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Application Note

CA850 Ver.2.50

C Compiler Package

Coding Technique

Target Devices V850 Series[™]

[MEMO]

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INTRODUCTION

Target Readers This manual is intended for users who wish to design and develop application systems using

the V850 Series.

Purpose This manual is intended to give users an understanding of the coding techniques for

increasing execution speed or reducing code size after specifying an optimization option

using the V850 Series C compiler CA850.

Organization This manual is divided into the following parts.

Overview

• Reducing code size

• Increasing execution speed

Defining variables

How to Read This Manual It is assumed that the reader of this manual has general knowledge in the fields of electrical engineering, logic circuits, microcontrollers, C language, and assembler programming.

To know the hardware functions of the V850 Series

ightarrow Refer to the hardware user's manual of each product

To know the instruction functions of the V850 Series

→ Refer to V850 Series Architecture User's Manual

Conventions Data significance: Higher digits on the left and lower digits on the right

Note: Footnote for item marked with **Note** in the text

Caution: Information requiring particular attention

Remark: Supplementary information

Numeric representation: Binary ... XXXX or XXXXB

Decimal ... XXXX

Hexadecimal ... XXXXH

Prefix indicating the power of 2 (address space, memory capacity):

K (kilo): $2^{10} = 1024$ M (mega): $2^{20} = 1024^2$

G (giga): $2^{30} = 1024^3$

Related Documents

Refer to the following manuals when using this manual.

The related documents indicated in this publication may include preliminary versions. However, preliminary versions are not marked as such.

Documents related to the V850 Series (user's manuals)

Docume	Document No.			
IE-703002-MC (In-Circuit Emulator for V853 [™] , V850/SV1 [™])	U11595E			
IE-V850E-MC (In-Circuit Emulator for V850E/IA1 TM , VIE-V850E-MC-A (In-Circuit Emulator for V850E/MA1 ^{TI}	/850E/IA2 TM) ^M , V850E/MA2 TM)		U14487E	
IE-703003-MC-EM1 (In-Circuit Emulator Option Board	d for V853)		U11596E	
IE-703017-MC-EM1 (In-Circuit Emulator Option Board	d for V850/SA1)		U12898E	
IE-703037-MC-EM1 (In-Circuit Emulator Option Board	d for V850/SB1, V850/SB2)		U14151E	
IE-703040-MC- EM1 (In-Circuit Emulator Option Boar	rd for V850/SV1)		U14337E	
IE-703079-MC- EM1 (In-Circuit Emulator Option Boar	d for V850/SF1)		U15447E	
IE-703102-MC (In-Circuit Emulator for V850E/MS1 [™])			U13875E	
IE-703102-MC-EM1, IE-703102-MC-EM1-A (In-Circui	t Emulator Option Board for	V850E/MS1)	U13876E	
IE-703107-MC-EM1 (In-Circuit Emulator Option Board	d for V850E/MA1)		U14481E	
IE-703116-MC-EM1 (In-Circuit Emulator Option Board	d for V850E/IA1)		U14700E	
CA850 Ver.2.50 C Compiler Package	Operation		U16053E	
	C Language		U16054E	
	PM plus		U16055E	
	Assembly Language	Assembly Language		
ID850 Ver.2.40 Integrated Debugger	Operation Windows TM B	ased	U15181E	
SM850 Ver.2.40 System Emulator	Operation Windows Bas	sed	U15182E	
SM850 Ver.2.00 or Later System Emulator	External Part User Oper	n Interface Specifications	U14873E	
RX850 Ver.3.13 or Later Real-time OS	Basics		U13430E	
	Installation		U13410E	
	Technical		U13431E	
RX850 Pro Ver.3.13 or Later Real-time OS	Basics		U13773E	
	Installation		U13774E	
	U13772E			
RD850 Ver.3.01 Task Debugger			U13737E	
RD850 Pro Ver.3.01 Task Debugger			U13916E	
AZ850 Ver.3.10 System Performance Analyzer	U14410E			
PG-FP4 Flash Memory Programmer			U15260E	
CA850 Ver.2.50 C Compiler Package (Application No	te)	Coding Technique	This manual	

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CHAPTER 1 OVERVIEW

This application note explains the coding techniques for further decreasing the code size or increasing the effective execution speed after an optimization option was specified using the CA850, which is the C compiler for the V850 Series.

