low level data1/data2 width <sup>Note</sup>	$t_{L1}, t_{L2}$		937.5 $\mu$ s	

**Note** The low-level width is the same as the 00H data width at 9,600 bps. (It includes the start bit and is therefore "0" data of 9 bits.)

$t_{L0}$  is the low-level width of the data transmitted from the 78K0R/Kx3-C firmware.  $t_{L1}$  and  $t_{L2}$  are the low-level widths of the data transmitted from the flash programmer.

## (2) Programming characteristics

Wait	Condition	Symbol	MIN.	MAX.
Between data frame transmissions	Data frame reception <sup>Note 1</sup>	$t_{DR}$	9.3 $\mu s$	
			1.9 $\mu s$	
	Data frame transmission	$t_{DT}$	0 <sup>Note 2</sup>	
From status frame transmission until data frame transmission	—	$t_{FD1}$	0 <sup>Note 2</sup>	
From status frame transmission until data frame reception (1)	Programming command	$t_{FD2}$	2.3 $\mu s$	
From status frame transmission until data frame reception (2)	Verify command	$t_{FD3}$	83.8 $\mu s$	
From status frame transmission until data frame reception (3)	Security setting command	$t_{FD4}$	236.2 $\mu s$	
From status frame transmission until command frame reception	<b>Note 3</b>	$t_{COM}$	13.2 $\mu s$	
			2.6 $\mu s$	

- Notes 1.** Use the value in the upper row after the flash memory programming mode is set and until the transmission of the Baud Rate Set command of the programmer finishes. Use the value in the lower row from the time of the transmission of the Reset command after the Baud Rate Set command is transmitted.
- 2.** Enable successive reception for the programmer. Also, set the time-out time for the programmer to at least 3 seconds.
- 3.** Use the value in the upper row after the flash memory programming mode is set and until the transmission of the Reset command finishes after the Baud Rate Set command of the programmer is transmitted. Use the value in the lower row from the time of the transmission of the Reset command after the Baud Rate Set command is transmitted.

**Remark** The waits are defined as follows.

< $t_{DR}$ ,  $t_{FD2}$ ,  $t_{FD3}$ ,  $t_{FD4}$ ,  $t_{COM}$ >

The 78K0R/Kx3-C can execute the next communication after the MIN. time has elapsed after completion of the current communication.

The programmer needs to transmit the following data within the period from the MIN. time to the MAX. time after completion of the current communication.

The MAX. time is not specified, but use approximately 3 seconds.

< $t_{DT}$ ,  $t_{FD1}$ >

The 78K0R/Kx3-C can execute the next communication after the MIN. time has elapsed after completion of the current communication.

The programmer needs to be ready for reception of the following data within the MIN. time after completion of the current communication.

The MAX. time is not specified, but set the timeout to approximately 3 seconds.

## (3) Command characteristics

Command	Symbol	Condition	MIN.	MAX.
Reset	t <sub>WT0</sub>	—	0 <sup>Note 1</sup>	
Chip Erase	t <sub>WT1</sub>	—	(34.8 + 1.8 × total number of blocks) ms	(877.8 + 56.3 × total number of blocks) ms
Block Erase	t <sub>WT2</sub> <sup>Note 2</sup>	—	10.6 ms	(0.8 + 251.9 × execution count of simultaneous selection and erasure + 55.0 × number of blocks to be erased) ms
Programming	t <sub>WT3</sub>	—	0 <sup>Note 1</sup>	
	t <sub>WT4</sub> <sup>Note 3</sup>	—	1.6 ms	41.9 ms
	t <sub>WT5</sub> <sup>Note 4</sup>	Block 0	30.7 ms	633.5 ms
Other than block 0		4.3 ms	6.7 ms	
Verify	t <sub>WT6</sub>	—	0 <sup>Note 1</sup>	
	t <sub>WT7</sub> <sup>Note 3</sup>	—	0 <sup>Note 1</sup>	
Block Blank Check	t <sub>WT8</sub> <sup>Note 4</sup>	—	2.0 ms	3.7 ms
Baud Rate Set	t <sub>WT10</sub>	—	205.3 μs	
Silicon Signature	t <sub>WT11</sub>	—	0 <sup>Note 1</sup>	
Version Get	t <sub>WT12</sub>	—	0 <sup>Note 1</sup>	
Security Set	t <sub>WT13</sub>	—	0 <sup>Note 1</sup>	
	t <sub>WT14</sub>	—	7.5 μs	14.1 μs
	t <sub>WT15</sub>	—	2.6 μs	626.8 ms
Checksum	t <sub>WT16</sub>	—	0 <sup>Note 1</sup>	

**Notes 1.** Reception must be enabled for the programmer before command frame transmission.

**2.** See 5.4 **Simultaneous selection and erasure performed by Block Erase command** for the calculation method of the execution count of simultaneous selection and erasure.

**3.** Time for 256-byte data transmission

**4.** Time for one block transmission

**Remark** The waits are defined as follows.

<t<sub>WT0</sub> to t<sub>WT8</sub>, t<sub>WT11</sub> to t<sub>WT16</sub>>

The 78K0R/Kx3-C completes command processing between the MIN. and MAX. times and transmits a status frame.

For commands with a specified MAX. time, the programmer must wait for the start bit of the reception frame until the MAX. time has elapsed and then perform time-out processing.

For commands without a specified MAX. time, set the timeout to approximately 3 seconds.

<t<sub>WT10</sub>>

The 78K0R/Kx3-C can execute the next communication after the MIN. time has elapsed after completion of the current communication.

The programmer needs to transmit the following data within the period from the MIN. time to the MAX. time after completion of the current communication.

The MAX. time is not specified, but use approximately 3 seconds.

