Old Company Name in Catalogs and Other Documents

On April 1st, 2010, NEC Electronics Corporation merged with Renesas Technology Corporation, and Renesas Electronics Corporation took over all the business of both companies. Therefore, although the old company name remains in this document, it is a valid Renesas Electronics document. We appreciate your understanding.

Renesas Electronics website: http://www.renesas.com

April 1st, 2010 Renesas Electronics Corporation

Issued by: Renesas Electronics Corporation (http://www.renesas.com)

Send any inquiries to http://www.renesas.com/inquiry.



Notice

- 1. All information included in this document is current as of the date this document is issued. Such information, however, is subject to change without any prior notice. Before purchasing or using any Renesas Electronics products listed herein, please confirm the latest product information with a Renesas Electronics sales office. Also, please pay regular and careful attention to additional and different information to be disclosed by Renesas Electronics such as that disclosed through our website.
- Renesas Electronics does not assume any liability for infringement of patents, copyrights, or other intellectual property rights
 of third parties by or arising from the use of Renesas Electronics products or technical information described in this document.
 No license, express, implied or otherwise, is granted hereby under any patents, copyrights or other intellectual property rights
 of Renesas Electronics or others.
- 3. You should not alter, modify, copy, or otherwise misappropriate any Renesas Electronics product, whether in whole or in part.
- 4. Descriptions of circuits, software and other related information in this document are provided only to illustrate the operation of semiconductor products and application examples. You are fully responsible for the incorporation of these circuits, software, and information in the design of your equipment. Renesas Electronics assumes no responsibility for any losses incurred by you or third parties arising from the use of these circuits, software, or information.
- 5. When exporting the products or technology described in this document, you should comply with the applicable export control laws and regulations and follow the procedures required by such laws and regulations. You should not use Renesas Electronics products or the technology described in this document for any purpose relating to military applications or use by the military, including but not limited to the development of weapons of mass destruction. Renesas Electronics products and technology may not be used for or incorporated into any products or systems whose manufacture, use, or sale is prohibited under any applicable domestic or foreign laws or regulations.
- 6. Renesas Electronics has used reasonable care in preparing the information included in this document, but Renesas Electronics does not warrant that such information is error free. Renesas Electronics assumes no liability whatsoever for any damages incurred by you resulting from errors in or omissions from the information included herein.
- 7. Renesas Electronics products are classified according to the following three quality grades: "Standard", "High Quality", and "Specific". The recommended applications for each Renesas Electronics product depends on the product's quality grade, as indicated below. You must check the quality grade of each Renesas Electronics product before using it in a particular application. You may not use any Renesas Electronics product for any application categorized as "Specific" without the prior written consent of Renesas Electronics. Further, you may not use any Renesas Electronics product for any application for which it is not intended without the prior written consent of Renesas Electronics. Renesas Electronics shall not be in any way liable for any damages or losses incurred by you or third parties arising from the use of any Renesas Electronics product for an application categorized as "Specific" or for which the product is not intended where you have failed to obtain the prior written consent of Renesas Electronics. The quality grade of each Renesas Electronics product is "Standard" unless otherwise expressly specified in a Renesas Electronics data sheets or data books, etc.
 - "Standard": Computers; office equipment; communications equipment; test and measurement equipment; audio and visual equipment; home electronic appliances; machine tools; personal electronic equipment; and industrial robots.
 - "High Quality": Transportation equipment (automobiles, trains, ships, etc.); traffic control systems; anti-disaster systems; anti-crime systems; safety equipment; and medical equipment not specifically designed for life support.
 - "Specific": Aircraft; aerospace equipment; submersible repeaters; nuclear reactor control systems; medical equipment or systems for life support (e.g. artificial life support devices or systems), surgical implantations, or healthcare intervention (e.g. excision, etc.), and any other applications or purposes that pose a direct threat to human life.
- 8. You should use the Renesas Electronics products described in this document within the range specified by Renesas Electronics, especially with respect to the maximum rating, operating supply voltage range, movement power voltage range, heat radiation characteristics, installation and other product characteristics. Renesas Electronics shall have no liability for malfunctions or damages arising out of the use of Renesas Electronics products beyond such specified ranges.
- 9. Although Renesas Electronics endeavors to improve the quality and reliability of its products, semiconductor products have specific characteristics such as the occurrence of failure at a certain rate and malfunctions under certain use conditions. Further, Renesas Electronics products are not subject to radiation resistance design. Please be sure to implement safety measures to guard them against the possibility of physical injury, and injury or damage caused by fire in the event of the failure of a Renesas Electronics product, such as safety design for hardware and software including but not limited to redundancy, fire control and malfunction prevention, appropriate treatment for aging degradation or any other appropriate measures. Because the evaluation of microcomputer software alone is very difficult, please evaluate the safety of the final products or system manufactured by you.
- 10. Please contact a Renesas Electronics sales office for details as to environmental matters such as the environmental compatibility of each Renesas Electronics product. Please use Renesas Electronics products in compliance with all applicable laws and regulations that regulate the inclusion or use of controlled substances, including without limitation, the EU RoHS Directive. Renesas Electronics assumes no liability for damages or losses occurring as a result of your noncompliance with applicable laws and regulations.
- 11. This document may not be reproduced or duplicated, in any form, in whole or in part, without prior written consent of Renesas Electronics
- 12. Please contact a Renesas Electronics sales office if you have any questions regarding the information contained in this document or Renesas Electronics products, or if you have any other inquiries.
- (Note 1) "Renesas Electronics" as used in this document means Renesas Electronics Corporation and also includes its majority-owned subsidiaries.
- (Note 2) "Renesas Electronics product(s)" means any product developed or manufactured by or for Renesas Electronics.



Application Note

78K0/Kx2-L, 78K0/Ix2

8-bit Single-Chip Microcontrollers

Flash Memory Programming (Programmer)

78K0/KY2-L: μ PD78F0550, 78F0551, 78F0552, 78F0555, 78F0556, 78F0557 78K0/KA2-L: μ PD78F0560, 78F0561, 78F0562, 78F0565, 78F0566, 78F0567 78K0/KB2-L: μ PD78F0571, 78F0572, 78F0573, 78F0576, 78F0577, 78F0578 78K0/KC2-L: μ PD78F0581, 78F0582, 78F0583, 78F0586, 78F0587, 78F0588

78K0/IY2: μPD78F0740, 78F0741, 78F0742, 78F0750, 78F0751, 78F0752

78K0/IA2: μ PD78F0743, 78F0744, 78F0753, 78F0754 78K0/IB2: μ PD78F0745, 78F0746, 78F0755, 78F0756

Document No. U19735EJ1V0AN00 (1st edition) Date Published August 2009 N

[MEMO]

NOTES FOR CMOS DEVICES —

(1) VOLTAGE APPLICATION WAVEFORM AT INPUT PIN

Waveform distortion due to input noise or a reflected wave may cause malfunction. If the input of the CMOS device stays in the area between $V_{\rm IL}$ (MAX) and $V_{\rm IH}$ (MIN) due to noise, etc., the device may malfunction. Take care to prevent chattering noise from entering the device when the input level is fixed, and also in the transition period when the input level passes through the area between $V_{\rm IL}$ (MAX) and $V_{\rm IH}$ (MIN).

(2) HANDLING OF UNUSED INPUT PINS

Unconnected CMOS device inputs can be cause of malfunction. If an input pin is unconnected, it is possible that an internal input level may be generated due to noise, etc., causing malfunction. CMOS devices behave differently than Bipolar or NMOS devices. Input levels of CMOS devices must be fixed high or low by using pull-up or pull-down circuitry. Each unused pin should be connected to VDD or GND via a resistor if there is a possibility that it will be an output pin. All handling related to unused pins must be judged separately for each device and according to related specifications governing the device.

③ PRECAUTION AGAINST ESD

A strong electric field, when exposed to a MOS device, can cause destruction of the gate oxide and ultimately degrade the device operation. Steps must be taken to stop generation of static electricity as much as possible, and quickly dissipate it when it has occurred. Environmental control must be adequate. When it is dry, a humidifier should be used. It is recommended to avoid using insulators that easily build up static electricity. Semiconductor devices must be stored and transported in an anti-static container, static shielding bag or conductive material. All test and measurement tools including work benches and floors should be grounded. The operator should be grounded using a wrist strap. Semiconductor devices must not be touched with bare hands. Similar precautions need to be taken for PW boards with mounted semiconductor devices.

(4) STATUS BEFORE INITIALIZATION

Power-on does not necessarily define the initial status of a MOS device. Immediately after the power source is turned ON, devices with reset functions have not yet been initialized. Hence, power-on does not guarantee output pin levels, I/O settings or contents of registers. A device is not initialized until the reset signal is received. A reset operation must be executed immediately after power-on for devices with reset functions.

5 POWER ON/OFF SEQUENCE

In the case of a device that uses different power supplies for the internal operation and external interface, as a rule, switch on the external power supply after switching on the internal power supply. When switching the power supply off, as a rule, switch off the external power supply and then the internal power supply. Use of the reverse power on/off sequences may result in the application of an overvoltage to the internal elements of the device, causing malfunction and degradation of internal elements due to the passage of an abnormal current.

The correct power on/off sequence must be judged separately for each device and according to related specifications governing the device.

6 INPUT OF SIGNAL DURING POWER OFF STATE

Do not input signals or an I/O pull-up power supply while the device is not powered. The current injection that results from input of such a signal or I/O pull-up power supply may cause malfunction and the abnormal current that passes in the device at this time may cause degradation of internal elements. Input of signals during the power off state must be judged separately for each device and according to related specifications governing the device.

- The information in this document is current as of May, 2009. The information is subject to change
 without notice. For actual design-in, refer to the latest publications of NEC Electronics data sheets,
 etc., for the most up-to-date specifications of NEC Electronics products. Not all products and/or
 types are available in every country. Please check with an NEC Electronics sales representative for
 availability and additional information.
- No part of this document may be copied or reproduced in any form or by any means without the prior
 written consent of NEC Electronics. NEC Electronics assumes no responsibility for any errors that may
 appear in this document.
- NEC Electronics does not assume any liability for infringement of patents, copyrights or other intellectual property rights of third parties by or arising from the use of NEC Electronics products listed in this document or any other liability arising from the use of such products. No license, express, implied or otherwise, is granted under any patents, copyrights or other intellectual property rights of NEC Electronics or others.
- Descriptions of circuits, software and other related information in this document are provided for illustrative
 purposes in semiconductor product operation and application examples. The incorporation of these
 circuits, software and information in the design of a customer's equipment shall be done under the full
 responsibility of the customer. NEC Electronics assumes no responsibility for any losses incurred by
 customers or third parties arising from the use of these circuits, software and information.
- While NEC Electronics endeavors to enhance the quality and safety of NEC Electronics products, customers agree and acknowledge that the possibility of defects thereof cannot be eliminated entirely. In addition, NEC Electronics products are not taken measures to prevent radioactive rays in the product design. When customers use NEC Electronics products with their products, customers shall, on their own responsibility, incorporate sufficient safety measures such as redundancy, fire-containment and anti-failure features to their products in order to avoid risks of the damages to property (including public or social property) or injury (including death) to persons, as the result of defects of NEC Electronics products.
- NEC Electronics products are classified into the following three quality grades: "Standard", "Special" and "Specific".
 - The "Specific" quality grade applies only to NEC Electronics products developed based on a customerdesignated "quality assurance program" for a specific application. The recommended applications of an NEC Electronics product depend on its quality grade, as indicated below. Customers must check the quality grade of each NEC Electronics product before using it in a particular application.
 - "Standard": Computers, office equipment, communications equipment, test and measurement equipment, audio and visual equipment, home electronic appliances, machine tools, personal electronic equipment and industrial robots.
 - "Special": Transportation equipment (automobiles, trains, ships, etc.), traffic control systems, anti-disaster systems, anti-crime systems, safety equipment and medical equipment (not specifically designed for life support).
 - "Specific": Aircraft, aerospace equipment, submersible repeaters, nuclear reactor control systems, life support systems and medical equipment for life support, etc.

The quality grade of NEC Electronics products is "Standard" unless otherwise expressly specified in NEC Electronics data sheets or data books, etc. If customers wish to use NEC Electronics products in applications not intended by NEC Electronics, they must contact an NEC Electronics sales representative in advance to determine NEC Electronics' willingness to support a given application.

(Note)

- (1) "NEC Electronics" as used in this statement means NEC Electronics Corporation and also includes its majority-owned subsidiaries.
- (2) "NEC Electronics products" means any product developed or manufactured by or for NEC Electronics (as defined above).

M8E0904E

INTRODUCTION

Caution

Explanations in this application note assume use of the 78K0/Kx2-L as the representative microcontroller. Users of a product other than the 78K0/Kx2-L should read 78K0/Kx2-L as referring to that product.

Target Readers

This application note is intended for users who understand the functions of the 78K0/Kx2-L and 78K0/Ix2 and who will use these products to design application systems.

Purpose

The purpose of this application note is to help users understand how to develop dedicated flash memory programmers for rewriting the internal flash memory of the 78K0/Kx2-L and 78K0/Ix2.

The sample programs and circuit diagrams shown in this document are for reference only and are not intended for use in actual design-ins.

Therefore, these sample programs must be used at the user's own risk. Correct operation is not guaranteed if these sample programs are used.

Organization

This manual consists of the following main sections.

- · Flash memory programming
- Command/data frame format
- Description of command processing
- · UART communication mode
- Flash memory programming parameter characteristics

How to Read This Manual

It is assumed that the reader of this manual has general knowledge in the fields of electrical engineering, logic circuits, and microcontrollers.

- To gain a general understanding of functions:
 - → Read this manual in the order of the **CONTENTS**.
- To learn more about hardware functions of the 78K0/Kx2-L and 78K0/Ix2:
 - → See the user's manual of the 78K0/Kx2-L and 78K0/Ix2.

Conventions

Data significance: Higher digits on the left and lower digits on the right

Active low representation: \overline{xxx} (overscore over pin or signal name)

Note: Footnote for item marked with **Note** in the text

Caution: Information requiring particular attention

Remark: Supplementary information

Numeral representation: Binary.....xxxx or xxxxB

Decimalxxxx
HexadecimalxxxH

Related Documents

The related documents indicated in this publication may include preliminary versions. However, preliminary versions are not marked as such.

Documents related to devices

Document Name	Document No.
78K0/Kx2-L User's Manual	U19111E
78K0/lx2 User's Manual	U19353E
78K0 Microcontrollers Instructions User's Manual	U12326E

Caution The related documents listed above are subject to change without notice. Be sure to use the latest version of each document for designing.

CONTENTS

CHAPT	ER 1 FLASH MEMORY PROGRAMMING	11
1.1	Overview	11
1.2	System Configuration	
1.3	Flash Memory Configuration	
1.4	Command List and Status List	
	1.4.1 Command list	15
	1.4.2 Status list	
1.5	Power Application and Setting Flash Memory Programming Mode	
	1.5.1 UART communication mode	
	1.5.2 Mode setting flowchart	19
	1.5.3 Sample program	
1.6	Shutting Down Target Power Supply	
1.7	Command Execution Flow at Flash Memory Rewriting	
	·	
CHAPT	ER 2 COMMAND/DATA FRAME FORMAT	27
2.1	Command Frame Transmission Processing	20
2.2	Data Frame Transmission Processing	
2.3	Data Frame Reception Processing	
2.0	Data i fame neception i rocessing	
CHAPT	ER 3 DESCRIPTION OF COMMAND PROCESSING	30
3.1	Reset Command	30
0.1	3.1.1 Description	
	3.1.2 Command frame and status frame	
3.2	Baud Rate Set Command	
	3.2.1 Description	
	3.2.2 Command frame and status frame	
3.3	Chip Erase Command	
	3.3.1 Description	
	3.3.2 Command frame and status frame	
3.4	Block Erase Command	
	3.4.1 Description	
	3.4.2 Command frame and status frame	33
3.5	Programming Command	34
	3.5.1 Description	32
	3.5.2 Command frame and status frame	32
	3.5.3 Data frame and status frame	32
	3.5.4 Completion of transferring all data and status frame	35
3.6	Verify Command	36
	3.6.1 Description	
	3.6.2 Command frame and status frame	
	3.6.3 Data frame and status frame	36
3.7	Block Blank Check Command	38
	3.7.1 Description	38

	3.7.2	Command frame and status frame	38
3.8	Silico	n Signature Command	39
	3.8.1	Description	39
	3.8.2	Command frame and status frame	39
	3.8.3	Silicon signature data frame	40
	3.8.4	Silicon signature list	42
3.9	Versio	on Get Command	47
	3.9.1	Description	47
	3.9.2	Command frame and status frame	47
	3.9.3	Version data frame	48
3.10	Check	ssum Command	49
	3.10.1	Description	49
	3.10.2	Command frame and status frame	49
	3.10.3	Checksum data frame	49
3.11	Secur	ity Set Command	50
	3.11.1	Description	50
	3.11.2	Command frame and status frame	50
	3.11.3	Data frame and status frame	51
	3.11.4	Internal verify check and status frame	51
CHAPTE	ER4U	UART COMMUNICATION MODE	53
4.1	Comm	nand Frame Transmission Processing Flowchart	53
4.2		Frame Transmission Processing Flowchart	
4.3		Frame Reception Processing Flowchart	
4.4		Command	
	4.4.1	Processing sequence chart	56
	4.4.2	Description of processing sequence	57
	4.4.3	Status at processing completion	57
	4.4.4	Flowchart	58
	4.4.5	Sample program	
4.5	Baud	Rate Set Command	
	4.5.1	Processing sequence chart	60
	4.5.2	Description of processing sequence	61
	4.5.3	Status at processing completion	61
	4.5.4	Flowchart	62
	4.5.5	Sample program	63
4.6	Chip E	Erase Command	64
	4.6.1	Processing sequence chart	64
	4.6.2	Description of processing sequence	65
	4.6.3	Status at processing completion	65
	4.6.4	Flowchart	66
	4.6.5	Sample program	67
4.7	Block	Erase Command	68
	4.7.1	Processing sequence chart	68
	4.7.2	Description of processing sequence	69
	4.7.3	Status at processing completion	69
	4.7.4	Flowchart	70
	4.7.5	Sample program	71

4.8	Progra	amming Command	72
	4.8.1	Processing sequence chart	72
	4.8.2	Description of processing sequence	73
	4.8.3	Status at processing completion	74
	4.8.4	Flowchart	75
	4.8.5	Sample program	76
4.9	Verify	Command	78
	4.9.1	Processing sequence chart	78
	4.9.2	Description of processing sequence	79
	4.9.3	Status at processing completion	79
	4.9.4	Flowchart	80
	4.9.5	Sample program	81
4.10	Block	Blank Check Command	83
	4.10.1	Processing sequence chart	83
	4.10.2	Description of processing sequence	84
	4.10.3	Status at processing completion	
	4.10.4	Flowchart	85
	4.10.5	Sample program	86
4.11	Silico	n Signature Command	
	4.11.1	Processing sequence chart	
	4.11.2	Description of processing sequence	88
	4.11.3	Status at processing completion	
	4.11.4	Flowchart	
	4.11.5	Sample program	90
4.12	Versio	on Get Command	
	4.12.1	Processing sequence chart	91
	4.12.2	Description of processing sequence	
	4.12.3	Status at processing completion	
	4.12.4	Flowchart	
	4.12.5	Sample program	
4.13	Check	sum Command	
	4.13.1	Processing sequence chart	95
	4.13.2	Description of processing sequence	
	4.13.3	Status at processing completion	
	4.13.4	Flowchart	
	4.13.5	Sample program	
4.14		ity Set Command	
	4.14.1	Processing sequence chart	
	4.14.2	Description of processing sequence	
	4.14.3	Status at processing completion	
	4.14.4	Flowchart	
	4.14.5	Sample program	
L A DT	=D & E	FLASH MEMORY PROGRAMMING PARAMETER CHARACTERISTICS	10/
ларіі 5.1		Memory Parameter Characteristics of 78K0/Kx2-L	
J. I	5.1.1	Flash memory programming mode setting time	
	5.1.1	Programming characteristics	
	5.1.2	Command characteristics	
	J. I.J	COMMINATIO CHATACLETISTICS	106

5.2	Flash	Memory Parameter Characteristics of 78K0/lx2	107
	5.2.1	Flash memory programming mode setting time	107
	5.2.2	Programming characteristics	108
	5.2.3	Command characteristics	109
5.3	Simu	taneous Selection and Erasure Performed by Block Erase Command	110
	5.3.1	Calculation of number of blocks to be selected and erased simultaneously	110
	5.3.2	Calculation of execution count (M) of simultaneous selection and erasure	111
5.4	UART	Communication Mode	118
APPEN	DIX A	CIRCUIT DIAGRAMS (REFERENCE)	121

CHAPTER 1 FLASH MEMORY PROGRAMMING

To rewrite the contents of the internal flash memory of the 78K0/Kx2-L, a dedicated flash memory programmer (hereafter referred to as the "programmer") is usually used.

This Application Note explains how to develop a dedicated programmer.

1.1 Overview

The 78K0/Kx2-L incorporates firmware that controls flash memory programming. The programming to the internal flash memory is performed by transmitting/receiving commands between the programmer and the 78K0/Kx2-L via serial communication.