For details concerning the CA850, refer to CA850-related user's manuals.

The amount that the coding size is reduced for the examples in each section is specific to those examples. The amount that the coding size will be reduced when that technique is applied elsewhere will differ somewhat depending on the individual case.

Also, the following points should be noted when modifying source code.

Since the register usage conditions will vary depending on the source code modification, register transfers that remain without having been eliminated by the optimization up to that point may be removed or, conversely, optimization that was performed may no longer be effective and lengthy register transfers may remain.

By adding temporary variables, registers for the new register variables may be used and this may cause coding for saving/restoring those registers to be added at function entries/exits. In this case, the code size will be increased by the amount of code required for saving/restoring the registers.

The output assembly list for the examples that appear in this application note shows assembler source code that was compiled by specifying size-priority optimization (-Os). Note that the results will differ when optimization other than size priority optimization is specified.

CHAPTER 2 REDUCING CODE SIZE

This chapter introduces coding techniques for reducing the code size.

Note that reducing the code size may also increase the effective execution speed.

2.1 Using the if~else Statement Instead of the switch Statement

If the following two conditions are satisfied simultaneously, the CA850 generates table branching format code for the switch statement.

- The number of case labels is at least four.
- The difference between the upper and lower limits of the label value is at most three times the number of case labels.

In this case, if the number of case labels is roughly 16 or less (however, this number differs according to the format of the switch expression or the distribution of label values), switching to the equivalent if~else statements and arranging a sequence of compare and branch instructions will produce a smaller code size.

If the switch expression is an external variable reference or complicated expression, you must assign a value to a temporary variable and change to a reference of that temporary variable in the if expression.

Note, however, that since the switch instruction is output for the V850E, the switch statement will produce a smaller code size.

Caution Even if the source description is not change the expansion code of a switch statement can be specified per file using the -Xcase option.

A sample program is shown below.

Remark x is assumed to be an auto variable.

Before Modification	After Modification
int x;	int x;
switch(x) {	if $(x == 0)$
case 0:	return(0);
return(0);	else if (x == 1)
case 1:	return(1);
return(1);	else if $(x == 2)$
case 2:	return(2);
return(2);	else if $(x == 3)$
case 3:	return(3);
return(3);	else if $(x == 4)$
case 4:	return(4);
return(4);	else if $(x == 5)$
case 5:	return(5);
return(5);	
}	