## 5.4 Simultaneous Selection and Erasure Performed by Block Erase Command

The Block Erase command of the 78K0R/Kx3-L, 78K0R/lx3, and 78K0R/Kx3-C is executed by repeating “simultaneous selection and erasure”, which erases multiple blocks simultaneously.

The wait time inserted during Block Erase command execution is therefore equal to the total execution time of “simultaneous selection and erasure”.

To calculate the “total execution time of simultaneous selection and erasure”, the execution count (M) of the simultaneous selection and erasure must first be calculated.

“M” is calculated by obtaining the number of blocks to be erased simultaneously (number of blocks to be selected and erased simultaneously).

The following describes the method for calculating the number of blocks to be selected and erased simultaneously and the execution count (M).

### 5.4.1 Calculation of number of blocks to be selected and erased simultaneously

The number of blocks to be selected and erased simultaneously should be 1, 2, 4, 8, 16, 32, 64, or 128, depending on which satisfies all of the following conditions.

**[Condition 1]**

(Number of blocks to be erased)  $\geq$  (Number of blocks to be selected and erased simultaneously)

**[Condition 2]**

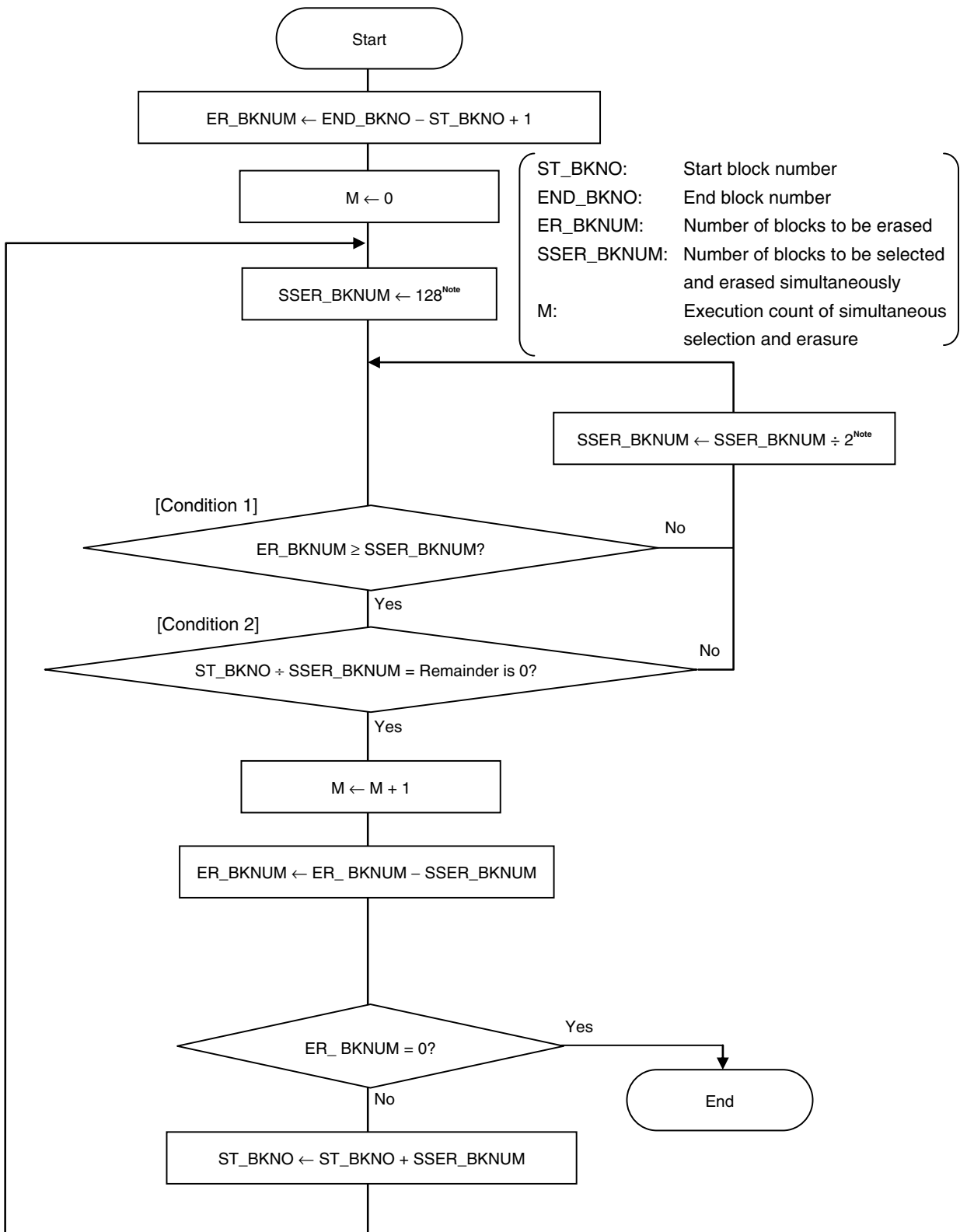
(Start block number)  $\div$  (Number of blocks to be selected and erased simultaneously) = Remainder is 0

**[Condition 3]**

The maximum value among the values that satisfy both Conditions 1 and 2

5.4.2 Calculation of the execution count (M) of simultaneous selection and erasure

Calculation of the execution count (M) is illustrated in the following flowchart.



**Note** Based on the maximum value of SSER\_BKNUM (128), obtain the value that satisfies Conditions 1 and 2 by executing SSER\_BKNUM ÷ 2; Condition 3 is then satisfied.