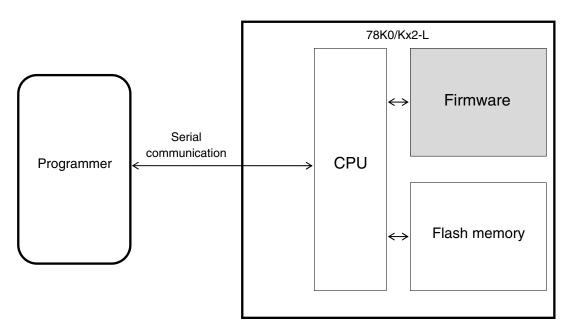


Figure 1-1. System Outline of Flash Memory Programming in 78K0/Kx2-L

1.2 System Configuration

Examples of the system configuration for programming the flash memory are illustrated in Figure 1-2.

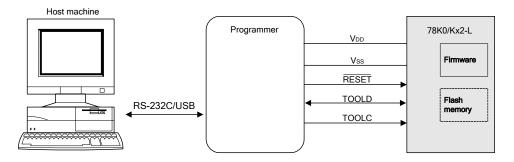
Figure 1-2 illustrates how to program the flash memory with the programmer, under control of a host machine.

Depending on how the programmer is connected, the programmer can be used in a standalone mode without using the host machine, if a user program has been downloaded to the programmer in advance.

For example, NEC Electronics' flash memory programmer PG-FP5 can execute programming either by using the GUI software with a host machine connected or by itself (standalone).

Figure 1-2. System Configuration Example

Single-wire UART communication mode (LSB-first transfer)



Remarks 1. The 78K0/Kx2-L can only communicate via the single-wire UART communication mode.

2. For the pins used by flash memory programming and the recommended connections of unused pins, see the user's manual of each product.

1.3 Flash Memory Configuration

The 78K0/Kx2-L and 78K0/Ix2 must manage product-specific information (such as device name and memory information) via the programmer.

Table 1-1 shows the flash memory size of the 78K0/Kx2-L and 78K0/Ix2 and Figure 1-3 shows the configuration of the flash memory.

Table 1-1. Size of Flash Memory for Each Product

(a) Size of flash memory for 78K0/Kx2-L

Device Name	Flash Memory Size
μPD78F0550, 78F0555, 78F0560, 78F0565	4 KB
μPD78F0551, 78F0556, 78F0561, 78F0566, 78F0571, 78F0576, 78F0581, 78F0586	8 KB
μPD78F0552, 78F0557, 78F0562, 78F0567, 78F0572, 78F0577, 78F0582, 78F0587	16 KB
μPD78F0573, 78F0578, 78F0583, 78F0588	32 KB

(b) Size of flash memory for 78K0/lx2

Device Name	Flash Memory Size
μPD78F0740, 78F0750	4 KB
μPD78F0741, 78F0743, 78F0745, 78F0751, 78F0753, 78F0755	8 KB
μPD78F0742, 78F0744, 78F0746, 78F0752, 78F0754, 78F0756	16 KB

Remark Products under development are included in the above tables.

<Address> <Flash memory size> <Block number> 07FFFH Block 1FH 32 KB 1 KB Block 10H 04000H 03FFFH 1 KB Block 0FH 16 KB Block 08H 1 KB 02000H 01FFFH 1 KB Block 07H 8 KB 1 KB Block 04H 01000H 00FFFH 1 KB Block 03H 4 KB 1 KB Block 00H 00000H

Figure 1-3. Flash Memory Configuration

Remark Each block consists of 1 KB (this figure only illustrates some blocks in the flash memory).

1.4 Command List and Status List

The flash memory incorporated in the 78K0/Kx2-L can be rewritten by using the commands listed in Table 1-2. The programmer transmits commands to control these functions to the 78K0/Kx2-L, and checks the response status sent from the 78K0/Kx2-L, to manipulate the flash memory.

1.4.1 Command list

The commands used by the programmer and their functions are listed below.

Table 1-2. List of Commands Transmitted from Programmer to 78K0/Kx2-L

Command Number	Command Name	Function Name	Function	
20H	Chip Erase	Erase	Erases the entire flash memory area.	
22H	Block Erase		Erases a specified area in the flash memory.	
40H	Programming	Write	Writes data to a specified area in the flash memory.	
13H	Verify	Verify	Compares the contents in a specified area in the flash memory with the data transmitted from the programmer.	
32H	Block Blank Check	Blank check	Checks the erase status of a specified block in the flash memory.	
C0H	Silicon Signature	Information acquisition	Acquires 78K0/Kx2-L information (product name, flash memory configuration, etc.).	
C5H	Version Get		Acquires version of the 78K0/Kx2-L and firmware.	
ВОН	Checksum		Acquires checksum data of a specified area.	
A0H	Security Set	Security	Sets security information.	
00H	Reset	Others	Detects synchronization in communication.	
9AH	Baud Rate Set		Sets the baud rate when UART communication mode is selected.	

1.4.2 Status list

The following table lists the status codes the programmer receives from the 78K0/Kx2-L.

Table 1-3. Status Code List

Status Code	Status	Description		
04H	Command number error	Error returned if a command not supported is received		
05H	Parameter error	Error returned if command information (parameter) is invalid		
06H	Normal acknowledgment (ACK)	Normal acknowledgment		
07H	Checksum error	Error returned if data in a frame transmitted from the programmer is abnormal		
0FH	Verify error	Error returned if a verify error has occurred upon verifying data transmitted from the programmer		
10H	Protect error	Error returned if an attempt is made to execute processing that is prohibited by the Security Set command		
15H	Negative acknowledgment (NACK)	Negative acknowledgment		
1AH	MRG10 error	Erase verify error		
1BH	MRG11 error	Internal verify error or blank check error during data write		
20H	Read error	Error returned if reading security information fails		
1CH	Write error	Write error		

Reception of a checksum error or NACK is treated as an immediate abnormal end in this manual. When a dedicated programmer is developed, however, the processing may be retried without problem from the wait immediately before transmission of the command that results a checksum error or NACK. In this event, limiting the retry count is recommended for preventing infinite repetition of the retry operation.

Although not listed in the above table, if a time-out error (BUSY time-out or time-out in data frame reception during UART communication) occurs, it is recommended to shutdown the power supply to the 78K0/Kx2-L (refer to 1.6 Shutting Down Target Power Supply) and then connect the power supply again.

1.5 Power Application and Setting Flash Memory Programming Mode

To rewrite the contents of the flash memory with the programmer, the 78K0/Kx2-L must first be set to the flash memory programming mode.

The RESET, TOOLD, and TOOLC pins are set to low level, and then the TOOLD and TOOLC pins are set to high level while keeping the RESET pin at low level. Next, the reset is ended by alternately inputting two pulses and seven pulses three times from the TOOLC and TOOLD pins, respectively, enabling the system to enter the flash memory programming mode.

The following figure illustrates a timing chart for setting the flash memory programming mode and selecting the communication mode.

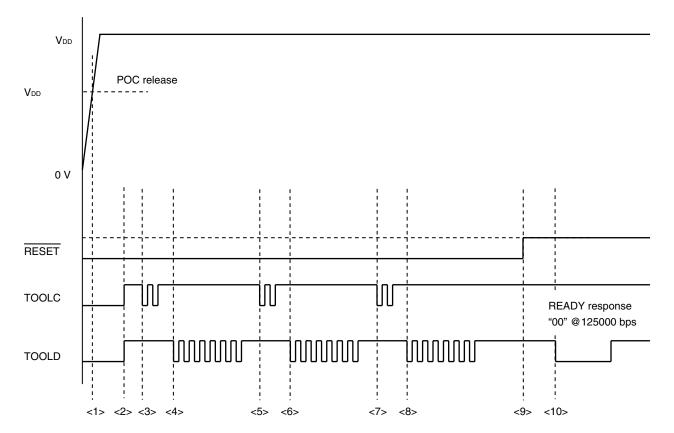


Figure 1-4. Setting Flash Memory Programming Mode and Selecting Communication Mode

- <1> Power application (VDD)
- <2> TOOLC/TOOLD = Low-level output → High-level output
- <3> TOOLC pulse output (2 pulses)
- <4> TOOLD pulse output (7 pulses)
- <5> TOOLC pulse output (2 pulses)
- <6> TOOLD pulse output (7 pulses)
- <7> TOOLC pulse output (2 pulses)
- <8> TOOLD pulse output (7 pulses)
- <9> Reset release (serial programming mode setting)
- <10> READY pulse ("00" @125000 bps input start (microcontroller \rightarrow programmer))

The following table shows the communication mode that can be selected with the 78K0/Kx2-L and the port to be used.

Table 1-4. 78K0/Kx2-L Communication Mode and Used Port

Communication Mode	Port Used for Communication
Single-wire UART	TOOLD0, TOOLD1

1.5.1 UART communication mode

The TOOLD pin is used for UART communication. The communication conditions are as shown below.

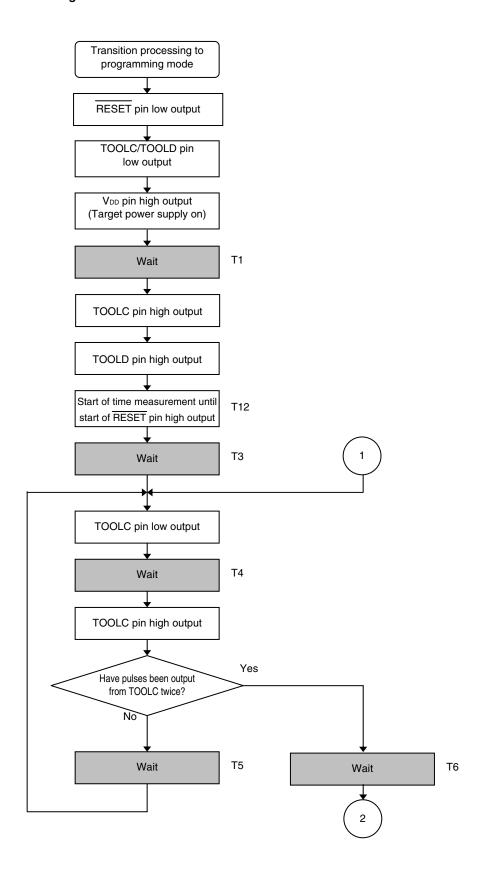
Table 1-6. UART Communication Conditions

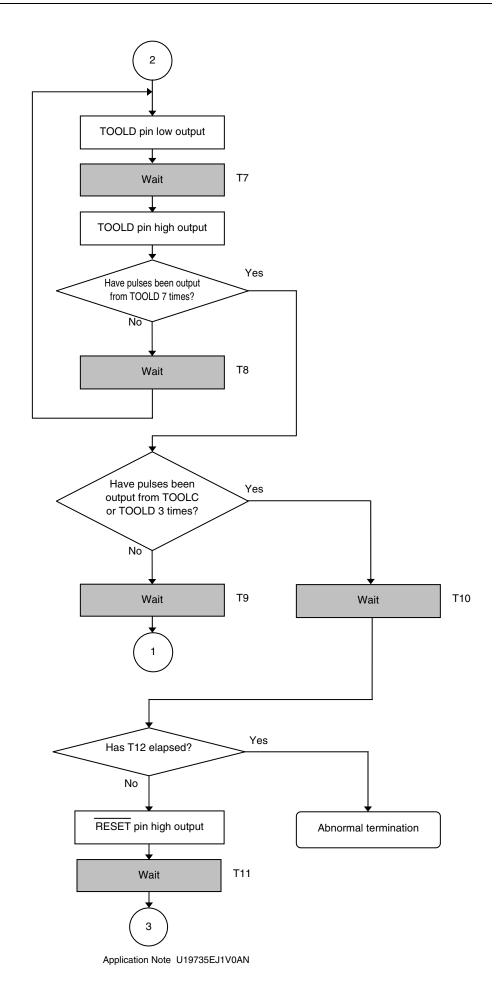
Item	Description	
Baud rate	Communication is performed at 125,000 bps until the Baud Rate Set command for baud ra setting command processing is transmitted. The transmission rate is changed to the baud ra set by the Baud Rate Set command from the transmission of the Reset command for baud ra command processing. For details of the settable baud rate, refer to 3.2 Baud Rate S Command.	
Parity bit	None	
Data length	8 bits (LSB first)	
Stop bit	2 bits (programmer → 78K0/Kx2-L)/1 bit (78K0/Kx2-L → programmer)	

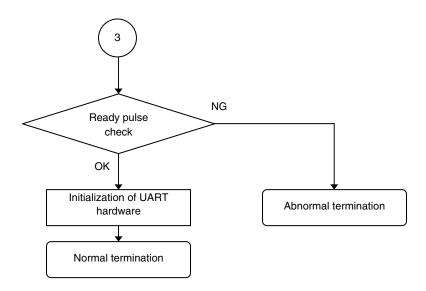
The role of the master and slave is occasionally exchanged during UART communication, so communication at the optimum timing is possible.

Caution Set the same baud rate to the master and slave devices when performing UART communication.

1.5.2 Mode setting flowchart







1.5.3 Sample program

The following shows a sample program for mode setting processing.

```
/*
                                                    * /
                                                    */
/* connect to Flash device
                                                    */
fl_con_dev(void)
extern void init_fl_uart(void);
extern void init_fl_csi(void);
extern void stop_UARTO(void);
 u16 rc = NO_ERROR;
      cnt1,cnt2;
 u8
 SRMK0 = true;  // disable UART Rx INT.
 UARTE0 = false;
                // disable UART H.W.
 stop_UART0();
                // TxD/RxD = Hi-Z
 OE_ENABLE()
                // TxD = output
 pFL_RES
          = low;
                     // RESET = low
 pmFL_FLMD0 = PM_OUT;
                     // FLMD0 = Low output
 pFL_FLMD0 = low;
 pmFL_TxD = PM_OUT;
                     // TxD = Low output
 pFL_TxD
           = low;
 //VDD low => high
                     // VDD = high
 FL_VDD_HI();
 fl_wait(t1);
                      // wait
 //TOOLC(FLMD0),TOOLD low => high
         = hi; // FLMD0 = high
 pFL_FLMD0
           = hi;
                     // TOOLD = high
 pFL_TxD
                     // start "t12" wait timer
 start_flto(t12);
                     // wait
 fl_wait(t3);
 cnt1 = 3;
 while(cnt1)
 {
 cnt1--;
      //TOOLC(FLMD0) output pulse (2 times)
      cnt2 = 2;
      while(cnt2)
           cnt2--;
                              // FLMD0 = low
           pFL_FLMD0 = low;
           fl_wait(t4);
                                 // wait
                            // FLMD0 = high
           pFL_FLMD0 = hi;
           if(cnt2){
                     fl_wait(t5);} // wait
           else{
                     fl_wait(t6);} // wait
      }
```

```
//TOOLD output pulse (7 times)
     cnt2 = 7;
     while(cnt2)
     {
          cnt2--;
                                   // TOOLD = low
          pFL_TxD
                         = low;
          fl_wait(t7);
                                     // wait
                                     // TOOLD = high
          pFL_TxD
                         = hi;
          if(cnt2){     fl_wait(t8);}
                                    // wait
     }
     if(cnt1){     fl_wait(t9);}
                                   // wait
     else{
              fl_wait(t10);}
                                    // wait
}
return FLC_COM_ERR; // yes
}
else{
     stop_fl_timer(); // no
}
//RESET low => high
pFL_RES = hi;
                    // RESET = high
OE_DISABLE()
                    // TxD = Hi-Z
pmFL_TxD = PM_IN;
                    // TxD = Hi-Z
fl_wait(t11_MIN);
                    // wait
rc = check_ready_pulse(); // check "READY PULSE" from target device
if (rc) {
    return rc;
                    // pulse width/timing error
}
UARTE0 = true;
               // enable UART h.w.
SRIF0 = false;
              // clear UART Rx IRQ flag
SRMK0 = false;
              // enable UART Rx INT.
return rc;
// start RESET command proc.
```

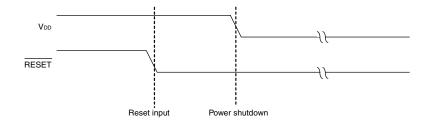
1.6 Shutting Down Target Power Supply

After each command execution is completed, shut down the power supply to the target after setting the $\overline{\text{RESET}}$ pin to low level, as shown below.

Set other pins to Hi-Z when shutting down the power supply to the target.

Caution Shutting down the power supply and inputting a reset during command processing are prohibited.

Figure 1-5. Timing for Terminating Flash Memory Programming Mode



1.7 Command Execution Flow at Flash Memory Rewriting

Figure 1-6 illustrates the basic flowchart when flash memory rewriting is performed with the programmer.

Other than commands shown in Figure 1-6, the Verify command and Checksum command are also supported.

Basic flow Power application to target (See Figure 1-4) Mode setting (reset release) (See 1.5) Selection of communication mode (pulse input) (See 1.5) Synchronization processing (Reset command) (See 3.1) Baud rate setting (see 3.2) Silicon signature acquisition (Silicon Signature command) (See 3.8) Command execution Processing No completed? Yes Reset input and power shutdown during rewriting is prohibited because security information may be lost. Target power shutdown processing (See 1.6) End

Figure 1-6. Basic Flowchart for Flash Memory Rewrite Processing

Remark The example of each command execution is shown in Figure 1-7.

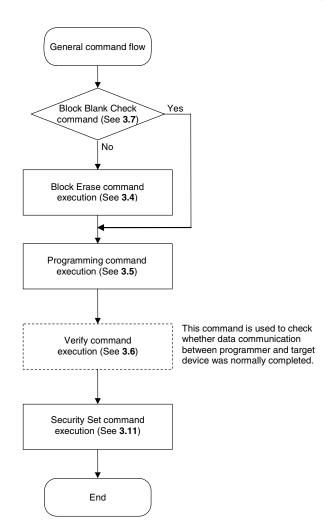


Figure 1-7. General Command Execution Flow at Flash Memory Rewriting

CHAPTER 2 COMMAND/DATA FRAME FORMAT

The programmer uses the command frame to transmit commands to the 78K0/Kx2-L. The 78K0/Kx2-L uses the data frame to transmit write data or verify data to the programmer. A header, footer, data length information, and checksum are appended to each frame to enhance the reliability of the transferred data.

The following shows the format of a command frame and data frame.

Figure 2-1. Command Frame Format

SOH	LEN	COM	Command information (variable length)	SUM	ETX
(1 byte)	(1 byte)	(1 byte)	(Max. 255 bytes)	(1 byte)	(1 byte)

Figure 2-2. Data Frame Format

STX	LEN	Data (variable length)	SUM	ETX or ETB
(1 byte)	(1 byte)	(Max. 256 bytes)	(1 byte)	(1 byte)

Table 2-1. Description of Symbols in Each Frame

Symbol	Value	Description
SOH	01H	Command frame header
STX	02H	Data frame header
LEN	-	Data length information (00H indicates 256) Command frame: COM + command information length Data frame: Data field length
COM	-	Command number
SUM	-	Checksum data for a frame Obtained by sequentially subtracting all of calculation target data from the initial value (00H) in 1-byte units (borrow is ignored). The calculation targets are as follows. Command frame: LEN + COM + all of command information Data frame: LEN + all of data
ETB	17H	Footer of data frame other than the last frame
ETX	03H	Command frame footer, or footer of last data frame

The following shows examples of calculating the checksum (SUM) for a frame.

[Command frame]

No command information is included in the following example of a Chip Erase command frame, so LEN and COM are targets of checksum calculation.

SOH	LEN	СОМ	SUM	ETX
01H	01H	20H	Checksum	03H
	Checksum cale	culation targets		

For this command frame, checksum data is obtained as follows.

00H (initial value) – 01H (LEN) – 20H (COM) = DFH (Borrow ignored. Lower 8 bits only.)

The command frame finally transmitted is as follows.

SOH	LEN	СОМ	SUM	ETX
01H	01H	20H	DFH	03H

[Data frame]

To transmit a data frame as shown below, LEN and D1 to D4 are targets of checksum calculation.

STX	LEN	D1	D2	D3	D4	SUM	ETX
02H	04H	FFH	80H	40H	22H	Checksum	03H

For this data frame, checksum data is obtained as follows.

00H (initial value)
$$-$$
 04H (LEN) $-$ FFH (D1) $-$ 80H (D2) $-$ 40H (D3) $-$ 22H (D4) $=$ 1BH (Borrow ignored. Lower 8 bits only.)

The data frame finally transmitted is as follows.

STX	LEN	D1	D2	D3	D4	SUM	ETX
02H	04H	FFH	80H	40H	22H	1BH	03H

When a data frame is received, the checksum data is calculated in the same manner, and the obtained value is used to detect a checksum error by judging whether the value is the same as that stored in the SUM field of the receive data. When a data frame as shown below is received, for example, a checksum error is detected.

STX	LEN	D1	D2	D3	D4	SUM	ETX
02H	04H	FFH	80H	40H	22H	1AH	03H

↑ Normally 1BH

2.1 Command Frame Transmission Processing

For details of the flowchart of processing to transmit command frames, read **4.1 Command Frame Transmission Processing Flowchart**.

2.2 Data Frame Transmission Processing

The write data frame (user program), verify data frame (user program), and security data frame (security flag) are transmitted as a data frame.

For details of the flowchart of processing to transmit data frames, read **4.2 Data Frame Transmission**Processing Flowchart.

2.3 Data Frame Reception Processing

The status frame, silicon signature data frame, version data frame, and checksum data frame are received as a data frame.

For details of the flowchart of processing to receive data frames, read **4.3 Data Frame Reception Processing Flowchart**.