	Before Modification		After Modification			
	Program	Size [bytes]			Program	Size [bytes]
ld.w	-8+0x8[sp], r10	4		ld.w	-8+0x8[sp], r10	4
cmp	5, r10	2		cmp	r0, r10	2
jh	.L1	2		jne	.L2	2
shl	1, r10	2		st.w	r0, -4+0x8[sp]	4
add	tp, r10	2		jbr	.L1	2
ld.h	.L10[r10], r11	8	.L2:			
add	.L10, r11	10		cmp	1, r10	2
add	tp, r11	2		jne	.L4	2
jmp	[r11]	2		mov	1, r11	2
.L10:				st.w	r11, -4+0x8[sp]	4
.hword	L .L4L10	2		jbr	.L1	2
.hword	L.L5L10	2	.L4:			
.hword	L .L6L10	2		cmp	2, r10	2
.hword	L .L7L10	2		jne	.L6	2
.hword	L .L8L10	2		mov	2, r12	2
.hword	L .L9L10	2		st.w	r12, -4+0x8[sp]	4
.L4:				jbr	.L1	2
st.w	r0, -4+0x8[sp]	4	.L6:	_		
jbr	.L1	2		cmp	3, r10	2
.L5:				jne	.L8	2
mov	1, r12	2		mov	3, r13	2
st.w	r12, -4+0x8[sp]	4		st.w	r13, -4+0x8[sp]	4
jbr	.L1	2		jbr	.L1	2
.L6:		_	.L8:	J~2	.22	-
mov	2, r13	2		cmp	4, r10	2
st.w	r13, -4+0x8[sp]	4		jne	.L10	2
jbr	.L1	2		mov	4, r14	2
.L7:	. 111			st.w	r14, -4+0x8[sp]	4
.L/:	3, r14	2		jbr	.L1	2
st.w	r14, -4+0x8[sp]	4	.L10:	-	• = =	
	· • •	2			5, r10	2
jbr	.L1	2		cmp	.L1	2
.L8:	415	2		jne		2
mov	4, r15			mov	5, r15	
st.w	r15, -4+0x8[sp]	4		st.w	r15, -4+0x8[sp]	4
jbr	.L1	2	.L1:	1.1	4 . 0 0	
.L9:				ld.w	-4+0x8[sp], r10	4
mov	5, r16	2				
st.w	r16, -4+0x8[sp]	4				
.L1:	4 0 0 0 1					
ld.w	-4+0x8[sp], r10	4				
Total code size		94 bytes	Total cod	de size		76 byte

	Before Modification		After Modification		
	Program	Size [bytes]		Program	Size [bytes]
ld.w	-8+0x8[sp], r10	4	ld.w	-8+0x8[sp], r10	4
cmp	5, r10	2	cmp	r0, r10	2
jh	.L1	2	jne	.L2	2
switc	h r10	2	st.w	r0, -4+0x8[sp]	4
.L10:			jbr	.L1	2
.shwo	rd .L4L10	2	.L2:		
.shwo	rd .L5L10	2	cmp	1, r10	2
.shwo	rd .L6L10	2	jne	.L4	2
.shwo	rd .L7L10	2	mov	1, r11	2
.shwo	rd .L8L10	2	st.w	r11, -4+0x8[sp]	4
.shwo	rd .L9L10	2	jbr	.L1	2
.L4:			.L4:		
st.w	r0, -4+0x8[sp]	4	cmp	2, r10	2
jbr	.L1	2	jne	.L6	2
.L5:			mov	2, r12	2
mov	1, r12	2	st.w	r12, -4+0x8[sp]	4
st.w	r12, -4+0x8[sp]	4	jbr	.L1	2
jbr	.L1	2	.L6:		
.L6:			cmp	3, r10	2
mov	2, r13	2	jne	.L8	2
st.w	r13, -4+0x8[sp]	4	mov	3, r13	2
jbr	.L1	2	st.w	r13, -4+0x8[sp]	4
.L7:		_	jbr	.L1	2
mov	3, r14	2	.L8:		
st.w	r14, -4+0x8[sp]	4	cmp	4, r10	2
jbr	.L1	2	jne	.L10	2
.L8:		_	mov	4, r14	2
mov	4, r15	2	st.w	r14, -4+0x8[sp]	4
st.w	r15, -4+0x8[sp]	4	jbr	.L1	2
jbr	.L1	2	.L10:	· 	-
.L9:			cmp	5, r10	2
mov	5, r16	2	jne	.L1	2
st.w	r16, -4+0x8[sp]	4	mov		2
.L1:	,		st.w	r15, -4+0x8[sp]	4
ld.w	-4+0x8[sp], r10	4	.L1:	, -: 3M0 [SP]	'
14.W	1.010 [02] / 110	'	ld.w	-4+0x8[sp], r10	4
Total code size		70 bytes	Total code size		76 bytes
iotal code size		70 bytes	TOTAL CODE SIZE		76 bytes

2.2 Making Assignments to the Same External Variable Via a Temporary Variable at the Branch Destinations of the switch Statement or if~else Statement

To assign a different value to the same external variable at each branch destination of the switch statement or if~else statement, the code size can be reduced by assigning a temporary variable at each location, and then assigning the value to the original external variable from the temporary variable after the control logic has merged together again.