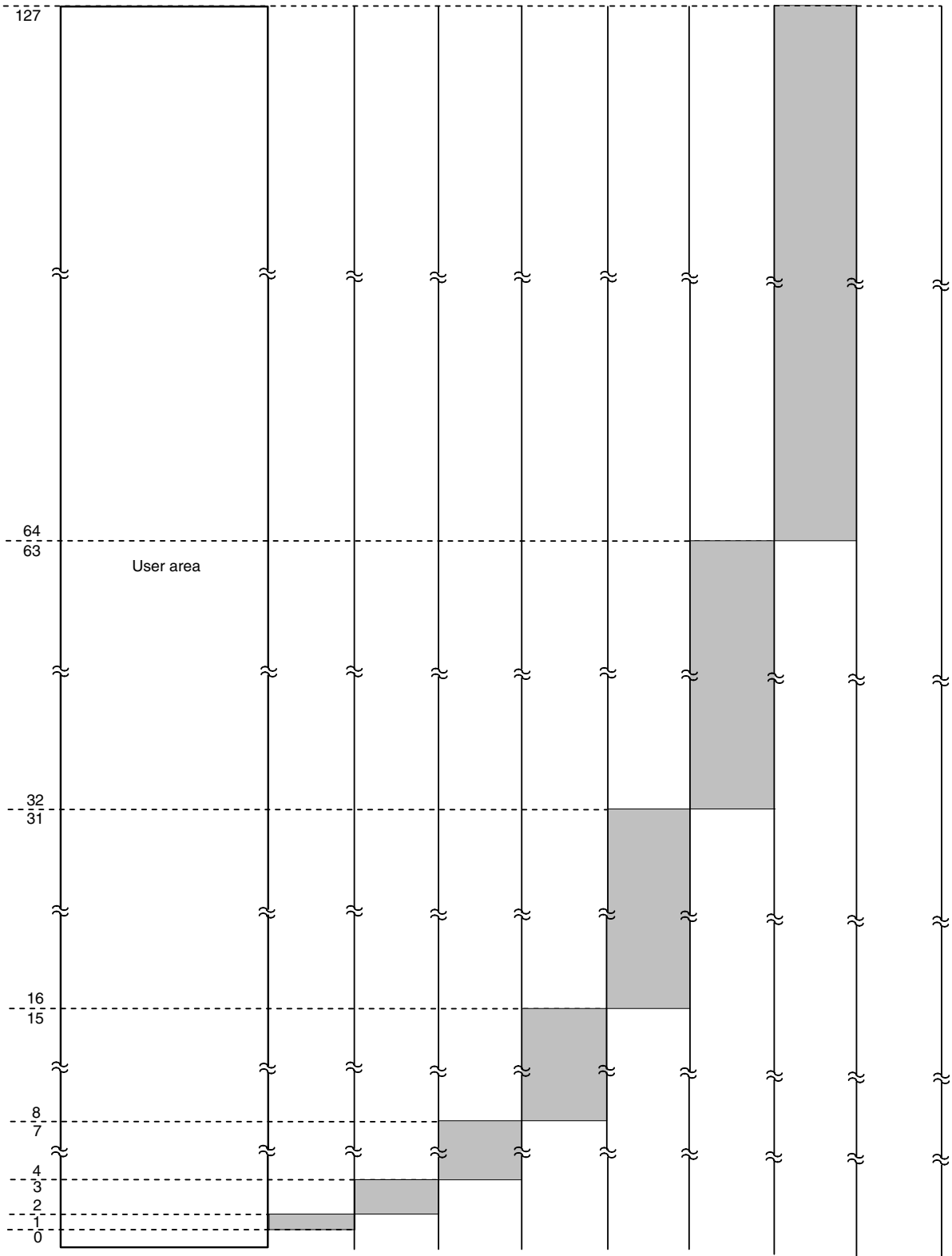
**Example 1** Erasing blocks 1 to 127 ( $N$  (number of blocks to be erased) = 127)

- <1> The first start block number is 1 and the number of blocks to be erased is 127; the values that satisfy Condition 1 are therefore 1, 2, 4, 8, 16, 32, and 64.  
Moreover, the value that satisfies Condition 2 is 1 and the value that satisfies Condition 3 is 1, so the number of blocks to be selected and erased simultaneously is 1; only block 1 is then erased.
- <2> After block 1 is erased, the next start block number is 2 and the number of blocks to be erased is 126; the values that satisfy Condition 1 are therefore 1, 2, 4, 8, 16, 32, and 64.  
Moreover, the values that satisfy Condition 2 are 1 and 2, the value that satisfies Condition 3 is 2, so the number of blocks to be selected and erased simultaneously is 2; blocks 2 and 3 are then erased.
- <3> After blocks 2 and 3 are erased, the next start block number is 4 and the number of blocks to be erased is 124; the values that satisfy Condition 1 are therefore 1, 2, 4, 8, 16, 32, and 64.  
Moreover, the values that satisfy Condition 2 are 1, 2, and 4, the value that satisfies Condition 3 is 4, so the number of blocks to be selected and erased simultaneously is 4; blocks 4 to 7 are then erased.
- <4> After blocks 4 to 7 are erased, the next start block number is 8 and the number of blocks to be erased is 120; the values that satisfy Condition 1 are therefore 1, 2, 4, 8, 16, 32, and 64.  
Moreover, the values that satisfy Condition 2 are 1, 2, 4, and 8, the value that satisfies Condition 3 is 8, so the number of blocks to be selected and erased simultaneously is 8; blocks 8 to 15 are then erased.
- <5> After blocks 8 to 15 are erased, the next start block number is 16 and the number of blocks to be erased is 112; the values that satisfy Condition 1 are therefore 1, 2, 4, 8, 16, 32, and 64.  
Moreover, the values that satisfy Condition 2 are 1, 2, 4, 8, and 16, the value that satisfies Condition 3 is 16, so the number of blocks to be selected and erased simultaneously is 16; blocks 16 to 31 are then erased.
- <6> After blocks 16 to 31 are erased, the next start block number is 32 and the number of blocks to be erased is 96; the values that satisfy Condition 1 are therefore 1, 2, 4, 8, 16, 32, and 64.  
Moreover, the values that satisfy Condition 2 are 1, 2, 4, 8, 16, and 32, the value that satisfies Condition 3 is 32, so the number of blocks to be selected and erased simultaneously is 32; blocks 32 to 63 are then erased.
- <7> After blocks 32 to 63 are erased, the next start block number is 64 and the number of blocks to be erased is 64; the values that satisfy Condition 1 are therefore 1, 2, 4, 8, 16, 32, and 64.  
Moreover, the values that satisfy Condition 2 are 1, 2, 4, 8, 16, 32, and 64, the value that satisfies Condition 3 is 64, so the number of blocks to be selected and erased simultaneously is 64; blocks 64 to 127 are then erased.