CHAPTER 3 DESCRIPTION OF COMMAND PROCESSING

3.1 Reset Command

3.1.1 Description

This command is used to check the establishment of communication between the programmer and the 78K0/Kx2-L after the communication mode is set.

The baud rate must be the same for the programmer and 78K0/Kx2-L. Therefore, synchronization is checked by executing the Reset command immediately after the serial programming mode transitions to the operating mode. The default baud rate for the 78K0/Kx2-L is 125,000 bps.

3.1.2 Command frame and status frame

Figure 3-1 shows the format of a command frame for the Reset command, and Figure 3-2 shows the status frame for the command.

Figure 3-1. Reset Command Frame (from Programmer to 78K0/Kx2-L)

SOH	LEN	СОМ	SUM	ETX
01H	01H	00H (Reset)	Checksum	03H

Figure 3-2. Status Frame for Reset Command (from 78K0/Kx2-L to Programmer)

STX	LEN	Data	SUM	ETX
02H	1	ST1	Checksum	03H

Remark ST1: Synchronization detection result

See **4.4 Reset Command** for details about the flowchart of the processing sequence between the programmer and the 78K0/Kx2-L, the flowchart of command processing, and the sample program.

3.2 Baud Rate Set Command

3.2.1 Description

This command is used to change the baud rate for UART communication (125,000 bps by default).

Even if not changing the baud rate, be sure to execute this command after synchronization processing. If the command is not executed, subsequent commands cannot be normally executed.

After the Baud Rate Set command has been executed, the Reset command must be executed to check synchronization at the changed baud rate.

3.2.2 Command frame and status frame

Figure 3-3 shows the format of a command frame for the Baud Rate Set command, and Figure 3-4 shows the status frame for the command.

Figure 3-3. Baud Rate Set Command Frame (from Programmer to 78K0/Kx2-L)

SOH	LEN	СОМ		Command Information ^{Note}					ETX
01H	06H	9AH	D01	D02H	D02L	D03	D04	Checksum	03H

Note For details of the command information setting, refer to **Table 3-1**. If data other than in Table 3-1 is set, a time-out error will occur.

If a time-out error has occurred, execute a hardware reset and re-set the flash memory programming mode.

Remark D01: Invalid data

D02H, D02L: Baud rate setting
D03: Invalid data
D04: Invalid data

Table 3-1. Command Information Setting

Baud Rate	D01	D02H	D02L	D03	D04
125,000 bps	Fixed to FFH	00H	00H	Fixed to FFH	Fixed to FFH
250,000 bps		00H	01H		
500,000 bps		00H	02H		

Figure 3-4. Status Frame for Baud Rate Set Command (from 78K0/Kx2-L to Programmer)

STX	LEN	Data	SUM	ETX
02H	01H	ST1	Checksum	03H

Remark ST1: Synchronization detection result

See **4.5 Baud Rate Set Command** for details about the flowchart of the processing sequence between the programmer and the 78K0/Kx2-L, the flowchart of command processing, and the sample program.

3.3 Chip Erase Command

3.3.1 Description

This command is used to erase the entire contents of the flash memory. In addition, all of the information that is set by security setting processing can be initialized by chip erase processing, as long as erasure is not prohibited by the security setting (see **3.11 Security Set Command**).

3.3.2 Command frame and status frame

Figure 3-5 shows the format of a command frame for the Chip Erase command, and Figure 3-6 shows the status frame for the command.

Figure 3-5. Chip Erase Command Frame (from Programmer to 78K0/Kx2-L)

SOH	LEN	СОМ	SUM	ETX
01H	01H	20H (Chip Erase)	Checksum	03H

Figure 3-6. Status Frame for Chip Erase Command (from 78K0/Kx2-L to Programmer)

STX	LEN	Data	SUM	ETX
02H	01H	ST1	Checksum	03H

Remark ST1: Chip erase result

See **4.6 Chip Erase Command** for details about the flowchart of the processing sequence between the programmer and the 78K0/Kx2-L, the flowchart of command processing, and the sample program.

3.4 Block Erase Command

3.4.1 Description

This command is used to erase the content of flash memory of the block with the specified number.

A block can be specified with the first address of the block where erasing starts and the last address where erasing ends. Successive multiple blocks can be specified.

Erasing cannot be performed, however, if erasing is prohibited due to the security setting (see 3.11 Security Set Command).

3.4.2 Command frame and status frame

Figure 3-7 shows the format of a command frame for the Block Erase command, and Figure 3-8 shows the status frame for the command.

Figure 3-7. Block Erase Command Frame (from Programmer to 78K0/Kx2-L)

SOH	LEN	СОМ	Command Information SUM	ETX
01H	07H	22H (Block Erase)	SAHSAMSALEAHEAMEAL Checksun	03H

Remark SAH, SAM, SAL: Block erase start address (start address of any block)

SAH: Start address, high (bits 23 to 16) SAM: Start address, middle (bits 15 to 8) SAL: Start address, low (bits 7 to 0)

EAH, EAM, EAL: Block erase end address (last address of any block)

EAH: End address, high (bits 23 to 16)
EAM: End address, middle (bits 15 to 8)
EAL: End address, low (bits 7 to 0)

Figure 3-8. Status Frame for Block Erase Command (from 78K0/Kx2-L to Programmer)

STX	LEN	Data	SUM	ETX
02H	01H	ST1	Checksum	03H

Remark ST1: Block erase result

See **4.7 Block Erase Command** for details about the flowchart of the processing sequence between the programmer and the 78K0/Kx2-L, the flowchart of command processing, and the sample program.

3.5 Programming Command

3.5.1 Description

This command is used to write the user program to the flash memory by transmitting write data after having transmitted the write start address and the write end address. Internal verification is then executed after the last data has been transmitted and writing has been completed.

The write start/end address can be set only in the block start/end address units.

If both of the status frames (ST1 and ST2) after the last data transmission indicate ACK, the 78K0/Kx2-L firmware automatically executes internal verify. Therefore, the Status command for this internal verify must be transmitted.

3.5.2 Command frame and status frame

Figure 3-9 shows the format of a command frame for the Programming command, and Figure 3-10 shows the status frame for the command.

Figure 3-9. Programming Command Frame (from Programmer to 78K0/Kx2-L)

SOH	LEN	СОМ	Command Information				SUM	ETX		
01H	07H	40H (Programming)	SAH	SAM	SAL	EAH	EAM	EAL	Checksum	03H

Remark SAH, SAM, SAL: Write start addresses

EAH, EAM, EAL: Write end addresses

Figure 3-10. Status Frame for Programming Command (from 78K0/Kx2-L to Programmer)

STX	LEN	Data	SUM	ETX
02H	01H	ST1 (a)	Checksum	03H

Remark ST1 (a): Command reception result

3.5.3 Data frame and status frame

Figure 3-11 shows the format of a frame that includes data to be written, and Figure 3-12 shows the status frame for the data.

Figure 3-11. Data Frame to Be Written (from Programmer to 78K0/Kx2-L)

STX	LEN	Data	SUM	ETX/ETB
02H	00H to FFH (00H = 256)	Write Data	Checksum	03H/17H

Remark Write Data: User program to be written

Figure 3-12. Status Frame for Data Frame (from 78K0/Kx2-L to Programmer)

STX	LEN	Data		SUM	ETX
02H	02H	ST1 (b)	ST2 (b)	Checksum	03H

Remark ST1 (b): Data reception check result

ST2 (b): Write result

3.5.4 Completion of transferring all data and status frame

Figure 3-13 shows the status frame after transfer of all data is completed.

Figure 3-13. Status Frame After Completion of Transferring All Data (from 78K0/Kx2-L to Programmer)

STX	LEN	Data	SUM	ETX
02H	01H	ST1 (c)	Checksum	03H

Remark ST1 (c): Internal verify result

See **4.8 Programming Command** for details about the flowchart of the processing sequence between the programmer and the 78K0/Kx2-L, the flowchart of command processing, and the sample program.

3.6 Verify Command

3.6.1 Description

This command is used to compare the data transmitted from the programmer with the data read from the 78K0/Kx2-L (read level) in the specified address range, and check whether they match.

The verify start/end address can be set only in the block start/end address units.

3.6.2 Command frame and status frame

Figure 3-14 shows the format of a command frame for the Verify command, and Figure 3-15 shows the status frame for the command.

Figure 3-14. Verify Command Frame (from Programmer to 78K0/Kx2-L)

SOH	LEN	СОМ		Command Information					SUM	ETX
01H	07H	13H (Verify)	SAH	SAM	SAL	EAH	EAM	EAL	Checksum	03H

Remark SAH, SAM, SAL: Verify start addresses

EAH, EAM, EAL: Verify end addresses

Figure 3-15. Status Frame for Verify Command (from 78K0/Kx2-L to Programmer)

STX	LEN	Data	SUM	ETX
02H	01H	ST1 (a)	Checksum	03H

Remark ST1 (a): Command reception result

3.6.3 Data frame and status frame

Figure 3-16 shows the format of a frame that includes data to be verified, and Figure 3-17 shows the status frame for the data.

Figure 3-16. Data Frame of Data to Be Verified (from Programmer to 78K0/Kx2-L)

STX	LEN	Data	SUM	ETX/ETB
02H	00H to FFH (00H = 256)	Verify Data	Checksum	03H/17H

Remark Verify Data: User program to be verified

Figure 3-17. Status Frame for Data Frame (from 78K0/Kx2-L to Programmer)

STX	LEN	Da	ata	SUM	ETX
02H	02H	ST1 (b)	ST2 (b)	Checksum	03H

Remark ST1 (b): Data reception check result

ST2 (b): Verify result^{Note}

Note Even if a verify error occurs in the specified address range, ACK is always returned as the verify result. The status of all verify errors are reflected in the verify result for the last data. Therefore, the occurrence of verify errors can be checked only when all the verify processing for the specified address range is completed.

See **4.9 Verify Command** for details about the flowchart of the processing sequence between the programmer and the 78K0/Kx2-L, the flowchart of command processing, and the sample program.

3.7 Block Blank Check Command

3.7.1 Description

This command is used to check if a block in the flash memory, with a specified block number, is blank (erased state).

A block can be specified with the start address of the blank check start block and the last address of the blank check end block. Successive multiple blocks can be specified.

3.7.2 Command frame and status frame

Figure 3-18 shows the format of a command frame for the Block Blank Check command, and Figure 3-19 shows the status frame for the command.

Figure 3-18. Block Blank Check Command Frame (from Programmer to 78K0/Kx2-L)

SOH	LEN	СОМ		С	omma	nd Info	rmatic	n		SUM	ETX
01H	08H	32H (Block Blank Check)	SAH	SAM	SAL	EAH	EAM	EAL	D01	Checksum	03H

Remark SAH, SAM, SAL: Block blank check start address (start address of any block)

SAH: Start address, high (bits 23 to 16)
SAM: Start address, middle (bits 15 to 8)
SAL: Start address, low (bits 7 to 0)

EAH, EAM, EAL: Block blank check end address (last address of any block)

EAH: End address, high (bits 23 to 16)
EAM: End address, middle (bits 15 to 8)
EAL: End address, low (bits 7 to 0)

D01: 00H: When performing a block blank check for a single block

01H: When performing a blank check for the complete area before

erasing the chip

Figure 3-19. Status Frame for Block Blank Check Command (from 78K0/Kx2-L to Programmer)

STX	LEN	Data	SUM	ETX	
02H	01H	ST1	Checksum	03H	

Remark ST1: Block blank check result

See **4.10 Block Blank Check Command** for details about the flowchart of the processing sequence between the programmer and the 78K0/Kx2-L, the flowchart of command processing, and the sample program.

3.8 Silicon Signature Command

3.8.1 Description

This command is used to read information such as the write protocol information (silicon signature) of the device and security flag information.

3.8.2 Command frame and status frame

Figure 3-20 shows the format of a command frame for the Silicon Signature command, and Figure 3-21 shows the status frame for the command.

Figure 3-20. Silicon Signature Command Frame (from Programmer to 78K0/Kx2-L)

SOH	LEN	СОМ	SUM	ETX
01H	01H	C0H (Silicon Signature)	Checksum	03H

Figure 3-21. Status Frame for Silicon Signature Command (from 78K0/Kx2-L to Programmer)

STX	LEN	Data	SUM	ETX
02H	01H	ST1	Checksum	03H

Remark ST1: Command reception result

3.8.3 Silicon signature data frame

Figure 3-22 shows the format of a frame that includes silicon signature data.

Figure 3-22. Silicon Signature Data Frame (from 78K0/Kx2-L to Programmer)

STX	LEN		Data						
02H	1BH	VEN	MET	MSC	DEC1	DEC2	DEC3	UAE(3)	DEV(10)

Da	ta (continue	SUM	ETX	
SCF	SCF BOT RES(6)		Checksum	03H

Remarks 1. LEN: Data length

VEN: Vendor code (NEC: 10H)
MET: Macro extension code
MSC: Macro function code
DEC1: Device extension code 1
DEC2: Device extension code 2
DEC3: Device extension code 3

UAE: User flash ROM last address (3 bytes)

DEV: Device name (10 bytes)
SCF: Security flag information
BOT: Boot block number
RES: Reserved (6 bytes)

2. For the vendor code (VEN), extension code (MET), function code (MSC), device extension code 1 (DEC1), device extension code 2 (DEC2), and device extension code 3 (DEC3), the lower 7 bits are used as data entity, and the highest bit is used as an odd parity. The following shows an example.

Table 3-2. Example of Silicon Signature Data (μPD78F0588 (78K0/KC2-L))

Field Name	Description	Length (Byte)	Example	e of Silicon Signature Data	Actual Value	Parity
VEN	Vendor code (NEC)	1	10H	(00010000B)	10H	Added
MET	Macro extension code	1	DFH	(11011111B)	DFH	Added
MSC	Macro function code	1	04H	(01000000B)	04H	Added
DEC1	Device extension code 1	1	FEH	(11111110B)	FEH	Added
DEC2	Device extension code 2	1	7FH	(01111111B)	7FH	Added
DEC3	Device extension code 3	1	7FH	(01111111B)	7FH	Added
UAE	User flash ROM last address	3	FFH	(11111111B)	007FFFH	Not added
			7FH	(01111111B)		
			00H	(0000000B)		
DEV	Device name	10	44H	(01000100B) = 'D'	'D'	Not added
			37H	(00110111B) = '7'	'7'	
			38H	(00111000B) = '8'	'8'	
			46H	(01000110B) = 'F'	'F'	
			30H	(00110001B) = '0'	·0'	
			35H	(00110000B) = '5'	'5'	
			38H	(00111000B) = '8'	'8'	
			38H	(00111000B) = '8'	'8'	
			20H	(00100000B) = ' '	٠,	
			20H	(00100000B) = ' '	٠,	
SCF	Security flag information	1	Any		Same as left column	Not added
ВОТ	Boot block number (fixed)	1	03H	(00000011B)	03H	Not added
RES	Reserved	6	FFFFFF	FFFFFH	FFFFFFFFFH	Not added

See **4.11 Silicon Signature Command** for details about the flowchart of the processing sequence between the programmer and the 78K0/Kx2-L, the flowchart of command processing, and the sample program.

3.8.4 Silicon signature list

(1) 78K0/Kx2-L silicon signature data list

Table 3-3. 78K0/Kx2-L Silicon Signature Data List

Item	Description	Length (Byte)	Data (Hex)
Vendor code	NEC	1	10
Extension code	Extension code	1	DF
Function code	Function information	1	04
Device information	Device information	3	FE
			7F
			7F
Internal flash ROM last address	Transmitted from lower bytes of address	3	Note 1
Device name (µPD)	78F0550, 78F0551, 78F0552, 78F0555, 78F0556, 78F0557 78F0560, 78F0561, 78F0562, 78F0565, 78F0566, 78F0567 78F0571, 78F0572, 78F0573, 78F0576, 78F0577, 78F0578 78F0581, 78F0582, 78F0583, 78F0586, 78F0587, 78F0588	10	Note 2
Security information	Security information	1	Any
Boot block number	The last block number of the boot cluster that is currently selected	1	03
Reserved	Reserved	6	FFFFFFFFF

Note 1. The list of internal flash ROM last addresses is as follows.

Item	Description	Length (Byte)	Data (Hex)
Internal flash ROM	4 KB (0FFFH)	3	FF0F00
last address	8 KB (1FFFH)		FF1F00
	16 KB (3FFFH)		FF3F00
	32 KB (7FFFH)		FF7F00

(Note 2 is on the next page.)

Note 2. The device names are listed below.

Device name list

												(1/2)
Generic Name	Part Number	Length					Actual	Value				
		(Bytes)				Upper	Row: S	Signatur	e Data			
						Lower	Row: C	haracte	r Code	J		
78K0/KY2-L	D78F0550	10	44	37	38	46	30	35	35	30	20	20
			D	7	8	F	0	5	5	0	_	_
	D78F0551		44	37	38	46	30	35	35	31	20	20
			D	7	8	F	0	5	5	1	_	_
	D78F0552		44	37	38	46	30	35	35	32	20	20
			D	7	8	F	0	5	5	2	_	_
	D78F0555		44	37	38	46	30	35	35	35	20	20
			D	7	8	F	0	5	5	5	_	_
	D78F0556		44	37	38	46	30	35	35	36	20	20
			D	7	8	F	0	5	5	6	-	_
	D78F0557		44	37	38	46	30	35	35	37	20	20
			D	7	8	F	0	5	5	7	_	_
78K0/KA2-L	D78F0560		44	37	38	46	30	35	36	30	20	20
			D	7	8	F	0	5	6	0	_	_
	D78F0561		44	37	38	46	30	35	36	31	20	20
			D	7	8	F	0	5	6	1	_	_
	D78F0562		44	37	38	46	30	35	36	32	20	20
			D	7	8	F	0	5	6	2	_	_
	D78F0565		44	37	38	46	30	35	36	35	20	20
			D	7	8	F	0	5	6	5	_	_
	D78F0566		44	37	38	46	30	35	36	36	20	20
			D	7	8	F	0	5	6	6	_	_
	D78F0567		44	37	38	46	30	35	36	37	20	20
			D	7	8	F	0	5	6	7	_	_
78K0/KB2-L	D78F0571		44	37	38	46	30	35	37	31	20	20
			D	7	8	F	0	5	7	1	-	_
	D78F0572		44	37	38	46	30	35	37	32	20	20
			D	7	8	F	0	5	7	2	_	_
	D78F0573		44	37	38	46	30	35	37	33	20	20
			D	7	8	F	0	5	7	3	_	_
	D78F0576		44	37	38	46	30	35	37	36	20	20
			D	7	8	F	0	5	7	6	_	_
	D78F0577		44	37	38	46	30	35	37	37	20	20
			D	7	8	F	0	5	7	7	_	_
	D78F0578		44	37	38	46	30	35	37	38	20	20
			D	7	8	F	0	5	7	8	_	_

(2/2)

Generic Name	Part Number	Length		Actual Value								
		(Bytes)		Upper Row: Signature Data								
			Lower Row: Character Code									
78K0/KC2-L	D78F0581	10	44	37	38	46	30	35	38	31	20	20
			D	7	8	F	0	5	8	1	-	-
	D78F0582		44	37	38	46	30	35	38	32	20	20
			D	7	8	F	0	5	8	2	-	-
	D78F0583		44	37	38	46	30	35	38	33	20	20
			D	7	8	F	0	5	8	3	-	-
	D78F0586		44	37	38	46	30	35	38	36	20	20
			D	7	8	F	0	5	8	6	-	-
	D78F0587		44	37	38	46	30	35	38	37	20	20
			D	7	8	F	0	5	8	7	-	-
	D78F0588		44	37	38	46	30	35	38	38	20	20
			D	7	8	F	0	5	8	8	_	_

(2) 78K0/lx2 silicon signature data list

Table 3-4. 78K0/lx2 Silicon Signature Data List

Item	Description	Length (Byte)	Data (Hex)
Vendor code	NEC	1	10
Extension code	Extension code	1	DF
Function code	Function information	1	04
Device information	Device information	3	FE
			7F
			7F
Internal flash ROM last address	Transmitted from lower bytes of address	3	Note 1
Device name (µPD)	78F0740, 78F0741, 78F0742, 78F0750, 78F0751, 78F0752 78F0743, 78F0744, 78F0753, 78F0754 78F0745, 78F0746, 78F0755, 78F0756	10	Note 2
Security information	Security information	1	Any
Boot block number	The last block number of the boot cluster that is currently selected	1	03
Reserved	Reserved	6	FFFFFFFFF

Note 1. The list of internal flash ROM last addresses is as follows.

Item	Description	Length (Byte)	Data (Hex)
Internal flash ROM	4 KB (0FFFH)	3	FF0F00
last address	8 KB (1FFFH)		FF1F00
	16 KB (3FFFH)		FF3F00

(Note 2 is on the next page.)

Note 2. The device names are listed below.