The code size is reduced for the following reason. Since external variables are rarely allocated to registers, an assignment to an external variable becomes a store instruction to memory (4 bytes). However, an assignment to a temporary variable often becomes a register transfer (2 bytes).

A sample program is shown below.

Remark s is assumed to be an external variable.

Before Modification	After Modification
int x;	int x;
	int tmp;
switch (x) {	
case 0:	if (x == 0) {
s = 0;	tmp = 0;
break;	}
case 1000:	else if $(x == 1000)$ {
s = 0x5555;	tmp = 0x5555;
break;	}
case 2000:	else if $(x == 2000)$ {
s = 0xAAAA;	tmp = 0xAAAA;
break;	}
case 3000:	else if $(x == 3000)$ {
s = 0xFFFF;	tmp = 0xFFFF;
}	}
	else {
	goto label;
	}
	s = tmp;
	label:
	i

	Before Modification			After Modification	
	Program	Size [bytes]		Program	Size [bytes]
ld.w	-4+0x4[sp], r10	4	ld.w	-4+0x4[sp], r10	4
cmp	r0, r10	2	cmp	r0, r10	2
je	. L4	2	jne	.L2	2
addi	-1000, r10, r0	4	mov	r0, r10	2
je	.L5	2	jbr	.L3	2
addi	-2000, r10, r0	4	.L2:		
ie	.L6	2	addi	-1000, r10, r0	4
addi	-3000, r10, r0	4	jne	. L4	2
je	.L7	2	mov	21845, r10	4
jbr	.L3	2	jbr	.L3	2
.L4:			.L4:		
st.w	r0, \$ s	4	addi	-2000, r10, r0	4
ibr	.L3	2	jne	.L6	
.L5:			mov	43690, r10	2 8
mov	21845, r11	4	jbr	.L3	2
st.w	r11, \$_s	4	.L6:		
ibr	.L3	2	addi	-3000, r10, r0	4
.L6:		_	jne	.L10	
mov	43690, r12	8	mov	65535, r10	2 8
st.w	r12, \$_s	4	.L3:	,	
ibr	.L3	2	st.w	r10, \$ s	4
.L7:		_	.L10:	· · · · = ·	
mov	65535, r13	8			
st.w	·	4			
.L3:	, 1	-			
Total code size	;	70 bytes	Total code size		58 bytes

Before Modification				After Modification	
	Program	Size [bytes]		Program	Size [bytes]
ld.w	-4+0x4[sp], r10	4	ld.w	-4+0x4[sp], r10	4
cmp	r0, r10	2	cmp	r0, r10	2
jе	.L4	2	jne	.L2	2
addi	-1000, r10, r0	4	mov	r0, r10	2
je	.L5	2	jbr	.L3	2
addi	-2000, r10, r0	4	.L2:		
jе	.L6	2	addi	-1000, r10, r0	4
addi	-3000, r10, r0	4	jne	.L4	2
jе	.L7	2	mov	21845, r10	4
jbr	.L3	2	jbr	.L3	2
.L4:			.L4:		
st.w	r0, \$_s	4	addi	-2000, r10, r0	4
jbr	.L3	2	jne	.L6	2
.L5:			mov	43690, r10	6
mov	21845, r11	4	jbr	.L3	2
st.w	r11, \$_s	4	.L6:		
jbr	.L3	2	addi	-3000, r10, r0	4
.L6:			jne	.L10	2
mov	43690, r12	6	mov	65535, r10	6
st.w	r12, \$ s	4	.L3:		
jbr		2	st.w	r10, \$ s	4
.L7:			.L10:	_	
mov	65535, r13	6			
st.w	r13, \$_s	4			
.L3:					
Total code size		66 bytes	Total code size		54 bytes

2.3 Moving One Assignment Statement in Front of the if Statement When the Branch Destinations of the if~else Statement Are Only Assignment Statements to the Same Variable

When the branch destinations of the if~else statement are only statements that assign different values to the same variable, the code size can be reduced by moving one of the assignment statements in front of the if statement and removing the else block, which eliminates the jump statement to the location following the else block from the if block.