Therefore, simultaneous selection and erasure is executed seven times (1, 2 and 3, 4 to 7, 8 to 15, 16 to 31, 32 to 63, and 64 to 127) to erase blocks 1 to 127, so  $M = 7$  is obtained.

Block configuration when executing simultaneous selection and erasure (when erasing blocks 1 to 127)

<Block number>



<Range of blocks that can be selected and erased simultaneously>

**Example 2** Erasing blocks 5 to 10 ( $N$  (number of blocks to be erased) = 6)

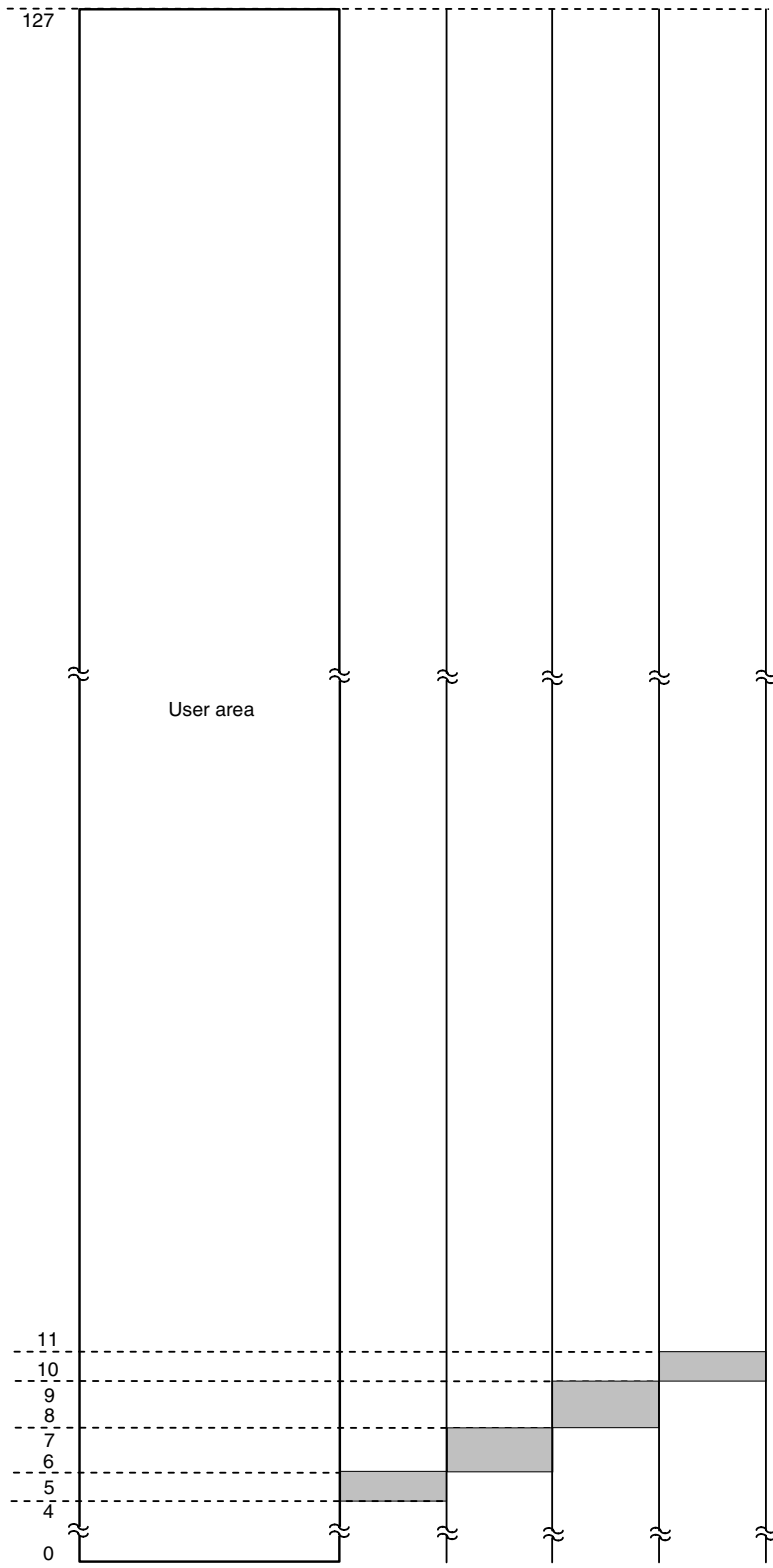
- <1> The first start block number is 5 and the number of blocks to be erased is 6; the values that satisfy Condition 1 are therefore 1, 2, and 4.  
Moreover, the value that satisfies Condition 2 is 1 and the value that satisfies Condition 3 is 1, so the number of blocks to be selected and erased simultaneously is 1; only block 5 is the erased.
  
- <2> After block 5 is erased, the next start block number is 6 and the number of blocks to be erased is 5; the values that satisfy Condition 1 are therefore 1, 2, and 4.  
Moreover, the values that satisfy Condition 2 are 1 and 2, the value that satisfies Condition 3 is 2, so the number of blocks to be selected and erased simultaneously is 2; blocks 6 and 7 are then erased.
  