Device name list

Generic Name	Part Number	Length					Actual	Value				
		(Bytes)				Upper	Row: S	Signatur	e Data			
						Lower	Row: C	haracte	r Code	J		
78K0/IY2	D78F0740	10	44	37	38	46	30	37	34	30	20	20
			D	7	8	F	0	7	4	0	-	_
	D78F0741		44	37	38	46	30	37	34	31	20	20
			D	7	8	F	0	7	4	1	_	_
	D78F0742		44	37	38	46	30	37	34	32	20	20
			D	7	8	F	0	7	4	2	_	_
	D78F0750		44	37	38	46	30	37	35	30	20	20
			D	7	8	F	0	7	5	0	_	_
	D78F0751		44	37	38	46	30	37	35	31	20	20
			D	7	8	F	0	7	5	1	_	_
	D78F0752		44	37	38	46	30	37	35	32	20	20
			D	7	8	F	0	7	5	2	_	_
78K0/IA2	D78F0743		44	37	38	46	30	37	34	33	20	20
			D	7	8	F	0	7	4	3	_	_
	D78F0744		44	37	38	46	30	37	34	34	20	20
			D	7	8	F	0	7	4	4	_	-
	D78F0753		44	37	38	46	30	37	35	33	20	20
			D	7	8	F	0	7	5	3	_	_
	D78F0754		44	37	38	46	30	37	35	34	20	20
			D	7	8	F	0	7	5	4	_	_
78K0/IB2	D78F0745		44	37	38	46	30	37	34	35	20	20
			D	7	8	F	0	7	4	5	_	_
	D78F0746		44	37	38	46	30	37	34	36	20	20
			D	7	8	F	0	7	4	6	_	_
	D78F0755		44	37	38	46	30	37	35	35	20	20
			D	7	8	F	0	7	5	5	_	_
	D78F0756		44	37	38	46	30	37	35	36	20	20
			D	7	8	F	0	7	5	6	_	_

3.9 Version Get Command

3.9.1 Description

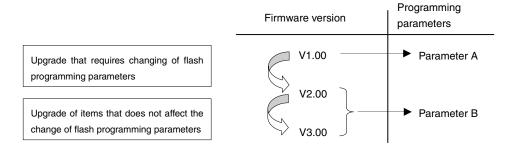
This command is used to acquire information on the 78K0/Kx2-L device version and firmware version.

The device version value is fixed to 00H.

Use this command when the programming parameters must be changed in accordance with the 78K0/Kx2-L firmware version.

Caution The firmware version may be updated during firmware update that does not affect the change of flash programming parameters (at this time, update of the firmware version is not reported).

Example: Firmware version and programming parameters



3.9.2 Command frame and status frame

Figure 3-23 shows the format of a command frame for the Version Get command, and Figure 3-24 shows the status frame for the command.

Figure 3-23. Version Get Command Frame (from Programmer to 78K0/Kx2-L)

SOH	LEN	СОМ	SUM	ETX
01H	01H	C5H (Version Get)	Checksum	03H

Figure 3-24. Status Frame for Version Get Command (from 78K0/Kx2-L to Programmer)

STX	LEN	Data	SUM	ETX
02H	01H	ST1	Checksum	03H

Remark ST1: Command reception result

3.9.3 Version data frame

Figure 3-25 shows the data frame of version data.

Figure 3-25. Version Data Frame (from 78K0/Kx2-L to Programmer)

STX	LEN			SUM	ETX				
02H	06H	DV1	DV2	DV3	FV1	FV2	FV3	Checksum	03H

Remark DV1: Integer of device version (fixed to 00H)

DV2: First decimal place of device version (fixed to 00H) DV3: Second decimal place of device version (fixed to 00H)

FV1: Integer of firmware version

FV2: First decimal place of firmware version FV3: Second decimal place of firmware version

See **4.12 Version Get Command** for details about the flowchart of the processing sequence between the programmer and the 78K0/Kx2-L, the flowchart of command processing, and the sample program.

3.10 Checksum Command

3.10.1 Description

This command is used to acquire the checksum data in the specified area.

For the checksum calculation start/end address, specify a fixed address in block units (2 KB) starting from the top of the flash memory.

Checksum data is obtained by sequentially subtracting data in the specified address range from the initial value (0000H) in 1-byte units.

3.10.2 Command frame and status frame

Figure 3-26 shows the format of a command frame for the Checksum command, and Figure 3-27 shows the status frame for the command.

Figure 3-26. Checksum Command Frame (from Programmer to 78K0/Kx2-L)

SOH	LEN	СОМ		С	SUM	ETX				
01H	07H	B0H (Checksum)	SAH	SAM	SAL	EAH	EAM	EAL	Checksum	03H

Remark SAH, SAM, SAL: Checksum calculation start addresses

EAH, EAM, EAL: Checksum calculation end addresses

Figure 3-27. Status Frame for Checksum Command (from 78K0/Kx2-L to Programmer)

STX	LEN	Data	SUM	ETX
02H	01H	ST1	Checksum	03H

Remark ST1: Command reception result

3.10.3 Checksum data frame

Figure 3-28 shows the format of a frame that includes checksum data.

Figure 3-28. Checksum Data Frame (from 78K0/Kx2-L to Programmer)

STX	LEN	Data		SUM	ETX
02H	02H	CK1	CK2	Checksum	03H

Remark CK1: Higher 8 bits of checksum data

CK2: Lower 8 bits of checksum data

See **4.13 Checksum Command** for details about the flowchart of the processing sequence between the programmer and the 78K0/Kx2-L, the flowchart of command processing, and the sample program.

3.11 Security Set Command

3.11.1 Description

This command is used to perform security settings (enabling/disabling of write, block erase, chip erase, and boot block cluster rewriting). By performing these settings with this command, rewriting of the flash memory by an unauthorized party can be restricted.

Caution Even after the security setting, additional setting of changing from enable to disable can be performed; however, changing from disable to enable is not possible. If an attempt is made to perform such a setting, a protect error (10H) will occur. If such setting is required, all of the security flags must first be initialized by executing the Chip Erase command (the Block Erase command cannot be used to initialize the security flags).

If chip erase or boot block rewrite has been disabled, however, chip erase itself will be impossible, so the settings cannot be erased from the programmer. Re-confirmation of security setting execution is therefore recommended before disabling chip erase, due to this programmer specification.

3.11.2 Command frame and status frame

Figure 3-29 shows the format of a command frame for the Security Set command, and Figure 3-30 shows the status frame for the command.

Figure 3-29. Security Set Command Frame (from Programmer to 78K0/Kx2-L)

SOH	LEN	СОМ	Command Information		SUM	ETX
01H	03H	A0H (Security Set)	00H (fixed)	00H (fixed)	Checksum	03H

Figure 3-30. Status Frame for Security Set Command (from 78K0/Kx2-L to Programmer)

STX	LEN	Data	SUM	ETX
02H	01H	ST1 (a)	Checksum	03H

Remark ST1 (a): Command reception result

3.11.3 Data frame and status frame

Figure 3-31 shows the format of a security data frame, and Figure 3-32 shows the status frame for the data.

Figure 3-31. Security Data Frame (from Programmer to 78K0/Kx2-L)

STX	LEN	Data				SUM	ETX				
02H	08H	FLG	вот	FFH	FFH	FFH	FFH	FFH	FFH	Checksum	03H
				(fixed)	(fixed)	(fixed)	(fixed)	(fixed)	(fixed)		

Remark FLG: Security flag

BOT: Boot block cluster last block number (fixed to 03H)

Figure 3-32. Status Frame for Security Data Writing (from 78K0/Kx2-L to Programmer)

STX	LEN	Data	SUM	ETX
02H	01H	ST1 (b)	Checksum	03H

Remark ST1 (b): Security data write result

3.11.4 Internal verify check and status frame

Figure 3-33 shows the status frame for internal verify check.

Figure 3-33. Status Frame for Internal Verify Check (from 78K0/Kx2-L to Programmer)

STX	LEN	Data	SUM	ETX
02H	01H	ST1 (c)	Checksum	03H

Remark ST1 (c): Internal verify result

The following table shows the contents in the security flag field.

Table 3-5. Contents of Security Flag Field

Item	Contents
Bit 7	Fixed to "1"
Bit 6	
Bit 5	
Bit 4	Boot block cluster rewrite disable flag (1: Enable, 0: Disable)
Bit 3	Fixed to "1"
Bit 2	Programming disable flag (1: Enable programming, 0: Disable programming)
Bit 1	Block erase disable flag (1: Enable block erase, 0: Disable block erase)
Bit 0	Chip erase disable flag (1: Enable chip erase, 0: Disable chip erase)

The following table shows the relationship between the security flag field settings and the enable/disable status of each operation.

Table 3-6. Security Flag Field and Enable/Disable Status of Each Operation

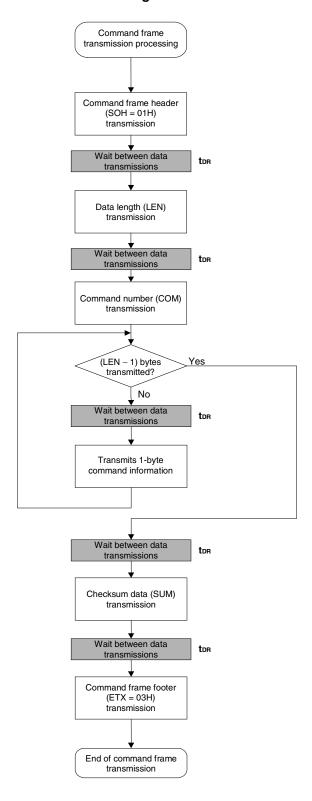
Operating Mode	Flash Memory Programming Mode			
Command	Command operation after security setting √: Execution possible, ×: Execution impossible △: Writing and block erase in boot block cluster area are impossible			
Setting Item	Programming	Chip Erase	Block Erase	
Disable programming	×	√	×	
Disable chip erase	√	×	×	
Disable block erase	√	√	×	
Boot block cluster rewrite disable flag	Δ	×	Δ	

Remark For the enable/disable status of each command in the self programming mode, see the user's manual of each device.

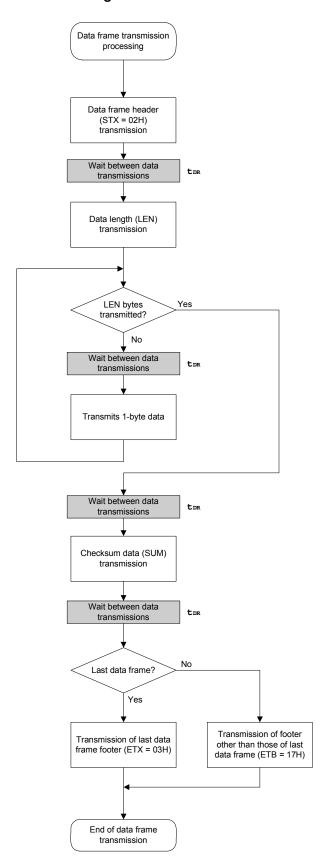
See **4.14 Security Set Command** for details about the flowchart of the processing sequence between the programmer and the 78K0/Kx2-L, the flowchart of command processing, and the sample program.

CHAPTER 4 UART COMMUNICATION MODE

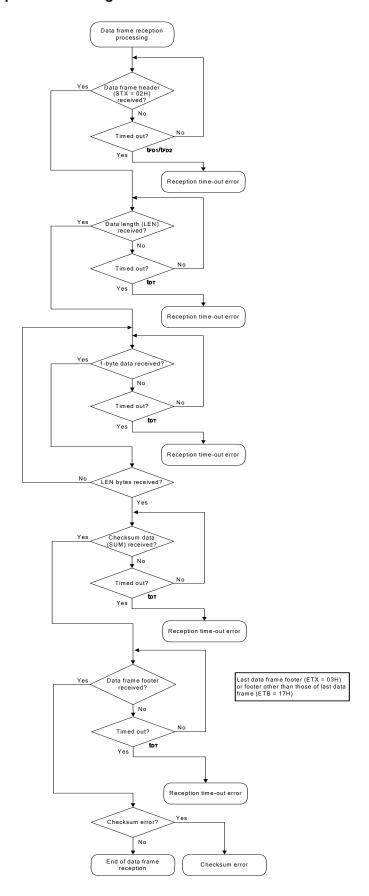
4.1 Command Frame Transmission Processing Flowchart



4.2 Data Frame Transmission Processing Flowchart



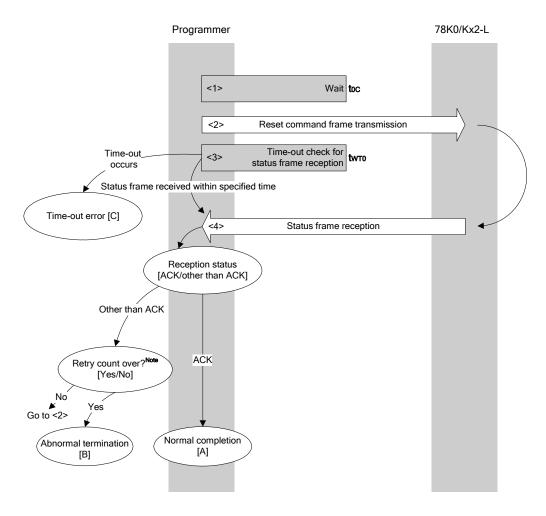
4.3 Data Frame Reception Processing Flowchart



4.4 Reset Command

4.4.1 Processing sequence chart

Reset command processing sequence



Note Do not exceed the retry count for the reset command transmission (up to 16 times).

4.4.2 Description of processing sequence

<1> Wait state (wait time toc).

<2> The Reset command is transmitted by command frame transmission processing.

<3> A time-out check is performed from command transmission until status frame reception. If a time-out occurs, a time-out error [C] is returned (time-out time t_{WTO}).

<4> The status code is checked.

When ST1 = ACK: Normal completion [A]

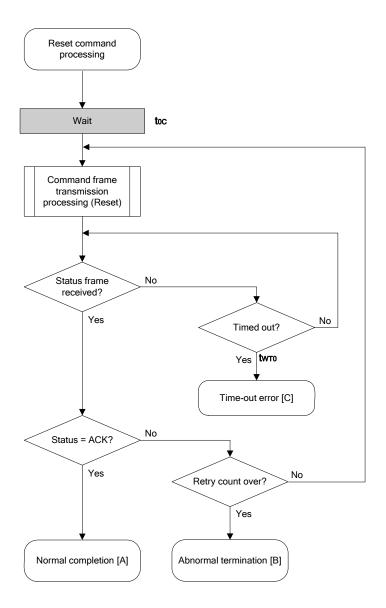
When ST1 \neq ACK: The retry count (trs) is checked.

The sequence is re-executed from <2> if the retry count is not over. If the retry count is over, the processing ends abnormally [B].

4.4.3 Status at processing completion

Status at F	Status at Processing Completion		Description
Normal completion [A]	Normal acknowledgment (ACK)	06H	The command was executed normally and synchronization between the programmer and the 78K0/Kx2-L has been established.
Abnormal termination [B]	Checksum error	07H	The checksum of the transmitted command frame does not match.
	Negative acknowledgment (NACK)	15H	Command frame data is abnormal (such as invalid data length (LEN) or no ETX).
Time-out error [0	0]	_	The status frame was not received within the specified time.

4.4.4 Flowchart



4.4.5 Sample program

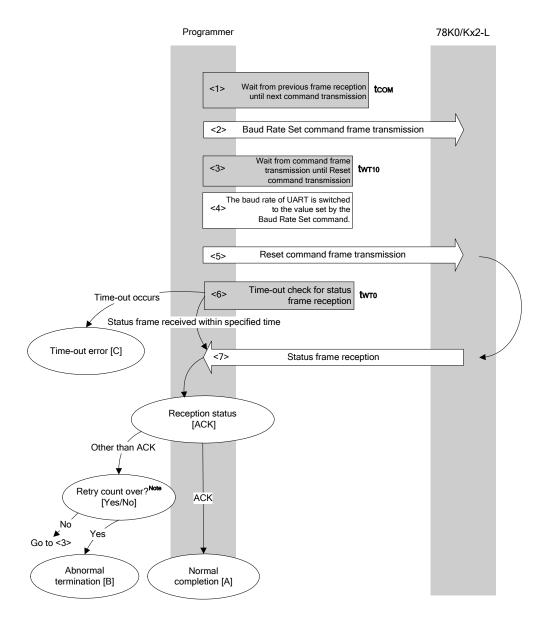
The following shows a sample program for Reset command processing.

```
*/
/*
  Reset command
                                             */
/*
                                             * /
[r] u16
                ... error code
u16 fl_ua_reset(void)
   u16
      rc;
   u32 retry;
   set_uart0_br(BR_125000); // change to 125000bps
// fl_wait(tCOM);
                     // wait
   set_ua_dir_rx();
                    // Change Mono-wire UART receive mode
   for (retry = 0; retry < tRS; retry++) {</pre>
       fl_wait(t0C_MIN); // wait
       rc = get_sfrm_ua(fl_ua_sfrm, tWT0_TO);
       if (rc == FLC_DFTO_ERR) // t.o. ?
                         // yes // case [C]
            break;
       if (rc == FLC_ACK) {
                        // ACK ?
                         // yes // case [A]
           break;
       }
       else{
            NOP();
       //continue;
                              // case [B] (if exit from loop)
   }
//
  switch(rc) {
//
       case FLC_NO_ERR: return rc; break; // case [A]
//
       case FLC_DFTO_ERR: return rc; break; // case [C]
//
//
       default: return rc; break; // case [B]
//
  return rc;
}
```

4.5 Baud Rate Set Command

4.5.1 Processing sequence chart

Baud Rate Set command processing sequence



Note Do not exceed the retry count for the reset command transmission (up to 16 times).

4.5.2 Description of processing sequence

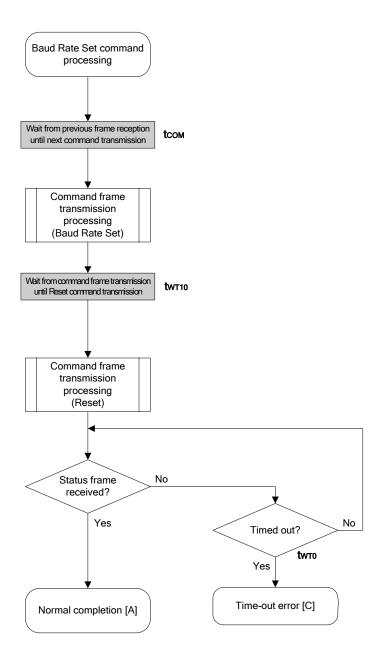
- <1> Waits from the previous frame reception until the next command transmission (wait time tcom).
- <2> The Baud Rate Set command is transmitted by command frame transmission processing.
- <3> Waits from command transmission until Reset command transmission (wait time twT10).
- <4> The baud rate of UART communication is switched to the value set by the Baud Rate Set command.
- <5> The Reset command is transmitted by command frame transmission processing.
- <6> A time-out check is performed from command transmission until status frame reception. If a time-out occurs, a time-out error [C] is returned (time-out time twro).
- <7> Since the status code should be ACK, the processing ends normally [A].

4.5.3 Status at processing completion

Status at F	Processing Completion	Status Code	Description
Normal completion [A]	Normal acknowledgment (ACK)	06H	The command was executed normally and the synchronization of the UART communication speed has been established between the programmer and the 78K0/Kx2-L.
Abnormal termination [B]	Checksum error	07H	The checksum of the transmitted command frame does not match.
	Negative acknowledgment (NACK)	15H	Command frame data is abnormal (such as invalid data length (LEN) or no ETX).
Time-out error [0	Time-out error [C] ^{Note}		Data frame reception was timed out. With the 78K0/Kx2-L, this command also results in errors in the following cases.
			Command information (D01, D02H, D02L, D03, D04) is invalid
			The Reset command was not detected after setting the baud rate and receiving command frame data for 16 times.

Note If a time-out error has occurred, execute a hardware reset and re-set to the flash memory programming mode.

4.5.4 Flowchart



4.5.5 Sample program

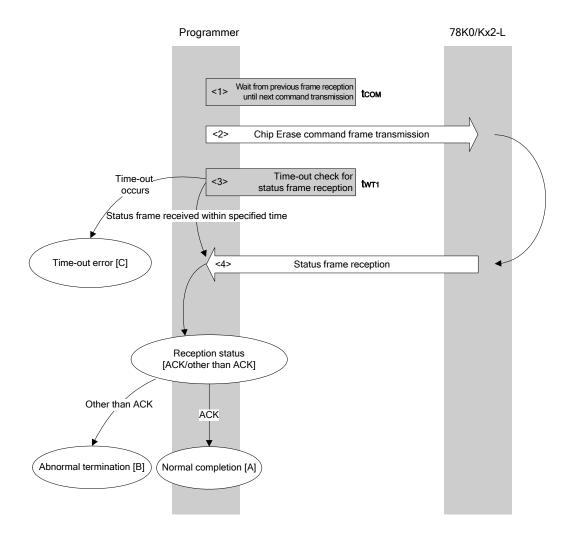
The following shows a sample program for Baud Rate Set command processing.

```
* /
/*Set baudrate command
                                           * /
                                           * /
* /
/* [i] u8 brid ... baudrate ID
/* [r] u16 ... error code
                                           */
u16 fl_ua_setbaud(u8 brid)
{
   u16
       rc;
   u8
      br;
   u32 retry;
  switch(brid) {
       default:
                            break;
       case 0x00: br = BR_125000;
       case 0x01: br = BR_250000; break;
       case 0x02: br = BR_500000;
                             break;
   }
   fl_wait(tCOM);
                     // wait before sending command
   put_cmd_ua(FL_COM_SET_BAUDRATE, 1+5, fl_cmd_prm); // send "Baudrate Set"
command
                                       // change baud-rate
   set_flbaud(br);
                                       // change baud-rate (h.w.)
   set_uart0_br(br);
   retry = tRS;
   while(1){
       fl_wait(tWT10);
       rc = get_sfrm_ua(fl_ua_sfrm, tWT0_TO);
                                      // get status frame
       if (rc) {
           if (retry--)
                continue;
            else
                return rc;
       break; // got ACK !!
   }
//
  switch(rc) {
       case FLC_NO_ERR: return rc; break; // case [A]
//
       case FLC_DFTO_ERR: return rc; break; // case [C]
//
//
       default:
               return rc; break; // case [B]
//
  }
  return rc;
}
```

4.6 Chip Erase Command

4.6.1 Processing sequence chart

Chip Erase command processing sequence



4.6.2 Description of processing sequence

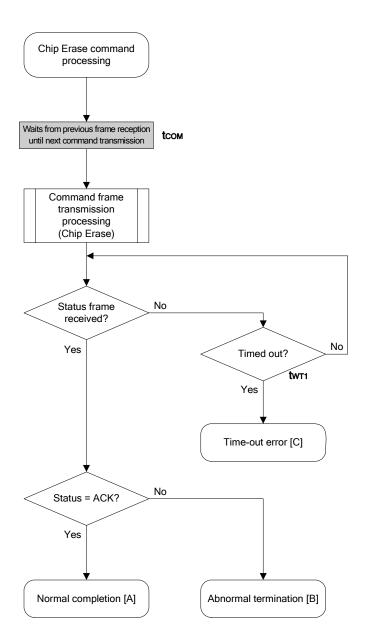
- <1> Waits from the previous frame reception until the next command transmission (wait time tcom).
- <2> The Chip Erase command is transmitted by command frame transmission processing.
- <3> A time-out check is performed from command transmission until status frame reception. If a time-out occurs, a time-out error [C] is returned (time-out time t_{WT1}).
- <4> The status code is checked.