A sample program is shown below.

Remark s is assumed to be an external variable.

Before Modification	After Modification
int x;	int x;
<pre>if (x == 10) { s = 1; } else { s = 0; }</pre>	<pre>s = 0; if (x == 10) { s = 1; }</pre>

[Output Assembly List for the V850]

	Before Modification		After Modification		
	Program	Size [bytes]		Program	Size [bytes]
ld.w	-4+0x4[sp], r10	4	st.w	r0, \$_s	4
cmp	10, r10	2	ld.w	-4+0x4[sp], r10	4
jne	.L2	2	cmp	10, r10	2
mov	1, r11	2	jne	.L2	2
st.w	r11, \$_s	4	mov	1, r11	2
jbr	.L3	2	st.w	r11, \$_s	4
.L2:			.L2:		
st.w	r0, \$_s	4			
.L3:					
Total code size		20 bytes	Total code size		18 bytes

Before Modification		After Modification		
	Program	Size [bytes]	Program	Size [bytes]
ld.w	-4+0x4[sp], r10	4	st.w r0, \$_s	4
cmp	10, r10	2	ld.w -4+0x4[sp], r10	4
jne	.L2	2	cmp 10, r10	2
mov	1, r11	2	jne .L2	2
st.w	r11, \$_s	4	mov 1, r11	2
jbr	.L3	2	st.w r11, \$_s	4
.L2:			.L2:	
st.w	r0, \$_s	4		
.L3:				
Total code size		20 bytes	Total code size	18 bytes

2.4 Replacing an Access to an External Variable with an Access to a Temporary Variable

Since an external variable access requires 4 bytes for both loading and storing, if the value of an external variable is assigned to a temporary variable and that temporary variable is used even in cases other than assignments like those described in Section 2.2, the code size can be reduced because the memory access is changed to a register access.

A sample program is shown below.

Remark s is assumed to be an external variable.

Before Modification	After Modification
int x;	int x;
	unsigned int tmp = s;
if (x != 0) {	
if ((s & 0x00F00F00) != 0x00E00E00) {	if (x != 0) {
return;	if ((tmp & 0x00F00F00) != 0x00E00E00) {
}	return;
s >>= 12;	}
s &= 0xFF;	tmp >>= 12;
}	tmp &= 0xFF;
else {	}
if ((s & 0x00FF0000) != 0x00EE0000) {	else {
return;	if ((tmp & 0x00FF0000) != 0x00EE0000) {
}	return;
s >>= 24;	}
}	tmp >>= 24;
	}
	s = tmp;

	Before Modification				After Modification	
	Program	Size [bytes]			Program	Size [bytes]
ld.w	-4+0x4[sp], r11	4		ld.w	\$_s, r10	4
cmp	r0, r11	2		ld.w	-4+0x4[sp], r11	4
je	.L2	2		cmp	r0, r11	2
ld.w	\$_s, r10	4		je	.L2	2
andi	0xf00f00, r10, r12	10		andi	0xf00f00, r10, r12	10
cmp	14683648, r12	10		cmp	14683648, r12	10
jne	.L1	2		jne	.L1	2
shr	12, r10	2		shr	12, r10	2
andi	0xff, r10, r13	4		and	0xff, r10	4
st.w	r13, \$_s	4		jbr	.L4	2
jbr	.L1	2	.L2:			
.L2:				andi	0xff0000, r10, r13	6
ld.w	\$_s, r10	4		cmp	15597568, r13	6
andi	0xff0000, r10, r14	6		jne	.L1	2
cmp	15597568, r14	6		shr	24, r10	2
jne	.L1	2	.L4:			
sar	24, r10	2		st.w	r10, \$_s	4
st.w	r10, \$_s	4	.L1:			
.L1:						
Total code size		70 bytes	Total cod	de size		62 bytes