- <3> After blocks 6 and 7 are erased, the next start block number is 8 and the number of blocks to be erased is 3; the values that satisfy Condition 1 are therefore 1 and 2.  
Moreover, the values that satisfy Condition 2 are 1 and 2, the value that satisfies Condition 3 is 2, so the number of blocks to be selected and erased simultaneously is 2; blocks 8 and 9 are then erased.
  
- <4> After blocks 8 and 9 are erased, the next start block number is 10 and the number of blocks to be erased is 1; the value that satisfies Condition 1 is therefore 1. This also satisfies Conditions 2 and 3, so the number of blocks to be selected and erased simultaneously is 1; block 10 is then erased.

Therefore, simultaneous selection and erasure is executed four times (5, 6 and 7, 8 and 9, and 10) to erase blocks 5 to 10, so  $M = 4$  is obtained.



Block configuration when executing simultaneous selection and erasure (when erasing blocks 5 to 10)

<Block number>



<Range of blocks that can be selected and erased simultaneously>

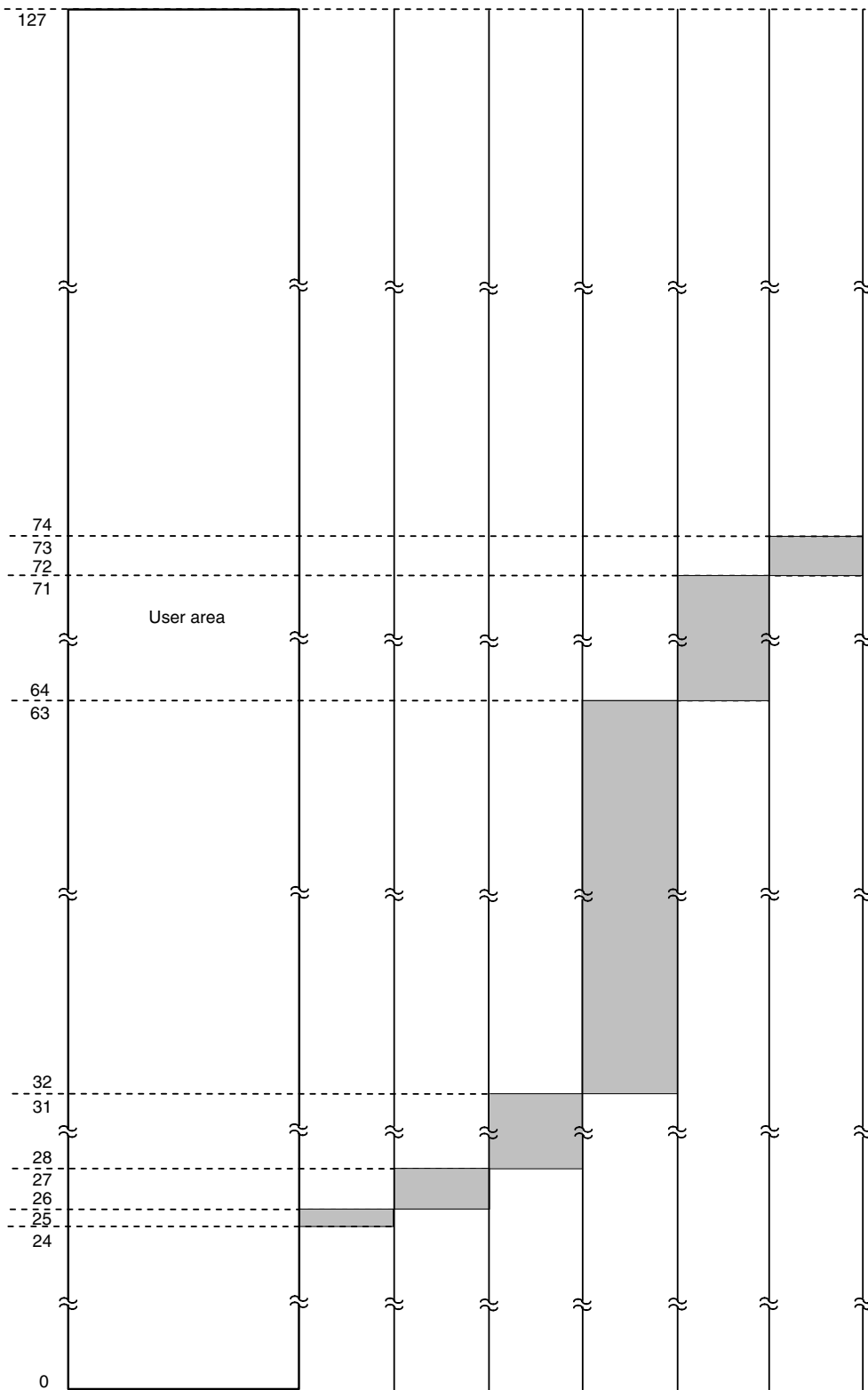
**Example 3** Erasing blocks 25 to 73 ( $N$  (number of blocks to be erased) = 49)