When ST1 = ACK: Normal completion [A] When ST1 ≠ ACK: Abnormal termination [B]

4.6.3 Status at processing completion

Status at F	Processing Completion	Status Code	Description
Normal completion [A]	Normal acknowledgment (ACK)	06H	The command was executed normally and chip erase was performed normally.
Abnormal termination [B]	Checksum error	07H	The checksum of the transmitted command frame does not match.
	Protect error	10H	Chip erase or boot block rewrite is prohibited in the security setting.
	Negative acknowledgment (NACK)	15H	Command frame data is abnormal (such as invalid data length (LEN) or no ETX).
	MRG10 error	1AH	An erase error has occurred.
	MRG11 error	1BH	
	Write error	1CH	
Time-out error [0)	-	The status frame was not received within the specified time.

4.6.4 Flowchart



4.6.5 Sample program

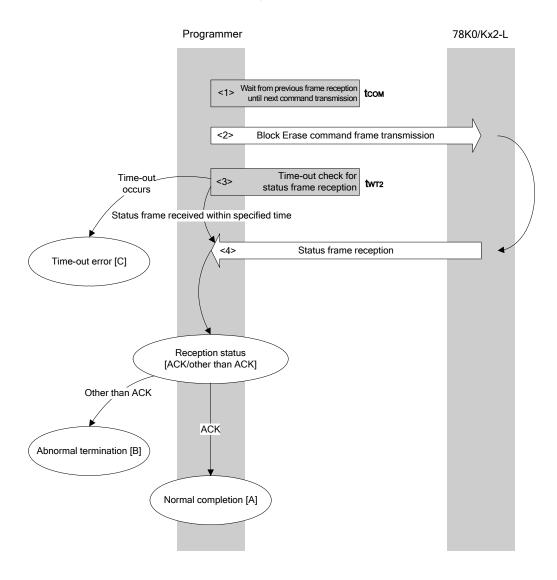
The following shows a sample program for Chip Erase command processing.

```
* /
/*
  Erase all(chip) command
                                           * /
/*
                                           * /
/* [r] u16
                ... error code
fl_ua_erase_all(void)
u16
{
   u16 rc;
   fl_wait(tCOM);
                    // wait before sending command
   put_cmd_ua(FL_COM_ERASE_CHIP, 1, fl_cmd_prm); // send ERASE CHIP command
   rc = get_sfrm_ua(fl_ua_sfrm, tWT1_MAX); // get status frame
//
  switch(rc) {
//
//
       case FLC_NO_ERR: return rc;
                                  break; // case [A]
//
       case FLC_DFTO_ERR: return rc;
                                 break; // case [C]
                                 break; // case [B]
//
       default:
                    return rc;
//
  }
  return rc;
}
```

4.7 Block Erase Command

4.7.1 Processing sequence chart

Block Erase command processing sequence



4.7.2 Description of processing sequence

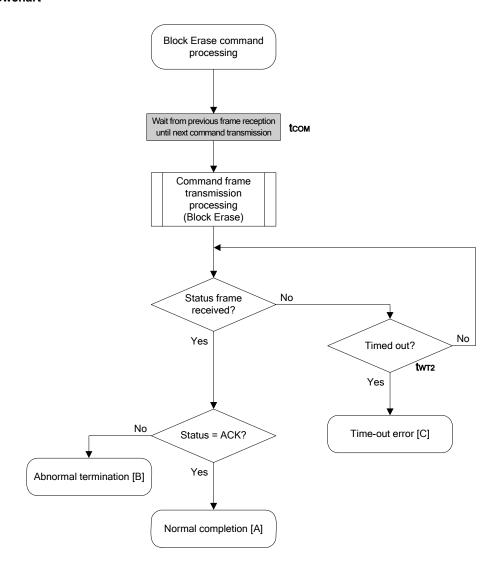
- <1> Waits from the previous frame reception until the next command transmission (wait time tcom).
- <2> The Block Erase command is transmitted by command frame transmission processing.
- <3> A time-out check is performed from command transmission until status frame reception. If a time-out occurs, a time-out error [C] is returned (time-out time t_{WT2}).
- <4> The status code is checked.

When ST1 = ACK: Normal completion [A]
When ST1 ≠ ACK: Abnormal termination [B]

4.7.3 Status at processing completion

Status at F	Processing Completion	Status Code	Description
Normal completion [A]	Normal acknowledgment (ACK)	06H	The command was executed normally and block erase was performed normally.
Abnormal termination [B]	Parameter error	05H	The specified end address is out of the flash memory range, or the specified start/end address is not the first/end address of the block.
	Checksum error	07H	The checksum of the transmitted command frame does not match.
	Protect error	10H	Write, block erase, or chip erase is prohibited in the security setting. A boot block is included in the specified range and boot block rewrite is prohibited.
	Negative acknowledgment (NACK)	15H	Command frame data is abnormal (such as invalid data length (LEN) or no ETX).
	MRG10 error	1AH	An erase error has occurred.
Time-out error [0	0]	_	The status frame was not received within the specified time.

4.7.4 Flowchart



4.7.5 Sample program

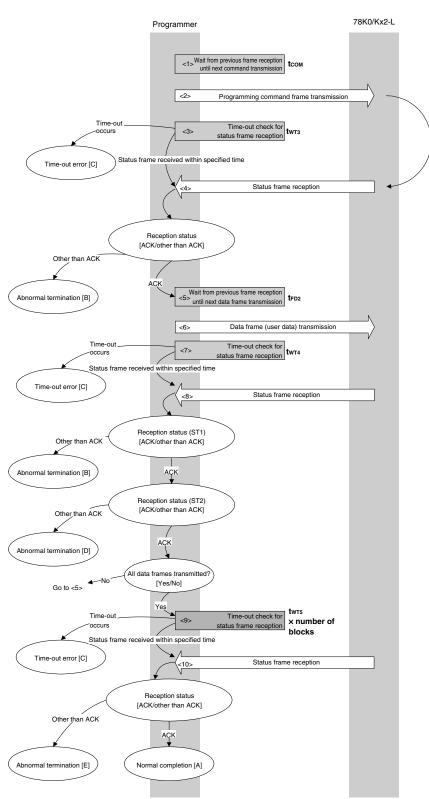
The following shows a sample program for Block Erase command processing.

```
* /
/*
  Erase block command
                                           * /
/*
                                           * /
[i] u8 block ... block number
/* [r] u16
               ... error code
                                           */
fl_ua_erase_blk(u16 sblk, u16 eblk)
{
   u16
      rc;
   u32
      wt2_max;
   u32 top, bottom;
  // get start address of start block
   set_range_prm(fl_cmd_prm, top, bottom); // set SAH/SAM/SAL, EAH/EAM/EAL
   wt2_max = make_wt2_max(sblk, eblk);
   fl_wait(tCOM);
                              // wait before sending command
   put_cmd_ua(FL_COM_ERASE_BLOCK, 1+6, fl_cmd_prm); // send ERASE_CHIP_command
   rc = get_sfrm_ua(fl_ua_sfrm, wt2_max); // get status frame
//
  switch(rc) {
//
//
       case FLC_NO_ERR: return rc;
                                 break; // case [A]
       case FLC_DFTO_ERR: return rc;
//
                                 break; // case [C]
       default: return rc;
                                 break; // case [B]
//
// }
  return rc;
}
```

4.8 Programming Command

4.8.1 Processing sequence chart

Programming command processing sequence



4.8.2 Description of processing sequence

- <1> Waits from the previous frame reception until the next command transmission (wait time tcom).
- <2> The Programming command is transmitted by command frame transmission processing.
- <3> A time-out check is performed from command transmission until status frame reception. If a time-out occurs, a time-out error [C] is returned (time-out time twrs).
- <4> The status code is checked.

When ST1 = ACK: Proceeds to <5>.

When ST1 ≠ ACK: Abnormal termination [B]

- <5> Waits from the previous frame reception until the next data frame transmission (wait time tfp2).
- <6> User data is transmitted by data frame transmission processing.
- <7> A time-out check is performed from user data transmission until data frame reception. If a time-out occurs, a time-out error [C] is returned (time-out time twv4).
- <8> The status code (ST1/ST2) is checked (also refer to the processing sequence chart and flowchart).

When ST1 ≠ ACK: Abnormal termination [B]

When ST1 = ACK: The following processing is performed according to the ST2 value.

- When ST2 = ACK: Proceeds to <9> when transmission of all data frames is completed.
 If there still remain data frames to be transmitted, the processing re-executes the sequence from <5>.
- When ST2 ≠ ACK: Abnormal termination [D]
- <9> A time-out check is performed until status frame reception.

If a time-out occurs, a time-out error [C] is returned (time-out time twis × number of blocks).

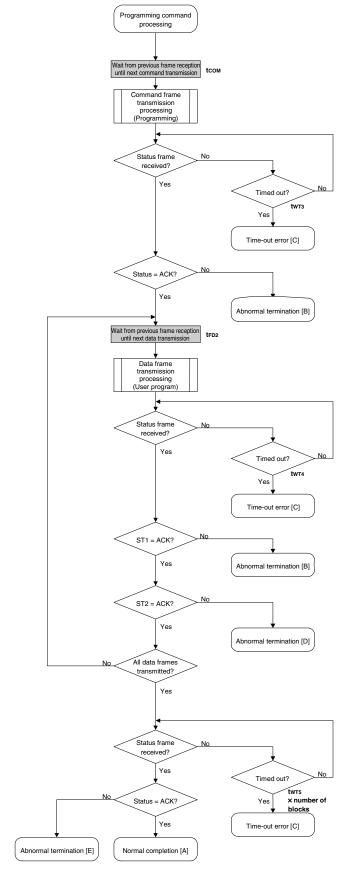
<10> The status code is checked.

When ST1 = ACK: Normal completion [A]
When ST1 ≠ ACK: Abnormal termination [E]

4.8.3 Status at processing completion

Status at F	Status at Processing Completion		Description
Normal completion [A]	Normal acknowledgment (ACK)	06H	The command was executed normally and the user data was written normally.
Abnormal termination [B]	Parameter error	05H	The start/end address is out of the flash memory range, or the specified start/end address is not the first/end address of the block.
	Checksum error	07H	The checksum of the transmitted command frame or data frame does not match.
	Protect error	10H	Write is prohibited in the security setting. A boot block is included in the specified range and boot block rewrite is prohibited.
	Negative acknowledgment (NACK)	15H	Command frame data or data frame data is abnormal (such as invalid data length (LEN) or no ETX).
Time-out error [0)]	_	The status frame was not received within the specified time.
Abnormal	MRG10 error	1AH	A write error has occurred.
termination	MRG11 error	1BH	
[D], [E]	Write error	1CH	

4.8.4 Flowchart



4.8.5 Sample program

The following shows a sample program for Programming command processing.

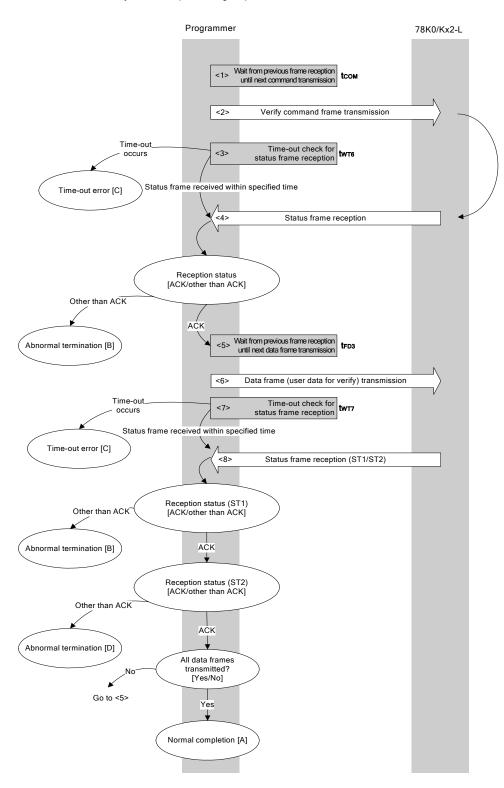
```
/************************
                                * /
/* Write command
                                */
                                * /
/* [i] u32 top ... start address
                                */
/* [i] u32 bottom ... end address
/* [r] u16 ... error code
                                * /
#define fl_st2_ua (fl_ua_sfrm[OFS_STA_PLD+1])
   fl_ua_write(u32 top, u32 bottom)
1116
{
u16 rc;
u32 send_head, send_size;
bool is_end;
u16 block_num;
set params
/***************
send command & check status
/****************
fl_wait(tCOM);
             // wait before sending command
rc = get_sfrm_ua(fl_ua_sfrm, tWT3_TO);
                       // get status frame
switch(rc) {
          FLC_NO_ERR:
                        break; // continue
   case
          FLC_DFTO_ERR: return rc; break; // case [C]
//
   case
   default:
                 return rc; break; // case [B]
}
/***************
/* send user data
/***************
send_head = top;
while(1){
    // make send data frame
    is_end = false;
                        // yes, not is_end frame
```

```
send_size = 256;
                                     // transmit size = 256 byte
         }
         else{
               is_end = true;
               send_size = bottom - send_head + 1;  // transmit size = (bottom
- send_head)+1 byte
        memcpy(fl_txdata_frm, rom_buf+send_head, send_size); // set data frame
payload
         send_head += send_size;
         fl_wait(tFD2);
                                             // wait before sending data frame
         put_dfrm_ua(send_size, fl_txdata_frm, is_end); // send user data
         rc = get_sfrm_ua(fl_ua_sfrm, tWT4_MAX);
                                                 // get status frame
         switch(rc) {
                          FLC_NO_ERR:
                                                   break; // continue
               case
               case
                          FLC_DFTO_ERR: return rc; break; // case [C]
               default:
                               return rc; break; // case [B]
         if (fl_st2_ua != FLST_ACK) {
                                             // ST2 = ACK ?
               rc = decode_status(fl_st2_ua);  // No
                                             // case [D]
               return rc;
         if (is_end)
               break;
   /****************/
          Check internally verify
   /****************
   // get status frame again
   rc = get_sfrm_ua(fl_ua_sfrm, tWT5_MAX*block_num);
   switch(rc) {
        case FLC_NO_ERR: return rc; break; // case [A]
 //
        case FLC_DFTO_ERR: return rc; break; // case [C]
         default:
                          return rc; break; // case [E]
   }
  return rc;
```

4.9 Verify Command

4.9.1 Processing sequence chart

Verify command processing sequence



4.9.2 Description of processing sequence

- <1> Waits from the previous frame reception until the next command transmission (wait time tcom).
- <2> The Verify command is transmitted by command frame transmission processing.
- <3> A time-out check is performed from command transmission until status frame reception. If a time-out occurs, a time-out error [C] is returned (time-out time twrf).
- <4> The status code is checked.

When ST1 = ACK: Proceeds to <5>.

When ST1 ≠ ACK: Abnormal termination [B]

- <5> Waits from the previous frame reception until the next data frame transmission (wait time tfp3).
- <6> User data for verifying is transmitted by data frame transmission processing.
- <7> A time-out check is performed from user data transmission until status frame reception. If a time-out occurs, a time-out error [C] is returned (time-out time twr7).
- <8> The status code (ST1/ST2) is checked (also refer to the processing sequence chart and flowchart).

When ST1 ≠ ACK: Abnormal termination [B]

When ST1 = ACK: The following processing is performed according to the ST2 value.

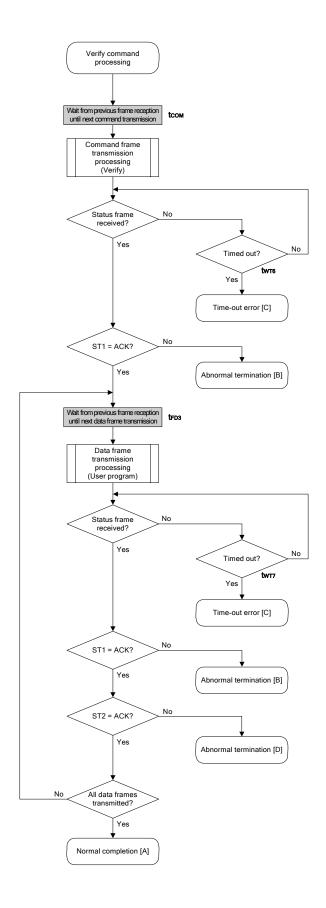
When ST2 = ACK: If transmission of all data frames is completed, the processing ends normally [A].
 If there still remain data frames to be transmitted, the processing re-executes the sequence from <5>.

• When ST2 ≠ ACK: Abnormal termination [D]

4.9.3 Status at processing completion

Status at F	Processing Completion	Status Code	Description
Normal completion [A]	Normal acknowledgment (ACK)	06H	The command was executed normally and the verify was completed normally.
Abnormal termination [B]	Parameter error	05H	The start/end address is out of the flash memory range, or the start/end address is not the start/end address of the block.
	Checksum error	07H	The checksum of the transmitted command frame or data frame does not match.
	Negative acknowledgment (NACK)	15H	Command frame data is abnormal (such as invalid data length (LEN) or no ETX).
Time-out error [C		_	The status frame was not received within the specified time.
Abnormal termination [D]	Verify error	0FH (ST2)	A verify error has occurred.

4.9.4 Flowchart



4.9.5 Sample program

The following shows a sample program for Verify command processing.

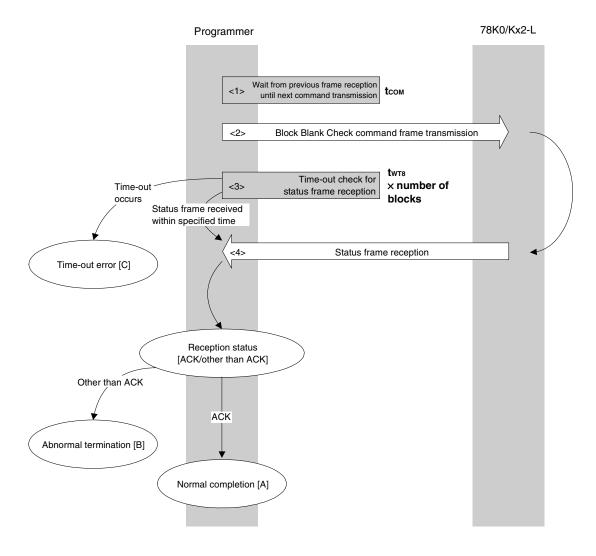
```
/* Verify command
/*
                                        * /
... start address
  [i] u32 top
                                        */
 [i] u32 bottom
              ... end address
/* [r] u16
                                        * /
               ... error code
u16
     fl_ua_verify(u32 top, u32 bottom, u8 *buf)
  u16
  u32 send_head, send_size;
  bool is_end;
  /***************/
     set params
  /****************
  set_range_prm(fl_cmd_prm, top, bottom); // set SAH/SAM/SAL, EAH/EAM/EAL
  /**********************************
      send command & check status
  /***************
  fl_wait(tCOM);
                                // wait before sending command
  rc = get_sfrm_ua(fl_ua_sfrm, tWT6_MAX);
                               // get status frame
  switch(rc) {
     case FLC_NO_ERR:
                               break; // continue
                               break; // case [C]
      case FLC_DFTO_ERR: return rc;
  //
      default:
                  return rc;
                                break; // case [B]
  }
  /**********************************
     send user data
  send_head = top;
  while(1){
       // make send data frame
       is_end = false;
                               // yes, not is_end frame
           send_size = 256;
                               // transmit size = 256 byte
       }
       else{
           is_end = true;
           send_size = bottom - send_head + 1;// transmit size = (bottom -
send_head)+1 byte
```

```
memcpy(fl_txdata_frm, buf+send_head, send_size);// set data frame payload
           send_head += send_size;
           fl_wait(tFD3);
          put_dfrm_ua(send_size, fl_txdata_frm, is_end); // send user data
           rc = get_sfrm_ua(fl_ua_sfrm, tWT7_MAX);
                                                       // get status frame
           switch(rc) {
                 case FLC_NO_ERR:
                                                        break; // continue
                 case FLC_DFTO_ERR: return rc;
           //
                                                        break; // case [C]
                 default:
                            return rc;
                                                        break; // case [B]
           if (fl_st2_ua != FLST_ACK) {
                                                 // ST2 = ACK ?
                 rc = decode_status(fl_st2_ua);
                                                 // No
                 return rc;
                                                 // case [D]
           }
           if (is_end)
                                      // send all user data ?
                 break;
                                       // yes
           //continue;
    return FLC_NO_ERR; // case [A]
}
```

4.10 Block Blank Check Command

4.10.1 Processing sequence chart

Block Blank Check command processing sequence



4.10.2 Description of processing sequence

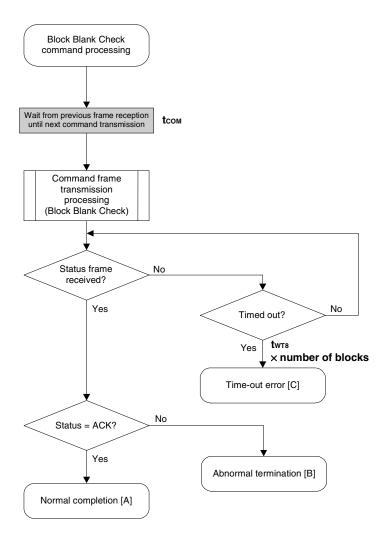
- <1> Waits from the previous frame reception until the next command transmission (wait time tcom).
- <2> The Block Blank Check command is transmitted by command frame transmission processing.
- <3> A time-out check is performed from command transmission until status frame reception. If a time-out occurs, a time-out error [C] is returned (time-out time $twr8 \times number$ of blocks).
- <4> The status code is checked.