	Before Modification		After Modification		
	Program	Size [bytes]		Program	Size [bytes]
ld.w	-4+0x4[sp], r11	4	ld.w	\$_s, r10	4
cmp	r0, r11	2	ld.w	-4+0x4[sp], r11	4
je	.L2	2	cmp	r0, r11	2
ld.w	\$_s, r10	4	je	.L2	2
andi	0xf00f00, r10, r12	8	andi	0xf00f00, r10, r12	8
cmp	14683648, r12	8	cmp	14683648, r12	8
jne	.L1	2	jne	.L1	2
shr	12, r10	2	shr	12, r10	2
zxb	r10	2	zxb	r10	2
st.w	r10, \$_s	4	jbr	.L4	2
jbr	.L1	2	.L2:		
.L2:			andi	0xff0000, r10, r13	6
ld.w	\$_s, r10	4	cmp	15597568, r13	6
andi	0xff0000, r10, r14	6	jne	.L1	2
cmp	15597568, r14	6	shr	24, r10	2
jne	.L1	2	.L4:		
sar	24, r10	2	st.w	r10, \$_s	4
st.w	r10, \$_s	4	.L1:		
.L1:					
Total code size		64 bytes	Total code size		56 bytes

2.5 Moving the Same Statement From After the Branch Clauses to Before the Branching Begins

If the same assignment statement or function call exists after each branch clause, move it in front of where the branching begins if it can be moved there.

If the evaluation result of that statement is referenced, assign it to a temporary variable and reference that temporary variable.

A sample program is shown below.

Remark s is assumed to be an external variable.

Before Modification	After Modification
int x;	int x;
	int tmp;
if $(x >= 0)$ {	
if (x > func(0, 1, 2)) {	tmp = func(0, 1, 2);
S++;	if $(x >= 0)$ {
}	if $(x > tmp)$ {
}	S++;
else {	}
if $(x < -func(0, 1, 2))$ {	}
s;	else {
}	if $(x < -tmp)$ {
}	s;
	}
	}

	Before Modification		After Modification			
	Program	Size [bytes]			Program	Size [bytes]
ld.w	-4+0x4[sp], r29	4	mo	ov	1, r7	2
cmp	r0, r29	2	mo	ov	2, r8	2
jlt	.L12	2	mo	ov	r0, r6	2
mov	1, r7	2	ja	arl	_func, lp	4
mov	2, r8	2	10	d.w	-4+0x4[sp], r11	4
mov	r0, r6	2	cr	mp	r0, r11	2
jarl	_func, lp	4	mo	ov	r10, r12	2
cmp	r10, r29	2	j:	lt	.L2	2
jle	.L4	2	cr	mp	r12, r11	2
ld.w	\$_s, r10	4	j:	le	.L4	2
add	1, r10	2	10	d.w	\$_s, r10	4
st.w	r10, \$_s	4	ad	dd	1, r10	2
jbr	L4	2	st	t.w	r10, \$_s	4
.L12:			jŀ	br	L4	2
mov	1, r7	2	.L2:			
mov	2, r8	2	no	ot	r12, r13	2
mov	r0, r6	2	a	dd	1, r13	2
jarl	_func, lp	4	CT	mp	r11, r13	2
not	r10, r12	2	j:	le	.L4	2
add	1, r12	2	10	d.w	\$ s, r14	4
cmp	r29, r12	2	ao	dd	- 4294967295, r14	2
jle	.L4	2	st	t.w	r14, \$_s	4
ld.w	\$_s, r13	4	.L4:		_	
add	4294967295, r13	2				
st.w	r13, \$ s	4				
.L4:						
Total code size		62 bytes	Total code	size		54 bytes