- <1> The first start block number is 25 and the number of blocks to be erased is 49; the values that satisfy Condition 1 are therefore 1, 2, 4, 8, 16, and 32.  
Moreover, the value that satisfies Condition 2 is 1 and the value that satisfies Condition 3 is 1, so the number of blocks to be selected and erased simultaneously is 1; only block 25 is then erased.
- <2> After block 25 is erased, the next start block number is 26 and the number of blocks to be erased is 48; the values that satisfy Condition 1 are therefore 1, 2, 4, 8, 16, and 32.  
Moreover, the values that satisfy Condition 2 are 1 and 2, the value that satisfies Condition 3 is 2, so the number of blocks to be selected and erased simultaneously is 2; blocks 26 and 27 are then erased.
- <3> After blocks 26 and 27 are erased, the next start block number is 28 and the number of blocks to be erased is 46; the values that satisfy Condition 1 are therefore 1, 2, 4, 8, 16, and 32.  
Moreover, the values that satisfy Condition 2 are 1, 2, and 4, the value that satisfies Condition 3 is 4, so the number of blocks to be selected and erased simultaneously is 4; blocks 28 to 31 are then erased.
- <4> After blocks 28 to 31 are erased, the next start block number is 32 and the number of blocks to be erased is 42; the values that satisfy Condition 1 are therefore 1, 2, 4, 8, 16, and 32.  
Moreover, the values that satisfy Condition 2 are 1, 2, 4, 8, and 32, the value that satisfies Condition 3 is 32, so the number of blocks to be selected and erased simultaneously is 32; blocks 32 to 63 are then erased.
- <5> After blocks 32 to 63 are erased, the next start block number is 64, and the number of blocks to be erased is 10; the values that satisfy Condition 1 are therefore 1, 2, 4, and 8.  
Moreover, the values that satisfy Condition 2 are 1, 2, 4, and 8, the value that satisfies Condition 3 is 8, so the number of blocks to be selected and erased simultaneously is 8; blocks 64 to 71 are then erased.
- <6> After blocks 64 to 71 are erased, the next start block number is 72, and the number of blocks to be erased is 2; the values that satisfy Condition 1 are therefore 1 and 2.  
Moreover, the values that satisfy Condition 2 are 1 and 2, the value that satisfies Condition 3 is 2, so the number of blocks to be selected and erased simultaneously is 2; blocks 72 and 73 are then erased.

Therefore, simultaneous selection and erasure is executed six times (25, 26 and 27, 28 to 31, 32 to 63, 64 to 71, and 72 and 73) to erase blocks 25 to 73, so  $M = 6$  is obtained.

Block configuration when executing simultaneous selection and erasure (when erasing blocks 25 to 73)

<Block number>

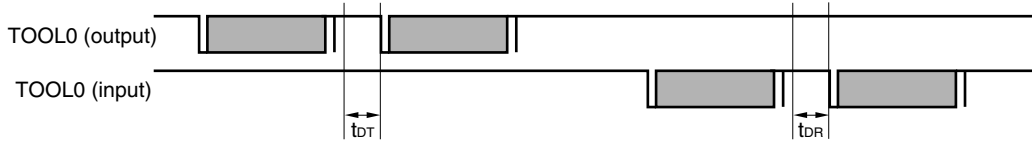


<Range of blocks that can be selected and erased simultaneously>

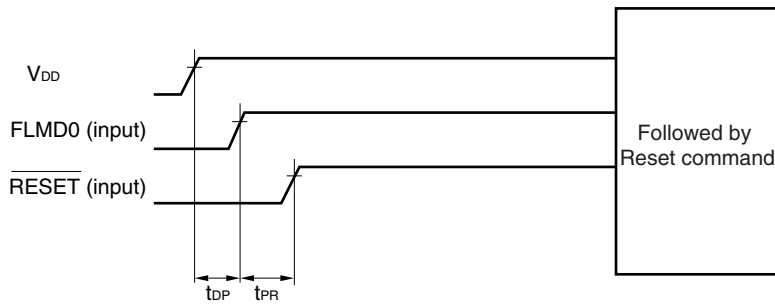
### 5.5 UART Communication Mode

In the figure below, TOOL0 is illustrated as two separate lines for the sake of description, but it is actually a single line. The  $V_{DD}$  level of TOOL0 can be achieved by using a pull-up resistor (the pin is Hi-Z).