When ST1 = ACK: Normal completion [A]
When ST1 ≠ ACK: Abnormal termination [B]

4.10.3 Status at processing completion

Status at F	Processing Completion	Status Code	Description
Normal completion [A]	Normal acknowledgment (ACK)	06H	The command was executed normally and block blank check was executed normally.
Abnormal termination [B]	Parameter error	05H	The end address is out of the flash memory range, the start/end address is not the first/end address of the block, or the value of parameter D01 is other than 00H or 01H.
	Checksum error	07H	The checksum of the transmitted command frame does not match.
	Negative acknowledgment (NACK)	15H	Command frame data is abnormal (such as invalid data length (LEN) or no ETX).
	MRG11 error	1BH	The flash memory of the specified block is not blank.
Time-out error [0	0]	_	The status frame was not received within the specified time.

4.10.4 Flowchart



4.10.5 Sample program

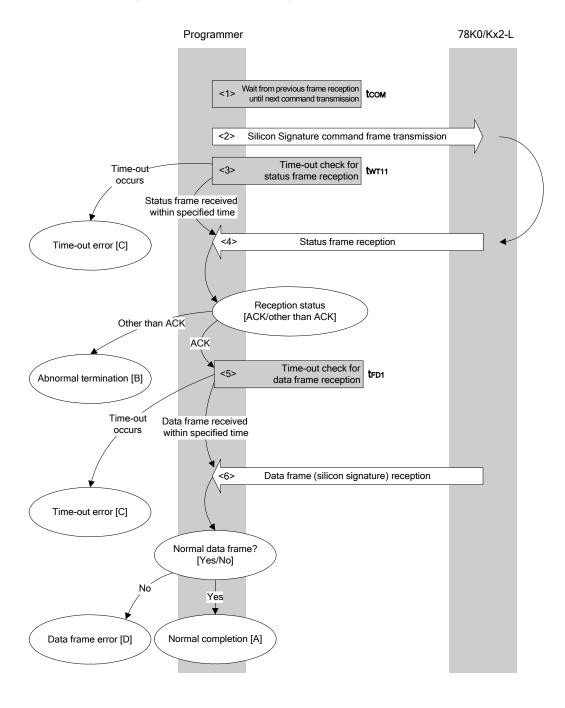
The following shows a sample program for Block Blank Check command processing.

```
*/
/* Block blank check command
                                                     * /
                                                     * /
... top address of blank check
/* [i] u32 top
/* [i] u32 bottom ... bottom address of blank check
/* [i] u8 whole ... <1>check w/NON user flash
                                                     * /
                  <0>chek only user flash
                                                     * /
/* [r] u16 ... error code
       fl_ua_blk_blank_chk(u32 top, u32 bottom, u8 whole)
{
   u16 rc;
   u16 block_num;
   set_range_prm(fl_cmd_prm, top, bottom); // set SAH/SAM/SAL, EAH/EAM/EAL
   block_num = get_block_num(top, bottom); // get block num
                                    // check only user area or not
   fl_cmd_prm[6] = whole;
   fl_wait(tCOM);
                                     // wait before sending command
   put_cmd_ua(FL_COM_BLOCK_BLANK_CHK, 7+1, fl_cmd_prm);
   rc = get_sfrm_ua(fl_ua_sfrm, tWT8_MAX * block_num); // get status frame
   switch(rc) {
//
//
                                          break; // case [A]
//
         case FLC_NO_ERR: return rc;
         case FLC_DFTO_ERR: return rc;
                                         break; // case [C]
//
                           return rc;
return rc;
//
         default:
                                         break; // case [B]
// }
   return rc;
}
```

4.11 Silicon Signature Command

4.11.1 Processing sequence chart

Silicon Signature command processing sequence



4.11.2 Description of processing sequence

- <1> Waits from the previous frame reception until the next command transmission (wait time tcom).
- <2> The Silicon Signature command is transmitted by command frame transmission processing.
- <3> A time-out check is performed from command transmission until status frame reception. If a time-out occurs, a time-out error [C] is returned (time-out time twrll).

<4> The status code is checked.

When ST1 = ACK: Proceeds to <5>.

When ST1 ≠ ACK: Abnormal termination [B]

<5> A time-out check is performed until data frame (silicon signature data) reception. If a time-out occurs, a time-out error [C] is returned (time-out time tfpl).

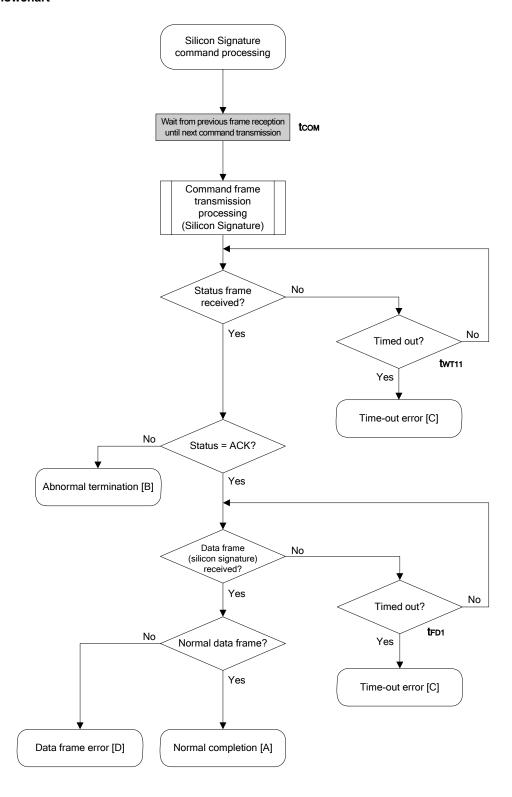
<6> The received data frame (silicon signature data) is checked.

If data frame is normal: Normal completion [A] If data frame is abnormal: Data frame error [D]

4.11.3 Status at processing completion

Status at F	Status at Processing Completion		Description
Normal completion [A]	Normal acknowledgment (ACK)	06H	The command was executed normally and silicon signature data was acquired normally.
Abnormal termination [B]	Checksum error	07H	The checksum of the transmitted command frame does not match.
	Negative acknowledgment (NACK)	15H	Command frame data is abnormal (such as invalid data length (LEN) or no ETX).
	Read error	20H	Reading security information failed.
Time-out error [C]		-	The status frame or data frame was not received within the specified time.
Data frame error	· [D]	_	The checksum of the data frame received as silicon signature data does not match.

4.11.4 Flowchart



4.11.5 Sample program

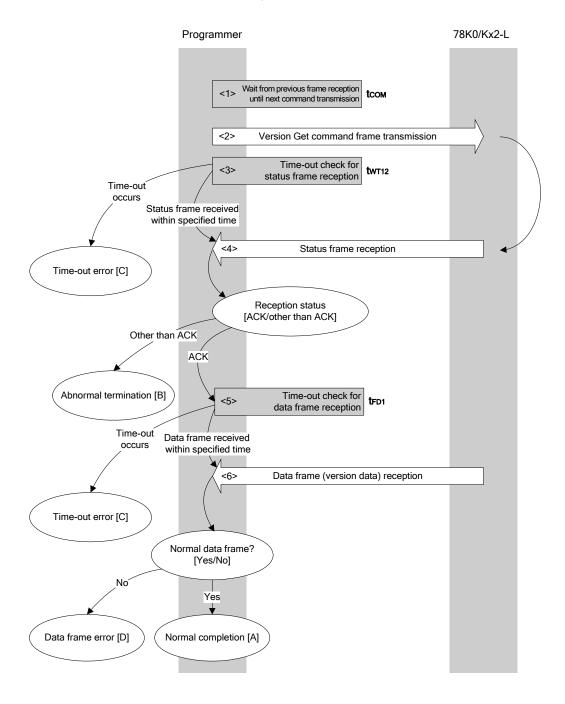
The following shows a sample program for Silicon Signature command processing.

```
*/
/* Get silicon signature command
                                                * /
                                                * /
/* [i] u8 *sig ... pointer to signature save area
/* [r] u16 ... error code
fl_ua_getsig(u8 *sig)
{
   u16 rc;
   fl_wait(tCOM);
                                       // wait before sending command
   put_cmd_ua(FL_COM_GET_SIGNATURE, 1, fl_cmd_prm); // send GET SIGNATURE command
   rc = get_sfrm_ua(fl_ua_sfrm, tWT11_MAX);
                                      // get status frame
   switch(rc) {
        case
                 FLC_NO_ERR:
                                      break; // continue
                 FLC_DFTO_ERR: return rc; break; // case [C]
        default:
                                      break; // case [B]
                            return rc;
   }
   rc = get_dfrm_ua(fl_rxdata_frm, tFD2_MAX);
                                      // get status frame
   if (rc) {
                                      // if error
                                      // case [D]
       return rc;
   memcpy(sig, fl_rxdata_frm+OFS_STA_PLD, fl_rxdata_frm[OFS_LEN]);// copy Signature
data
   return rc;
                                       // case [A]
}
```

4.12 Version Get Command

4.12.1 Processing sequence chart

Version Get command processing sequence



4.12.2 Description of processing sequence

- <1> Waits from the previous frame reception until the next command transmission (wait time tcom).
- <2> The Version Get command is transmitted by command frame transmission processing.
- <3> A time-out check is performed from command transmission until status frame reception. If a time-out occurs, a time-out error [C] is returned (time-out time twm12).

<4> The status code is checked.

When ST1 = ACK: Proceeds to <5>.

When ST1 ≠ ACK: Abnormal termination [B]

<5> A time-out check is performed until data frame (version data) reception. If a time-out occurs, a time-out error [C] is returned (time-out time t_{FD1}).

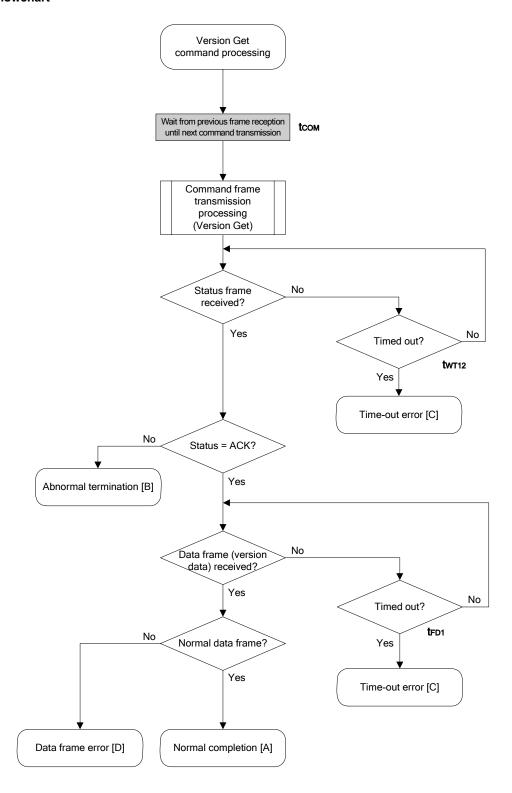
<6> The received data frame (version data) is checked.

If data frame is normal: Normal completion [A] If data frame is abnormal: Data frame error [D]

4.12.3 Status at processing completion

Status at Processing Completion		Status Code	Description
Normal completion [A]	Normal acknowledgment (ACK)	06H	The command was executed normally and version data was acquired normally.
Abnormal termination [B]	Checksum error	07H	The checksum of the transmitted command frame does not match.
	Negative acknowledgment (NACK)	15H	Command frame data is abnormal (such as invalid data length (LEN) or no ETX).
Time-out error [C]		-	The status frame or data frame was not received within the specified time.
Data frame error	Data frame error [D]		The checksum of the data frame received as version data does not match.

4.12.4 Flowchart



4.12.5 Sample program

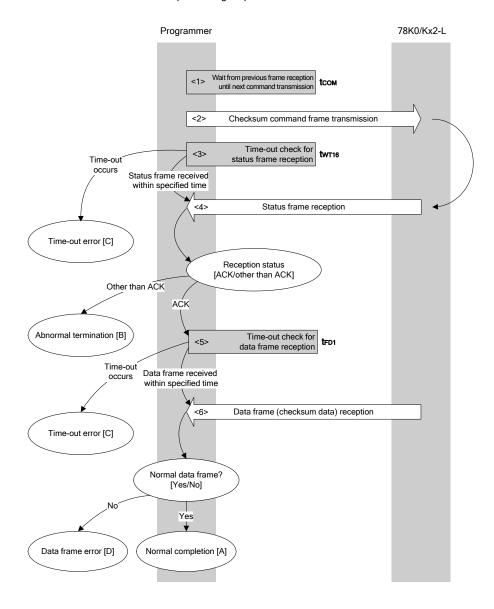
The following shows a sample program for Version Get command processing.

```
/*
   Get device/firmware version command
/* [i] u8 *buf ... pointer to version date save area
          ... error code
   [r] u16
u16 fl_ua_getver(u8 *buf)
{
   u16 rc;
   fl_wait(tCOM);
                                        // wait before sending command
   put_cmd_ua(FL_COM_GET_VERSION, 1, fl_cmd_prm); // send GET VERSION command
   rc = get_sfrm_ua(fl_ua_sfrm, tWT12_T0);
                                       // get status frame
   switch(rc) {
        case FLC_NO_ERR:
                                   break; // continue
   //
        case FLC_DFTO_ERR: return rc;
                                  break; // case [C]
         default:
                  return rc; break; // case [B]
   }
   rc = get_dfrm_ua(fl_rxdata_frm, tFD1_T0);
                                        // get data frame
   if (rc) {
        return rc;
                                        // case [D]
   }
   memcpy(buf, fl_rxdata_frm+OFS_STA_PLD, DFV_LEN);// copy version data
                                        // case [A]
   return rc;
}
u16 fl_ua_getver(u8 *buf)
{
   u16
        rc;
   fl_wait(tCOM);
                                        // wait before sending command
   put_cmd_ua(FL_COM_GET_VERSION, 1, fl_cmd_prm); // send GET VERSION command
   switch(rc) {
        case
             FLC_NO_ERR:
                                   break; // continue
        case FLC_DFTO_ERR: return rc;
   //
                                   break; // case [C]
                                 break; // case [B]
         default:
                 return rc;
   }
   rc = get_dfrm_ua(fl_rxdata_frm, tFD2_MAX);
                                        // get data frame
   if (rc) {
                                        // case [D]
        return rc:
   }
   memcpy(buf, fl_rxdata_frm+OFS_STA_PLD, DFV_LEN);// copy version data
   return rc;
                                        // case [A]
}
```

4.13 Checksum Command

4.13.1 Processing sequence chart

Checksum command processing sequence



4.13.2 Description of processing sequence

- <1> Waits from the previous frame reception until the next command transmission (wait time tcom).
- <2> The Checksum command is transmitted by command frame transmission processing.
- <3> A time-out check is performed from command transmission until status frame reception. If a time-out occurs, a time-out error [C] is returned (time-out time twr16).

<4> The status code is checked.

When ST1 = ACK: Proceeds to <5>.

When ST1 ≠ ACK: Abnormal termination [B]

<5> A time-out check is performed until data frame (checksum data) reception. If a time-out occurs, a time-out error [C] is returned (time-out time tfpl).

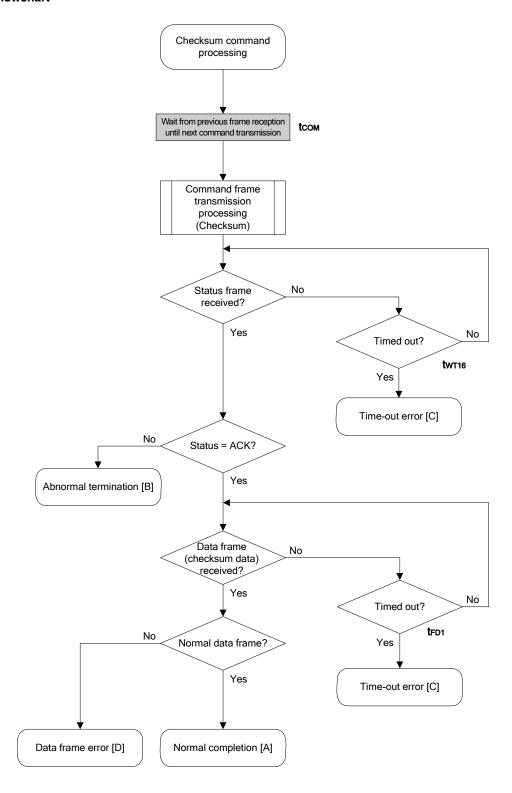
<6> The received data frame (checksum data) is checked.

If data frame is normal: Normal completion [A] If data frame is abnormal: Data frame error [D]

4.13.3 Status at processing completion

Status at F	Status at Processing Completion		Description
Normal completion [A]	Normal acknowledgment (ACK)	06H	The command was executed normally and checksum data was acquired normally.
Abnormal termination [B]	Parameter error	05H	The specified start/end address is out of the flash memory range, or the start/end address is not the start/end address of the block.
Checksum error Negative acknowledgment (NACK)		07H	The checksum of the transmitted command frame does not match.
		15H	Command frame data is abnormal (such as invalid data length (LEN) or no ETX).
Time-out error [C]		-	The status frame or data frame was not received within the specified time.
Data frame error	(D)	-	The checksum of the data frame received as checksum data does not match.

4.13.4 Flowchart



4.13.5 Sample program

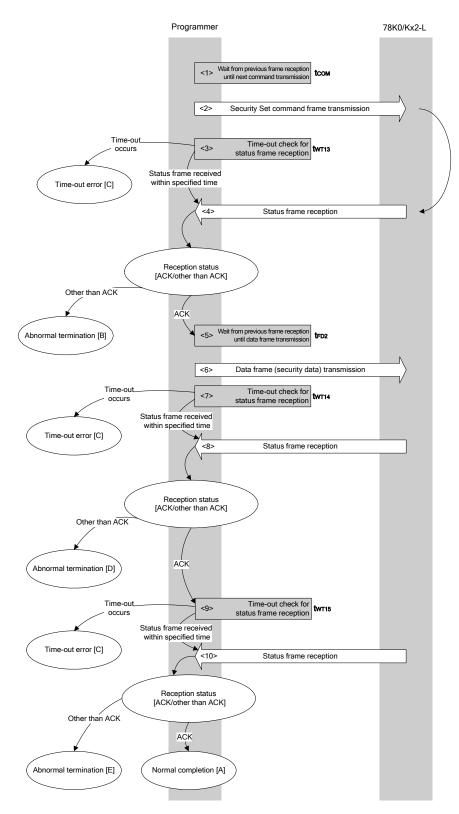
The following shows a sample program for Checksum command processing.

```
/*
                                           */
  Get checksum command
/************************
  [i] u16 *sum ... pointer to checksum save area
/*
   [i] u32 top ... start address
                                           * /
   [i] u32 bottom ... end address
/*
                                           * /
   [r] u16 ... error code
fl_ua_getsum(u16 *sum, u32 top, u32 bottom)
{
   u16 rc;
   set params
   // set params
   set_range_prm(fl_cmd_prm, top, bottom); // set SAH/SAM/SAL, EAH/EAM/EAL
   send command
   fl wait(tCOM);
                                  // wait before sending command
   put_cmd_ua(FL_COM_GET_CHECK_SUM, 7, fl_cmd_prm); // send GET VERSION command
   rc = get_sfrm_ua(fl_ua_sfrm, tWT16_MAX);
                                  // get status frame
   switch(rc) {
                                  break; // continue
      case
               FLC_NO_ERR:
      case FLC_DFTO_ERR: return rc; break; // case [C]
                         return rc; break; // case [B]
       default:
   }
   /****************
   /* get data frame (Checksum data)
                                  * /
   /***********************************
   rc = get_dfrm_ua(fl_rxdata_frm, tFD1_MAX);
                                  // get status frame
   if (rc) {
                                  // if no error,
      return rc;
                                  // case [D]
   }
   *sum = (fl_rxdata_frm[OFS_STA_PLD] << 8) + fl_rxdata_frm[OFS_STA_PLD+1]; // set
SUM data
                                  // case [A]
  return rc;
}
```

4.14 Security Set Command

4.14.1 Processing sequence chart

Security Set command processing sequence



4.14.2 Description of processing sequence

- <1> Waits from the previous frame reception until the next command transmission (wait time tcom).
- <2> The Security Set command is transmitted by command frame transmission processing.
- <3> A time-out check is performed from command transmission until status frame reception. If a time-out occurs, a time-out error [C] is returned (time-out time twrl3).