	Before Modification			After Modification	
	Program	Size [bytes]		Program	Size [bytes]
ld.w	-4+0x4[sp], r29	4	mov	1, r7	2
cmp	r0, r29	2	mov	2, r8	2
jlt	.L12	2	mov	r0, r6	2
mov	1, r7	2	jarl	_func, lp	4
mov	2, r8	2	ld.w	-4+0x4[sp], r11	4
mov	r0, r6	2	cmp	r0, r11	2
jarl	_func, lp	4	mov	r10, r12	2
cmp	r10, r29	2	jlt	.L2	2
jle	.L4	2	cmp	r12, r11	2
ld.w	\$_s, r10	4	jle	.L4	2
add	1, r10	2	ld.w	\$_s, r10	4
st.w	r10, \$_s	4	add	1, r10	2
jbr	L4	2	st.w	r10, \$ s	4
.L12:			jbr	_ .L4	2
mov	1, r7	2	.L2:		
mov	2, r8	2	not	r12, r13	2
mov	r0, r6	2	add	1, r13	2
jarl	_func, lp	4	cmp	r11, r13	2
not	r10, r12	2	jle	.L4	2
add	1, r12	2	ld.w	\$ s, r14	4
cmp	r29, r12	2	add	- 4294967295, r14	2
jle	.L4	2	st.w	r14, \$ s	4
ld.w	\$ s, r13	4	.L4:	_	
add	4294967295, r13	2			
st.w	·	4			
.L4:	· <u> </u>				
Total code size		62 bytes	Total code size		54 bytes

2.6 Moving the Same Statement From Before the Control Logic Merges Together to After It Merges Together

If the same assignment statement or function call exists after each branch clause and it cannot be moved in front of where the branching begins as described in Section 2.5, then move it after the control logic merges together if it can be moved there.

A sample program is shown below.

Remark s and t are assumed to be external variables.

Before Modification	After Modification
int tmp;	int tmp;
<pre>if (tmp & 0xff00ff00) { t++; s++; } else { t; s++; }</pre>	<pre>if (tmp & 0xff00ff00) { t++; } else { t; } s++;</pre>

	Before Modification			After Modification	
	Program	Size [bytes]		Program	Size [bytes]
ld.w	\$_s, r10	4	ld.w	-4+0x4[sp], r10	4
add	1, r10	2	and	0xff00ff00, r10	10
ld.w	-4+0x4[sp], r12	4	je	.L2	2
and	0xff00ff00, r12	10	ld.w	\$_t, r11	4
je	.L2	2	add	1, r11	2
ld.w	\$_t, r13	4	st.w	r11, \$_t	4
add	1, r13	2	jbr	.L3	2
st.w	r13, \$_t	4	.L2:		
st.w	r10, \$_s	4	ld.w	\$_t, r12	4
jbr	.L3	2	add	4294967295, r12	2
.L2:			st.w	r12, \$_t	4
ld.w	\$_t, r14	4	.L3:		
add	4294967295, r14	2	ld.w	\$_s, r13	4
st.w	r14, \$_t	4	add	1, r13	2
st.w	r10, \$_s	4	st.w	r13, \$_s	4
.L3:	_			_	
Total code size		52 bytes	Total code size		48 bytes

	Before Modification		After Modification		
	Program	Size [bytes]		Program	Size [bytes]
ld.w	\$_s, r10	4	ld.w	-4+0x4[sp], r10	4
add	1, r10	2	and	0xff00ff00, r10	8
ld.w	-4+0x4[sp], r12	4	je	.L2	2
and	0xff00ff00, r12	8	ld.w	\$_t, r11	4
je	.L2	2	add	1, r11	2
ld.w	\$_t, r13	4	st.w	r11, \$_t	4
add	1, r13	2	jbr	.L3	2
st.w	r13, \$_t	4	.L2:		
st.w	r10, \$_s	4	ld.w	\$_t, r12	4
jbr	.L3	2	add	4294967295, r12	2
.L2:			st.w	r12, \$_t	4
ld.w	\$ t, r14	4	.L3:	_	
add	4294967295, r14	2	ld.w	\$_s, r13	4
st.w	r14, \$_t	4	add	1, r13	2
st.w	r10, \$_s	4	st.w	r13, \$_s	4
.L3:					
Total code size		50 bytes	Total code size		46 bytes