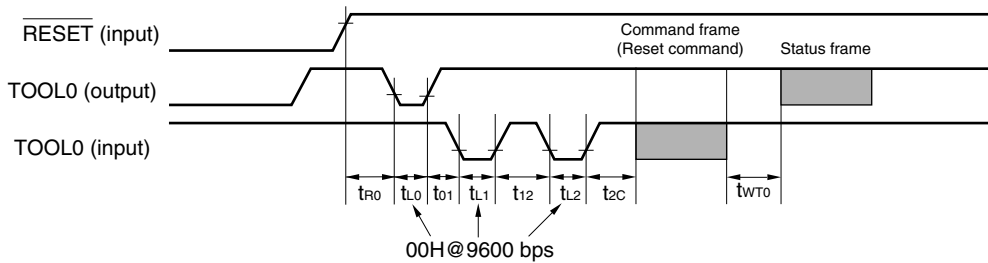
**(a) Data frame**



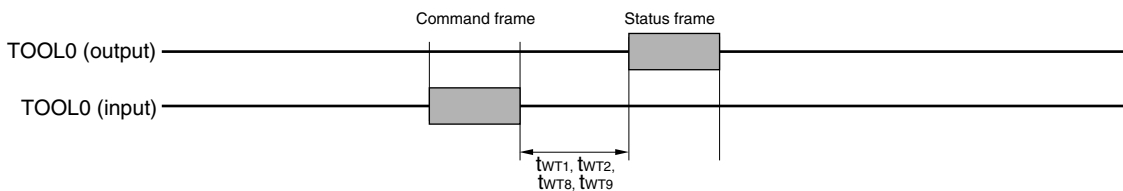
**(b) Programming mode setting**



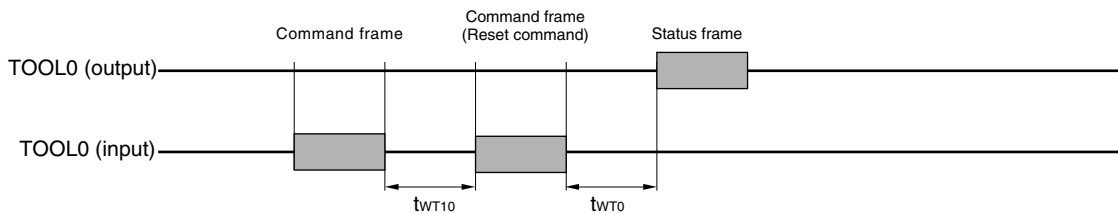
**(c) Reset command**



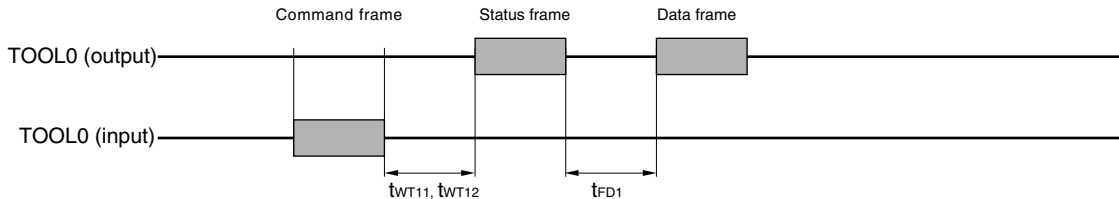
**(d) Chip Erase command/Block Erase command/Block Blank Check command/Oscillating Frequency Set command**



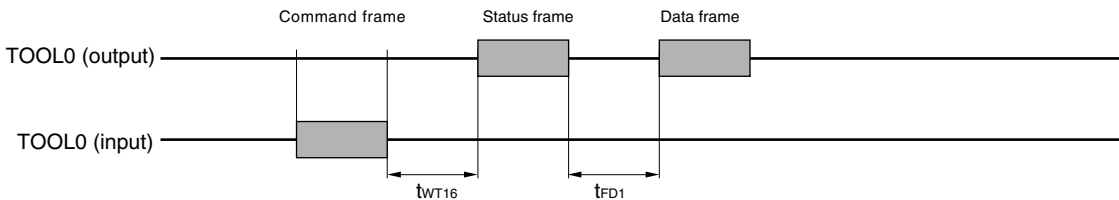
**(e) Baud Rate Set command**



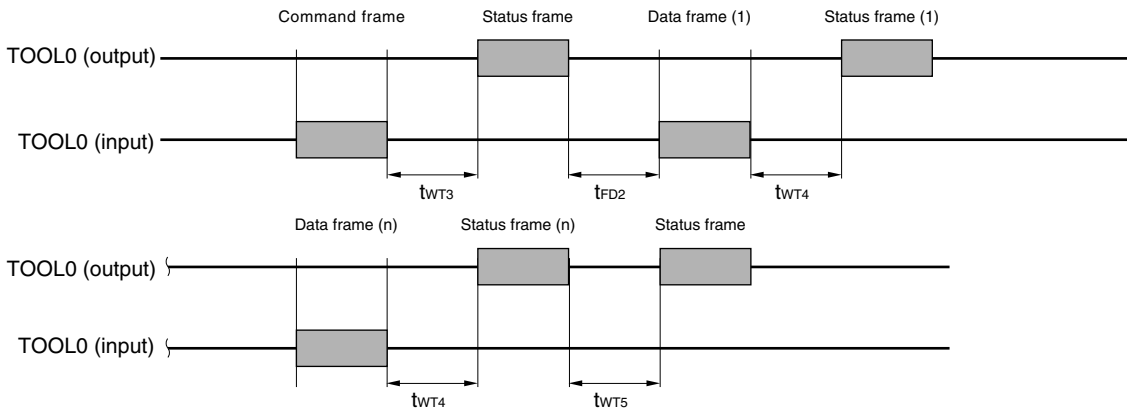
**(f) Silicon Signature command/Version Get command**