<4> The status code is checked.

When ST1 = ACK: Proceeds to <5>.

When ST1 ≠ ACK: Abnormal termination [B]

- <5> Waits from the previous frame reception until the next data frame transmission (wait time tfp3).
- <6> The data frame (security setting data) is transmitted by data frame transmission processing.
- <7> A time-out check is performed until status frame reception.

If a time-out occurs, a time-out error [C] is returned (time-out time twm14).

<8> The status code is checked.

When ST1 = ACK: Proceeds to <9>.

When ST1 ≠ ACK: Abnormal termination [D]

<9> A time-out check is performed until status frame reception.

If a time-out occurs, a time-out error [C] is returned (time-out time twels).

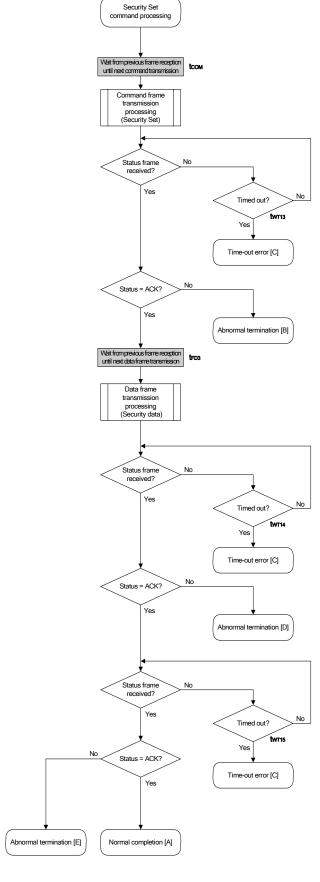
<10> The status code is checked.

When ST1 = ACK: Normal completion [A]
When ST1 ≠ ACK: Abnormal termination [E]

4.14.3 Status at processing completion

Status at F	Processing Completion	Status Code	Description
Normal completion [A]	Normal acknowledgment (ACK)	06H	The command was executed normally and security setting data was set normally.
Abnormal termination [B]	Parameter error	05H	The command information (parameter) is not 00H or the parameter BOT is not 03H.
	Checksum error		The checksum of the transmitted command frame or data frame does not match.
	Protect error	10H	An already prohibited flag is to be enabled.
	Negative acknowledgment (NACK)		Command frame data is abnormal (such as invalid data length (LEN) or no ETX).
Time-out error [0	Time-out error [C]		The status frame or data frame was not received within the specified time.
Abnormal	MRG10 error	1AH	Writing security data has failed.
termination	MRG11 error	1BH	
[D], [E]	Write error	1CH	

4.14.4 Flowchart



4.14.5 Sample program

The following shows a sample program for Security Set command processing.

```
/**************************
/*
/*
                                             * /
       Set security flag command
                                             * /
/* [i] u8 scf
              ... Security flag data
                                             * /
/* [r] u16
              ... error code
u16 fl_ua_setscf(u8 scf, u8 bot)
{
   u16
      rc;
   /****************
     set params
   /***************/
                              // "BLK" (must be 0x00)
   fl\_cmd\_prm[0] = 0x00;
                               // "PAG" (must be 0x00)
   fl_cmd_prm[1] = 0x00;
   fl_txdata_frm[0] = scf|= 0b11101000;  // "FLG" (bit 7,6,5,3 must be '1')
   fl_txdata_frm[1] = bot;
                              // "BOT"
   fl_txdata_frm[2] = 0xff;
                              // (must be 0xff)
   fl_txdata_frm[3] = 0xff;
                               // (must be 0xff)
                              // (must be 0xff)
   fl_txdata_frm[4] = 0xff;
   fl_txdata_frm[5] = 0xff;
                              // (must be 0xff)
   fl_txdata_frm[6] = 0xff;
                              // (must be 0xff)
                               // (must be 0xff)
   fl_txdata_frm[7] = 0xff;
   /***************
      send command
   /****************
   fl_wait(tCOM);
                                    // wait before sending command
   put_cmd_ua(FL_COM_SET_SECURITY, 3, fl_cmd_prm);
   rc = get_sfrm_ua(fl_ua_sfrm, tWT13_T0);
                                        // get status frame
   switch(rc) {
       case FLC_NO_ERR:
                               break; // continue
      case FLC_DFTO_ERR: return rc; break; // case [C]
       default: return rc; break; // case [B]
   }
   /* send data frame (security setting data) */
   /****************
   fl_wait(tFD3);
   rc = get_sfrm_ua(fl_ua_sfrm, tWT14_MAX);
                                        // get status frame
   switch(rc) {
       case FLC_NO_ERR:
                               break; // continue
       case FLC_DFTO_ERR: return rc; break; // case [C]
```

```
default:
                     return rc; break; // case [B]
   }
   /* Check internally verify
   /***************
                                        // get status frame
   rc = get_sfrm_ua(fl_ua_sfrm, tWT15_MAX);
// switch(rc) {
//
//
      case FLC_NO_ERR: return rc; break; // case [A]
      case FLC_DFTO_ERR: return rc; break; // case [C]
//
//
      default: return rc; break; // case [B]
// }
  return rc;
}
```

CHAPTER 5 FLASH MEMORY PROGRAMMING PARAMETER CHARACTERISTICS

This chapter describes the characteristics of parameter transmitted between the programmer and the devices (78K0/Kx2-L and 78K0/Ix2) in the flash memory programming mode. Refer to the user's manual of the 78K0/Kx2-L or 78K0/Ix2 for electrical specifications when designing a programmer.

5.1 Flash Memory Parameter Characteristics of 78K0/Kx2-L

5.1.1 Flash memory programming mode setting time

Parameter	Symbol	MIN.	TYP.	MAX.
Ready start time from RESET1	t R0	1.5 ms		3 s
Wait for Reset command	toc	14.0 <i>μ</i> s		3 s

(1) Detailed mode specification time

Parameter	Symbol	MIN.	TYP.	MAX.
Release of POC to TOOLC↑	T1	1.0 <i>μ</i> s		
TOOLD↑ to TOOLC pulse start	Т3	110.0 <i>μ</i> s		
Minimum Low width of TOOLC	T4	110.0 <i>μ</i> s		
Minimum High width of TOOLC	T5	110.0 <i>μ</i> s		
Setup time of TOOLC to TOOLD for Mode Setting	Т6	55.0 <i>μ</i> s		
Minimum Low width of TOOLD	T7	110.0 <i>μ</i> s		
Minimum High width of TOOLD	Т8	110.0 <i>μ</i> s		
Setup time of TOOLD to TOOLC for Mode Setting	Т9	55.0 <i>μ</i> s		
Setup time of TOOLD to RESET↑ for Mode Setting	T10	55.0 <i>μ</i> s		
Ready start time from RESET	T11 (= t _{R0})	1.5 ms		3 s
Maximum Setup time	T12			130.0 ms

(2) Detailed mode transition time (transition from the programming mode to normal mode)

Parameter	Symbol	MIN.	TYP.	MAX.
RESET↓ to TOOLC, TOOLD↓	T20	500.0 <i>μ</i> s		
Minimum Low width TOOLC and TOOLD	T21	110.0 <i>μ</i> s		
Minimum High width TOOLC and TOOLD	T22	55.0 <i>μ</i> s		

5.1.2 Programming characteristics

Wait	Condition	Symbol	MIN.	MAX.
Between data frame transmissions	Data frame reception	t DR	21.0 <i>μ</i> s	
	Data frame transmission	tот	O ^{Note 1}	
From status frame transmission until data frame transmission (1)		tFD1 ^{Note 2}	O ^{Note 1}	
From status frame transmission until data frame transmission (2)		t _{FD2}	O ^{Note 1}	
From status frame transmission until data frame reception (3)		t _{FD3}	20.0 <i>μ</i> s	
From status frame transmission until command frame reception		tсом	30.8 μs	

Notes 1. Enable successive reception for the programmer.

2. Time for one block transmission

Remark The waits are defined as follows.

tdr, tfd3, tcom>

The 78K0/Kx2-L can execute the next communication after the MIN. time has elapsed after completion of the current communication.

The programmer needs to transmit the following data within the period from the MIN. time to the MAX. time after completion of the current communication.

The MAX. time is not specified, but use approximately 3 seconds.

<tot, tfd1, tfd2>

The 78K0/Kx2-L can execute the next communication after the MIN. time has elapsed after completion of the current communication.

The programmer needs to be ready for reception of the following data within the MIN. time after completion of the current communication.

The MAX. time is not specified, but set the timeout to approximately 3 seconds.

5.1.3 Command characteristics

Command	Symbol	Condition	MIN.	MAX.
Reset	twто	_	Note 1	
Chip Erase	twr1	_	(54.4 + 6.2 × total number of	(883.6 + 136.1 × total number of
			blocks) ms	blocks) ms
Block Erase	twr2 ^{Note 2}	_	15.2 ms	(0.3 + 190.3 × execution count of
				simultaneous selection and erasure
				+ 139.9 \times number of blocks to be
				erased) ms
Programming	twтз	_	Note 1	
	twt4 ^{Note 3}	_	7.4 ms	132.9 ms
	twT5 ^{Note 4}	_	Note 1	723.6 ms
Verify	twт6	_	Note 1	
	twr7 ^{Note 3}	_	Note 1	
Block Blank Check	twt8 ^{Note 4}	_	Note 1	9.9 ms
Baud Rate Set	twr10	_	62.7 <i>μ</i> s	
Silicon Signature	twr11	_	Note 1	
Version Get	t WT12	_	Note 1	
Security Set	t wT13	_	Note 1	
	twT14	_	16.9 ms	351.3 ms
	t wT15	_	Note 1	352.5 ms
Checksum	t wT16	_	Note 1	

Notes 1. Reception must be enabled for the programmer before command frame transmission.

- 2. See 5.3 Simultaneous Selection and Erasure Performed by Block Erase Command for the calculation method of the execution count of simultaneous selection and erasure.
- 3. Time for 256-byte data transmission
- 4. Time for one block transmission

Remark The waits are defined as follows.

<twто to twт16>

The 78K0/Kx2-L completes command processing between the MIN. and MAX. times.

The programmer must check the status of the start bit of the reception frame until the MAX. time has elapsed.

For commands without a specified MAX. time, set the time to approximately 3 seconds.

5.2 Flash Memory Parameter Characteristics of 78K0/lx2

5.2.1 Flash memory programming mode setting time

Parameter	Symbol	MIN.	TYP.	MAX.
Ready start time from RESET↑Note	t R0	1.5 ms		3 s
Wait for Reset command	toc	14.0 <i>μ</i> s		3 s

Note See the table below. The values in the table are reference values. The values depend on the used oscillator and might therefore exceed the maximum value or fall below the minimum value.

(1) Detailed mode specification time

Parameter	Symbol	MIN.	TYP.	MAX.
Release of POC to TOOLC↑	T1	1.0 <i>μ</i> s		
TOOLD↑ to TOOLC pulse start	Т3	110.0 <i>μ</i> s		
Minimum Low width of TOOLC	T4	110.0 <i>μ</i> s		
Minimum High width of TOOLC	T5	110.0 <i>μ</i> s		
Setup time of TOOLC to TOOLD for Mode Setting	T6	55.0 <i>μ</i> s		
Minimum Low width of TOOLD	T7	110.0 <i>μ</i> s		
Minimum High width of TOOLD	Т8	110.0 <i>μ</i> s		
Setup time of TOOLD to TOOLC for Mode Setting	Т9	55.0 <i>μ</i> s		
Setup time of TOOLD to RESET↑ for Mode Setting	T10	55.0 <i>μ</i> s		
Ready start time from RESET	T11 (= tro)	1.5 ms		3 s
Maximum Setup time (13-bit counter)	T12			130.0 ms ^{Note}

Note @30 kHz ±50%

(2) Detailed mode transition time (transition from the programming mode to normal mode)

Parameter	Symbol	MIN.	TYP.	MAX.
RESET↓ to TOOLC, TOOLD↓	T20	500.0 <i>μ</i> s		
Minimum Low width TOOLC and TOOLD	T21	110.0 <i>μ</i> s		
Minimum High width TOOLC and TOOLD	T22	55.0 <i>μ</i> s		

5.2.2 Programming characteristics

Wait	Condition	Symbol	MIN.	MAX.
Between data frame transmissions	Data frame reception	tor	21.0 <i>μ</i> s	
	Data frame transmission	t ot	O ^{Note 1}	
From status frame transmission until data frame transmission (1)		tFD1 ^{Note 2}	O ^{Note 1}	
From status frame transmission until data frame transmission (2)		t _{FD2}	O ^{Note 1}	
From status frame transmission until data frame reception (3)		tгдз	20.0 <i>μ</i> s	
From status frame transmission until command frame reception		tсом	30.8 μs	

Notes 1. Enable successive reception for the programmer.

2. Time for one block transmission

Remark The waits are defined as follows.

<tDR, tFD3, tCOM>

The 78K0/lx2 can execute the next communication after the MIN. time has elapsed after completion of the current communication.

The programmer needs to transmit the following data within the period from the MIN. time to the MAX. time after completion of the current communication.

The MAX. time is not specified, but use approximately 3 seconds.

<tot, tfd1, tfd2>

The 78K0/lx2 can execute the next communication after the MIN. time has elapsed after completion of the current communication.

The programmer needs to be ready for reception of the following data within the MIN. time after completion of the current communication.

The MAX. time is not specified, but set the timeout to approximately 3 seconds.

5.2.3 Command characteristics

Command	Symbol	Condition	MIN.	MAX.
Reset	twтo	_	Note 1	
Chip Erase	twT1	_	(54.4 + 6.2 × total number of	(883.6 + 136.1 × total number of
			blocks) ms	blocks) ms
Block Erase	twr2 ^{Note 2}	_	15.2 ms	(0.3 + 190.3 × execution count of
				simultaneous selection and erasure
				+ 139.9 × number of blocks to be
				erased) ms
Programming	twтз	_	Note 1	
	twt4 ^{Note 3}	_	7.4 ms	132.9 ms
	twr5 ^{Note 4}	_	Note 1	723.6 ms
Verify	t wт6	_	Note 1	
	twr7 ^{Note 3}	_	Note 1	
Block Blank Check	twt8 ^{Note 4}	_	Note 1	9.9 ms
Baud Rate Set	twT10	_	62.7 <i>μ</i> s	
Silicon Signature	twr11	_	Note 1	
Version Get	t WT12	_	Note 1	
Security Set	twT13	_	Note 1	
	twT14	_	16.9 ms	351.3 ms
	t wT15	_	Note 1	352.5 ms
Checksum	t wT16	_	Note 1	

Notes 1. Reception must be enabled for the programmer before command frame transmission.

- 2. See 5.3 Simultaneous Selection and Erasure Performed by Block Erase Command for the calculation method of the execution count of simultaneous selection and erasure.
- 3. Time for 256-byte data transmission
- 4. Time for one block transmission

Remark The waits are defined as follows.

<twro to twr16>

The 78K0/lx2 completes command processing between the MIN. and MAX. times.

The programmer must check the status of the start bit of the reception frame until the MAX. time has elapsed.

For commands without a specified MAX. time, set the time to approximately 3 seconds.

5.3 Simultaneous Selection and Erasure Performed by Block Erase Command

The Block Erase command of the 78K0/Kx2-L and 78K0/lx2 is executed by repeating "simultaneous selection and erasure", which erases multiple blocks simultaneously.

The wait time inserted during Block Erase command execution is therefore equal to the total execution time of "simultaneous selection and erasure".

To calculate the "total execution time of simultaneous selection and erasure", the execution count (M) of the simultaneous selection and erasure must first be calculated.

"M" is calculated by obtaining the number of blocks to be erased simultaneously (number of blocks to be selected and erased simultaneously).

The following describes the method for calculating the number of blocks to be selected and erased simultaneously and the execution count (M).

5.3.1 Calculation of number of blocks to be selected and erased simultaneously

The number of blocks to be selected and erased simultaneously should be 1, 2, 4, 8, 16, 32, 64, or 128, depending on which satisfies all of the following conditions.

[Condition 1]

(Number of blocks to be erased) ≥ (Number of blocks to be selected and erased simultaneously)

[Condition 2]

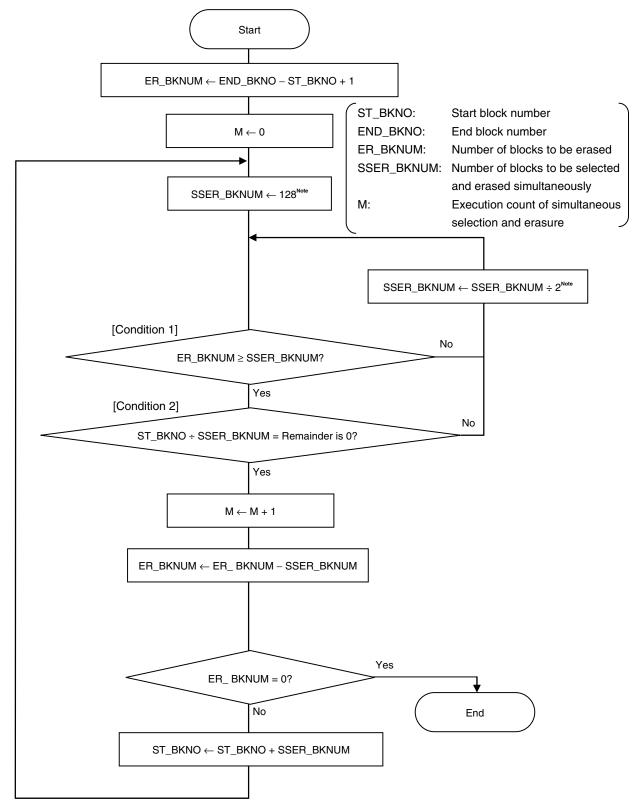
(Start block number) ÷ (Number of blocks to be selected and erased simultaneously) = Remainder is 0

[Condition 3]

The maximum value among the values that satisfy both Conditions 1 and 2

5.3.2 Calculation of execution count (M) of simultaneous selection and erasure

Calculation of the execution count (M) is illustrated in the following flowchart.



Note Based on the maximum value of SSER_BKNUM (128), obtain the value that satisfies Conditions 1 and 2 by executing SSER_BKNUM ÷ 2; Condition 3 is then satisfied.

Example 1 Erasing blocks 1 to 127 (N (number of blocks to be erased) = 127)

- <1> The first start block number is 1 and the number of blocks to be erased is 127; the values that satisfy Condition 1 are therefore 1, 2, 4, 8, 16, 32, and 64.

 Moreover, the value that satisfies Condition 2 is 1 and the value that satisfies Condition 3 is 1, so the number of blocks to be selected and erased simultaneously is 1; only block 1 is then erased.
- <2> After block 1 is erased, the next start block number is 2 and the number of blocks to be erased is 126; the values that satisfy Condition 1 are therefore 1, 2, 4, 8, 16, 32, and 64.
 Moreover, the values that satisfy Condition 2 are 1 and 2, the value that satisfies Condition 3 is 2, so the number of blocks to be selected and erased simultaneously is 2; blocks 2 and 3 are then erased.
- <3> After blocks 2 and 3 are erased, the next start block number is 4 and the number of blocks to be erased is 124; the values that satisfy Condition 1 are therefore 1, 2, 4, 8, 16, 32, and 64. Moreover, the values that satisfy Condition 2 are 1, 2, and 4, the value that satisfies Condition 3 is 4, so the number of blocks to be selected and erased simultaneously is 4; blocks 4 to 7 are then erased.
- <4> After blocks 4 to 7 are erased, the next start block number is 8 and the number of blocks to be erased is 120; the values that satisfy Condition 1 are therefore 1, 2, 4, 8, 16, 32, and 64. Moreover, the values that satisfy Condition 2 are 1, 2, 4, and 8, the value that satisfies Condition 3 is 8, so the number of blocks to be selected and erased simultaneously is 8; blocks 8 to 15 are then erased.
- <5> After blocks 8 to 15 are erased, the next start block number is 16 and the number of blocks to be erased is 112; the values that satisfy Condition 1 are therefore 1, 2, 4, 8, 16, 32, and 64. Moreover, the values that satisfy Condition 2 are 1, 2, 4, 8, and 16, the value that satisfies Condition 3 is 16, so the number of blocks to be selected and erased simultaneously is 16; blocks 16 to 31 are then erased.
- <6> After blocks 16 to 31 are erased, the next start block number is 32 and the number of blocks to be erased is 96; the values that satisfy Condition 1 are therefore 1, 2, 4, 8, 16, 32, and 64. Moreover, the values that satisfy Condition 2 are 1, 2, 4, 8, 16, and 32, the value that satisfies Condition 3 is 32, so the number of blocks to be selected and erased simultaneously is 32; blocks 32 to 63 are then erased.
- <7> After blocks 32 to 63 are erased, the next start block number is 64 and the number of blocks to be erased is 64; the values that satisfy Condition 1 are therefore 1, 2, 4, 8, 16, 32, and 64. Moreover, the values that satisfy Condition 2 are 1, 2, 4, 8, 16, 32, and 64, the value that satisfies Condition 3 is 64, so the number of blocks to be selected and erased simultaneously is 64; blocks 64 to 127 are then erased.

Therefore, simultaneous selection and erasure is executed seven times (1, 2 and 3, 4 to 7, 8 to 15, 16 to 31, 32 to 63, and 64 to 127) to erase blocks 1 to 127, so M = 7 is obtained.

Block configuration when executing simultaneous selection and erasure (when erasing blocks 1 to 127)

<Block number> 127 - <u>64</u> 63 User area --<u>16</u> 15

<Range of blocks that can be selected and erased simultaneously>

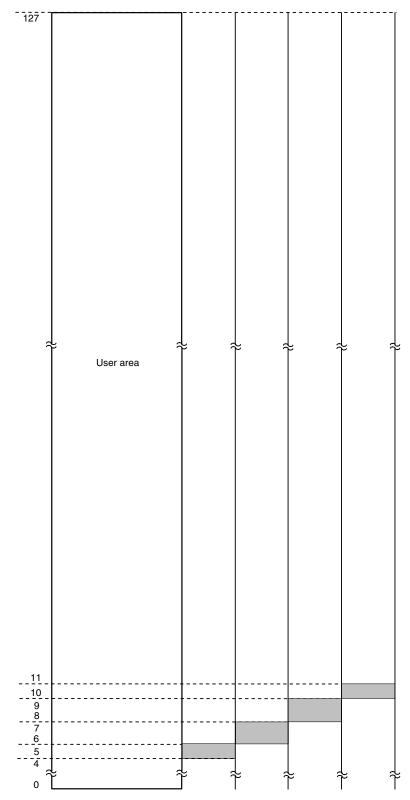
Example 2 Erasing blocks 5 to 10 (N (number of blocks to be erased) = 6)

- <1> The first start block number is 5 and the number of blocks to be erased is 6; the values that satisfy Condition 1 are therefore 1, 2, and 4.
 Moreover, the value that satisfies Condition 2 is 1 and the value that satisfies Condition 3 is 1, so the number of blocks to be selected and erased simultaneously is 1; only block 5 is the erased.
- <2> After block 5 is erased, the next start block number is 6 and the number of blocks to be erased is 5; the values that satisfy Condition 1 are therefore 1, 2, and 4.
 Moreover, the values that satisfy Condition 2 are 1 and 2, the value that satisfies Condition 3 is 2, so the number of blocks to be selected and erased simultaneously is 2; blocks 6 and 7 are then erased.
- <3> After blocks 6 and 7 are erased, the next start block number is 8 and the number of blocks to be erased is 3; the values that satisfy Condition 1 are therefore 1 and 2.
 Moreover, the values that satisfy Condition 2 are 1 and 2, the value that satisfies Condition 3 is 2, so the number of blocks to be selected and erased simultaneously is 2; blocks 8 and 9 are then erased.
- <4> After blocks 8 and 9 are erased, the next start block number is 10 and the number of blocks to be erased is 1; the value that satisfies Condition 1 is therefore 1. This also satisfies Conditions 2 and 3, so the number of blocks to be selected and erased simultaneously is 1; block 10 is then erased.

Therefore, simultaneous selection and erasure is executed four times (5, 6 and 7, 8 and 9, and 10) to erase blocks 5 to 10, so M = 4 is obtained.

Block configuration when executing simultaneous selection and erasure (when erasing blocks 5 to 10)





<Range of blocks that can be selected and erased simultaneously>

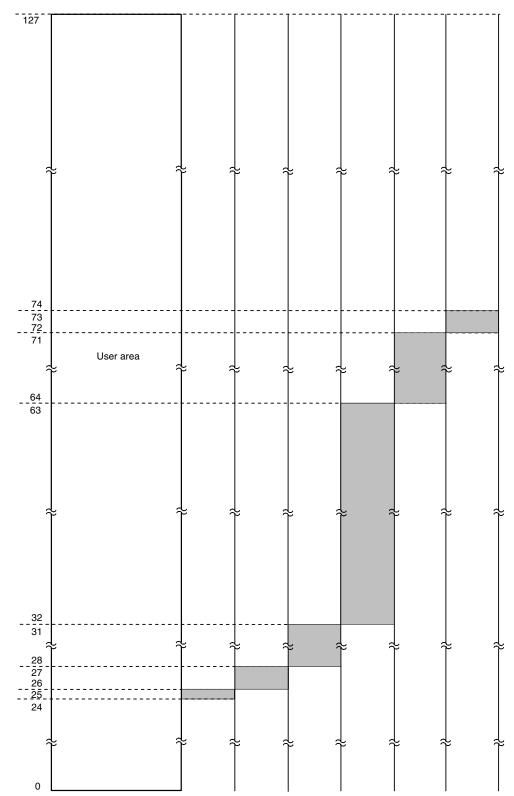
Example 3 Erasing blocks 25 to 73 (N (number of blocks to be erased) = 49)

- <1> The first start block number is 25 and the number of blocks to be erased is 49; the values that satisfy Condition 1 are therefore 1, 2, 4, 8, 16, and 32.
 Moreover, the value that satisfies Condition 2 is 1 and the value that satisfies Condition 3 is 1, so the number of blocks to be selected and erased simultaneously is 1; only block 25 is then erased.
- <2> After block 25 is erased, the next start block number is 26 and the number of blocks to be erased is 48; the values that satisfy Condition 1 are therefore 1, 2, 4, 8, 16, and 32. Moreover, the values that satisfy Condition 2 are 1 and 2, the value that satisfies Condition 3 is 2, so the number of blocks to be selected and erased simultaneously is 2; blocks 26 and 27 are then erased.
- <3> After blocks 26 and 27 are erased, the next start block number is 28 and the number of blocks to be erased is 46; the values that satisfy Condition 1 are therefore 1, 2, 4, 8, 16, and 32. Moreover, the values that satisfy Condition 2 are 1, 2, and 4, the value that satisfies Condition 3 is 4, so the number of blocks to be selected and erased simultaneously is 4; blocks 28 to 31 are then erased.
- <4> After blocks 28 to 31 are erased, the next start block number is 32 and the number of blocks to be erased is 42; the values that satisfy Condition 1 are therefore 1, 2, 4, 8, 16, and 32. Moreover, the values that satisfy Condition 2 are 1, 2, 4, 8, and 32, the value that satisfies Condition 3 is 32, so the number of blocks to be selected and erased simultaneously is 32; blocks 32 to 63 are then erased.
- <5> After blocks 32 to 63 are erased, the next start block number is 64, and the number of blocks to be erased is 10; the values that satisfy Condition 1 are therefore 1, 2, 4, and 8.
 Moreover, the values that satisfy Condition 2 are 1, 2, 4, and 8, the value that satisfies Condition 3 is 8, so the number of blocks to be selected and erased simultaneously is 8; blocks 64 to 71 are then erased.
- <6> After blocks 64 to 71 are erased, the next start block number is 72, and the number of blocks to be erased is 2; the values that satisfy Condition 1 are therefore 1 and 2.
 Moreover, the values that satisfy Condition 2 are 1 and 2, the value that satisfies Condition 3 is 2, so the number of blocks to be selected and erased simultaneously is 2; blocks 72 and 73 are then erased.

Therefore, simultaneous selection and erasure is executed six times (25, 26 and 27, 28 to 31, 32 to 63, 64 to 71, and 72 and 73) to erase blocks 25 to 73, so M = 6 is obtained.

Block configuration when executing simultaneous selection and erasure (when erasing blocks 25 to 73)



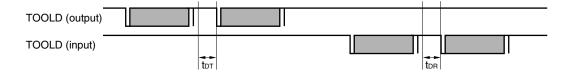


< Range of blocks that can be selected and erased simultaneously>

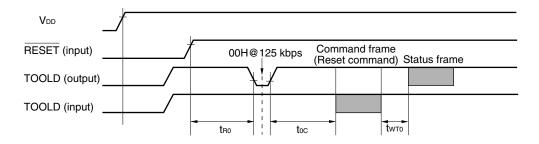
5.4 UART Communication Mode

In the figure below, TOOLD is illustrated as two separate lines for the sake of description, but it is actually a single line.

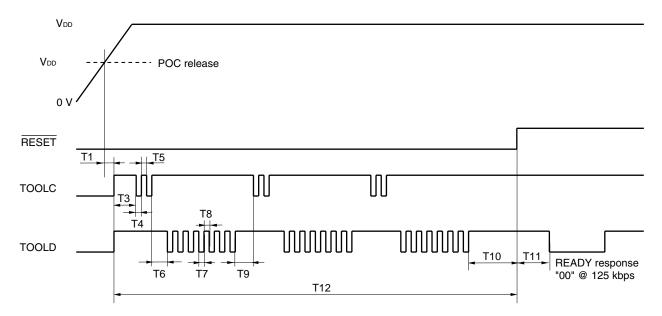
(1) Data frame



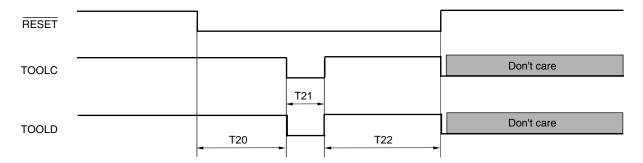
(2) Programming mode setting/Reset command



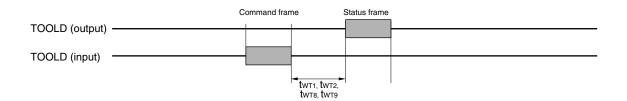
(a) Details of programming mode setting



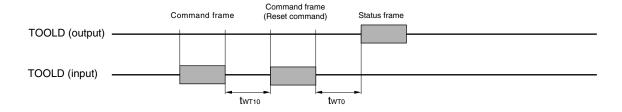
(3) Transition from programming mode to normal operating mode



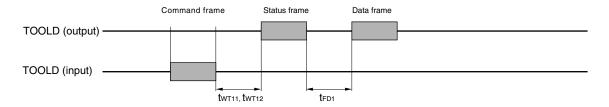
(4) Chip Erase command/Block Erase command/Block Blank Check command/Oscillating Frequency Set command



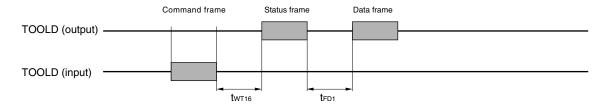
(5) Baud Rate Set command



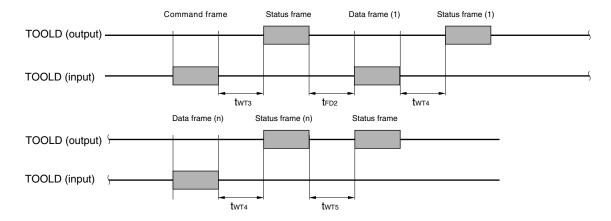
(6) Silicon Signature command/Version Get command



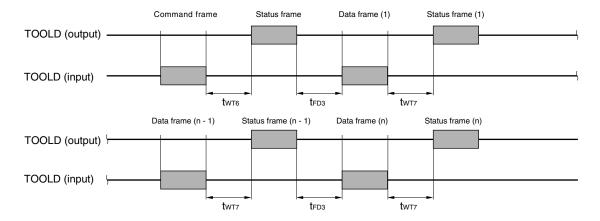
(7) Checksum command



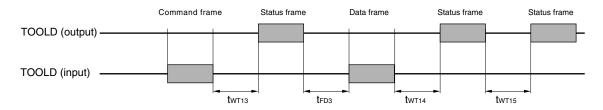
(8) Programming command



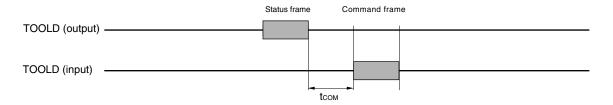
(9) Verify command



(10) Security Set command



(11) Wait before command frame transmission

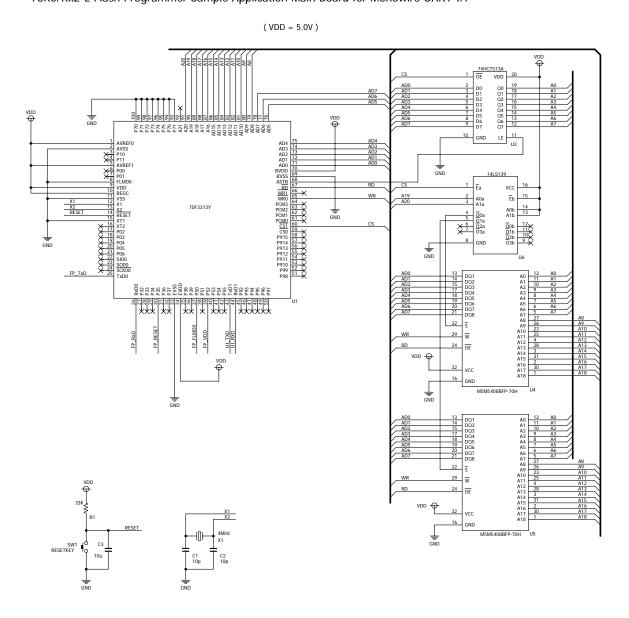


APPENDIX A CIRCUIT DIAGRAMS (REFERENCE)

Figures A-1 and A-2 show circuit diagrams of the programmer and the 78K0/Kx2-L, for reference.

Figure A-1. Reference Circuit Diagram of Programmer and 78K0/Kx2-L (Main Board)

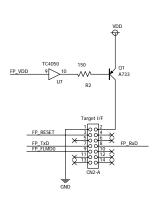
78KO/Kx2-L Flash Programmer Sample Application Main Board for Monowire UART I/F



78K0/Kx2-L Target Board for Monowire UART I/F

Remark For details about connecting unused pins shown in this circuit diagram, see the user's manual of each product.

122 Application Note U19735EJ1V0AN



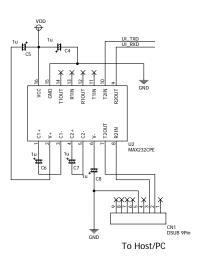
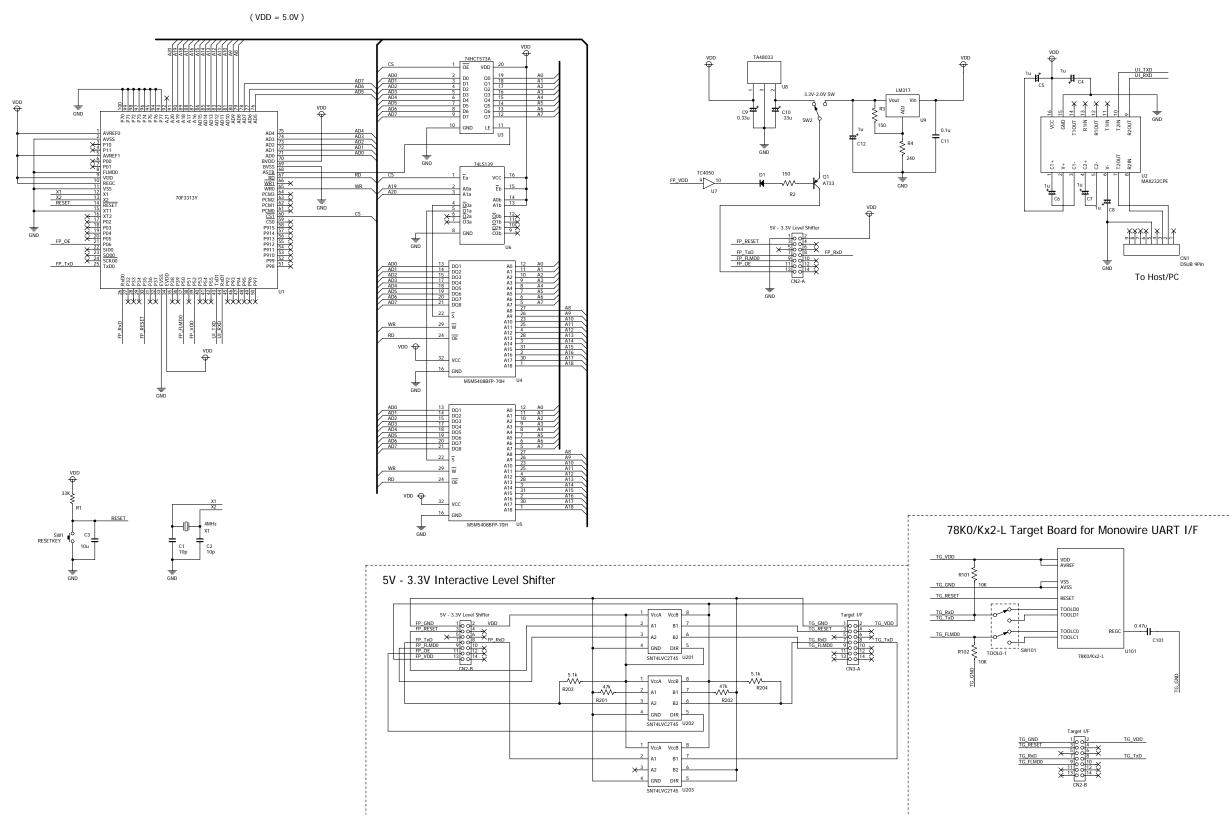


Figure A-2. Reference Circuit Diagram of Programmer and 78K0/Kx2-L (Main Board) (When Using 3.3 V Interactive Level Shifter)

78KO/Kx2-L Flash Programmer Sample Application Main Board for Monowire UART I/F



Remark For details about connecting unused pins shown in this circuit diagram, see the user's manual of each product.

124 Application Note U19735EJ1V0AN

For further information, please contact:

NEC Electronics Corporation

1753, Shimonumabe, Nakahara-ku, Kawasaki, Kanagawa 211-8668, Japan Tel: 044-435-5111 http://www.necel.com/

[America]

NEC Electronics America, Inc.

2880 Scott Blvd.
Santa Clara, CA 95050-2554, U.S.A.
Tel: 408-588-6000
800-366-9782
http://www.am.necel.com/

[Europe]

NEC Electronics (Europe) GmbH

Arcadiastrasse 10 40472 Düsseldorf, Germany Tel: 0211-65030 http://www.eu.necel.com/

Hanover Office

Podbielskistrasse 166 B 30177 Hannover Tel: 0 511 33 40 2-0

Munich Office

Werner-Eckert-Strasse 9 81829 München Tel: 0 89 92 10 03-0

Stuttgart Office

Industriestrasse 3 70565 Stuttgart Tel: 0 711 99 01 0-0

United Kingdom Branch

Cygnus House, Sunrise Parkway Linford Wood, Milton Keynes MK14 6NP, U.K. Tel: 01908-691-133

Succursale Française

9, rue Paul Dautier, B.P. 52 78142 Velizy-Villacoublay Cédex France

Tel: 01-3067-5800

Sucursal en España

Juan Esplandiu, 15 28007 Madrid, Spain Tel: 091-504-2787

Tyskland Filial

Täby Centrum
Entrance S (7th floor)
18322 Täby, Sweden
Tel: 08 638 72 00

Filiale Italiana

Via Fabio Filzi, 25/A 20124 Milano, Italy Tel: 02-667541

Branch The Netherlands

Steijgerweg 6 5616 HS Eindhoven The Netherlands Tel: 040 265 40 10

[Asia & Oceania]

NEC Electronics (China) Co., Ltd

7th Floor, Quantum Plaza, No. 27 ZhiChunLu Haidian District, Beijing 100083, P.R.China Tel: 010-8235-1155 http://www.cn.necel.com/

Shanghai Branch

Room 2509-2510, Bank of China Tower, 200 Yincheng Road Central, Pudong New Area, Shanghai, P.R.China P.C:200120 Tel:021-5888-5400 http://www.cn.necel.com/

Shenzhen Branch

Unit 01, 39/F, Excellence Times Square Building, No. 4068 Yi Tian Road, Futian District, Shenzhen, P.R.China P.C:518048 Tel:0755-8282-9800 http://www.cn.necel.com/

NEC Electronics Hong Kong Ltd.

Unit 1601-1613, 16/F., Tower 2, Grand Century Place, 193 Prince Edward Road West, Mongkok, Kowloon, Hong Kong Tel: 2886-9318 http://www.hk.necel.com/

NEC Electronics Taiwan Ltd.

7F, No. 363 Fu Shing North Road Taipei, Taiwan, R. O. C. Tel: 02-8175-9600 http://www.tw.necel.com/

NEC Electronics Singapore Pte. Ltd.

238A Thomson Road, #12-08 Novena Square, Singapore 307684 Tel: 6253-8311 http://www.sg.necel.com/

NEC Electronics Korea Ltd.

11F., Samik Lavied'or Bldg., 720-2, Yeoksam-Dong, Kangnam-Ku, Seoul, 135-080, Korea Tel: 02-558-3737 http://www.kr.necel.com/

G0706