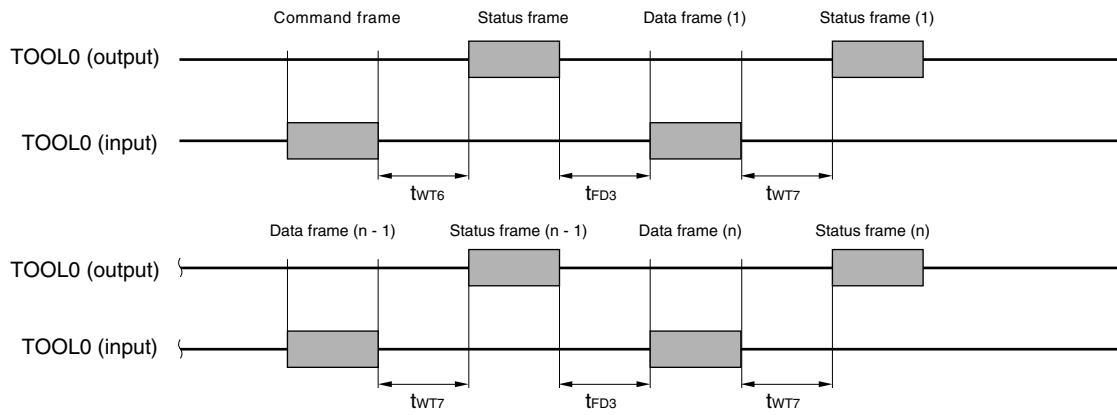
**(g) Checksum command**



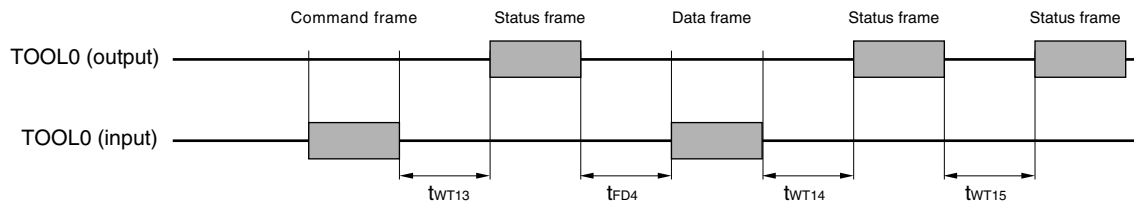
**(h) Programming command**



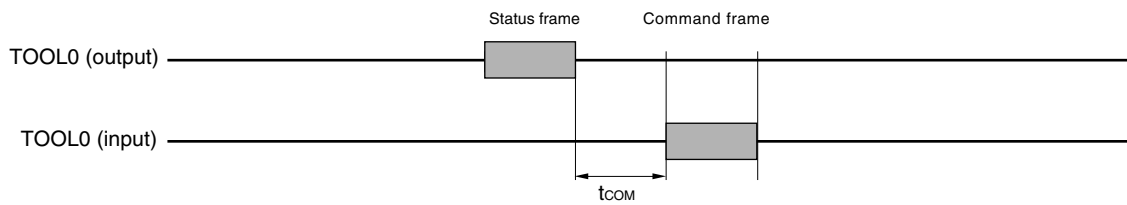
**(i) Verify command**



**(j) Security Set command**



**(k) Wait before command frame transmission**



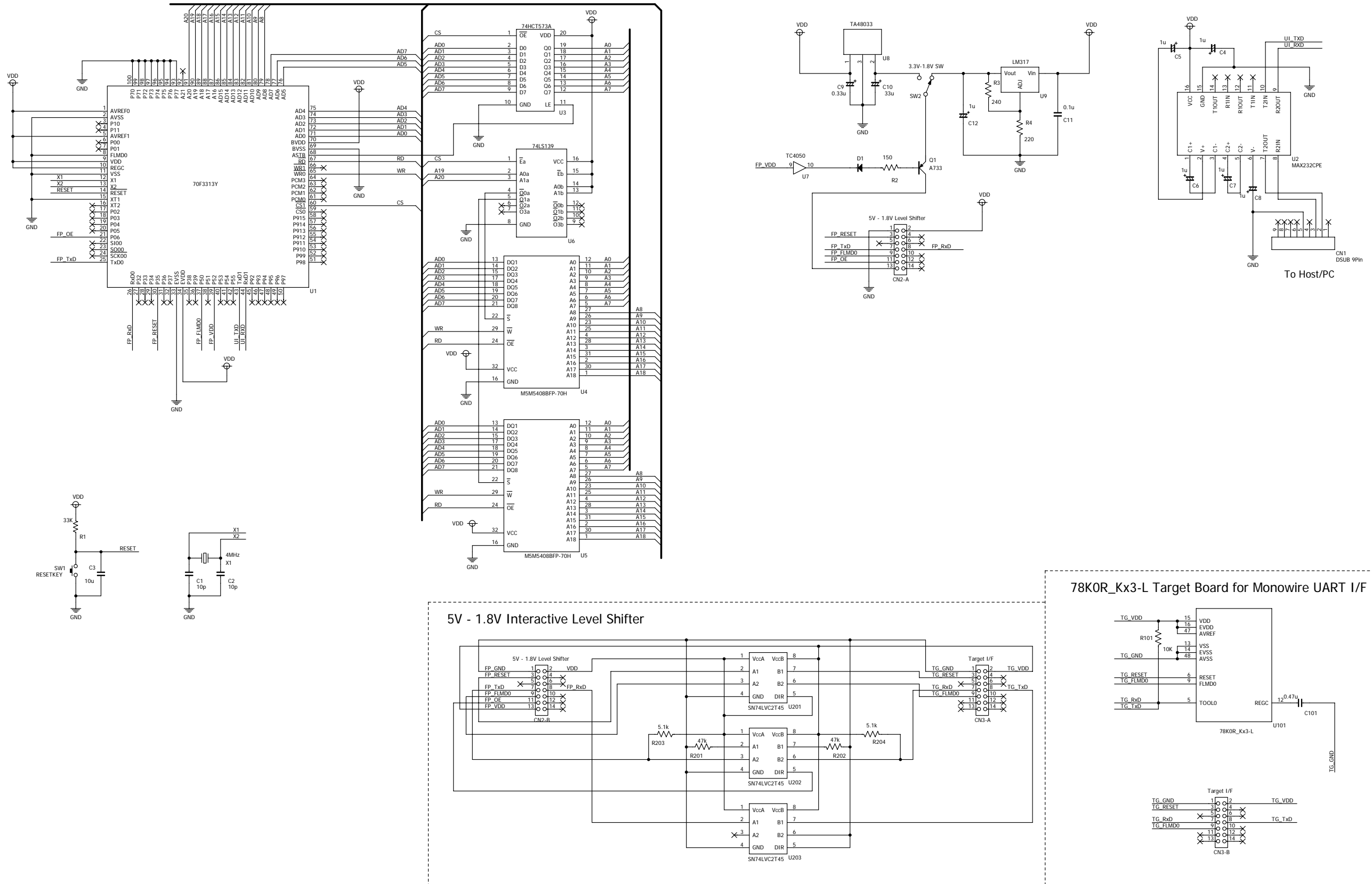
## APPENDIX A CIRCUIT DIAGRAMS (REFERENCE)

Figure A-1 shows a circuit diagram of the programmer and the 78K0R/Kx3-L, for reference.

Figure A-1. Reference Circuit Diagram of Programmer and 78K0R/Kx3-L (Main Board)

78K0R\_Kx3-L Flash Programmer Sample Application Main Board for Monowire UART I/F  
(Wide-voltage mode)

(VDD = 5.0V)



**Remark** For details about connecting unused pins shown in this circuit diagram, see the user's manual of each product.



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