2.7 Using a Temporary Variable to Consolidate the Calls to the Same Function with Different Arguments That Appear After Each Branch Clause at the Location After the Control Logic Merges Together

If a different argument is used when calling the same function after each branch clause, move the function call after the location where the control logic merges together again if it can be moved there. At this time, at each call location, assign the argument that differs to a temporary variable and use that temporary variable as the argument in the function call.

A sample program is shown below.

Remark s is assumed to be an external variable.

Before Modification	After Modification
if (s) {	int tmp;
func(0, 1, 2);	
}	if (s) {
else {	tmp = 2;
func(0, 1, 3);	}
}	else {
	tmp = 3;
	}
	func(0, 1, tmp);

	Before Modification		After Modification		
	Program	Size [bytes]		Program	Size [bytes]
ld.w	\$_s, r10	4	ld.w	\$_s, r10	4
cmp	r0, r10	2	cmp	r0, r10	2
je	.L9	2	je	.L2	2
mov	1, r7	2	mov	2, r11	2
mov	2, r8	2	jbr	.L7	2
mov	r0, r6	2	.L2:		
jarl	_func, lp	4	mov	3, r11	2
jbr	.L3	2	.L7:		
.L9:			mov	1, r7	2
mov	1, r7	2	mov	r0, r6	2
mov	3, r8	2	mov	r11, r8	2
mov	r0, r6	2	jarl	_func, lp	4
jarl	_func, lp	4			
.L3:					
Total code size		30 bytes	Total code size		24 bytes

	Before Modification			After Modificati	ion
	Program	Size [bytes]		Program	Size [bytes]
ld.w	\$_s, r10	4	ld.	w \$_s, r10	4
cmp	r0, r10	2	cmp	r0, r10	2
je	.L9	2	je	.L2	2
mov	1, r7	2	mov	2, r11	2
mov	2, r8	2	jbr	.L7	2
mov	r0, r6	2	.L2:		
jarl	_func, lp	4	mov	7 3, r11	2
jbr	.L3	2	.L7:		
.L9:			mov	7 1, r7	2
mov	1, r7	2	mov	r0, r6	2
mov	3, r8	2	mov	r11, r8	2
mov	r0, r6	2	jar	cl _func, lp	4
jarl	_func, lp	4			
.L3:					
Total code size		30 bytes	Total code s	ize	24 bytes

2.8 Replacing a Complex if Statement with One That Is Logically Equivalent

When the same processing is executed according to multiple cases in a combination of if~else statements, if the "multiple cases" can be consolidated by using a separate condition, consolidate them and remove the excess portions.

A sample program is shown below.

Remark When the conditions that the initial value of x is 0 and the values of s and t are either 0 or 1 are consolidated, they can be transformed as shown in this example. s, t, u, and v are assumed to be external variables.

Before Modification	After Modification
int x;	int x;
<pre>int x; if (!s) { if (t) { x = 1; } } else { if (!t) { x = 1; } } if (x) { if ((++u) >= v) { u = 0; }</pre>	<pre>int x; if (! (s ^ t)) { if ((++u) >= v) { u = 0; x = 1; } }</pre>
else {	
x = 0;	
}	

	Before Modification		After Modification		
	Program	Size [bytes]	Program	Size [bytes]	
ld.w	\$_s, r11	4	ld.w \$_s, r11	4	
cmp	r0, r11	2	ld.w \$_t, r12	4	
jne	.L2	2	xor r12, r11	2	
ld.w	\$_t, r12	4	jne .L2	2	
cmp	r0, r12	2	ld.w \$_u, r10	4	
je	.L4	2	add 1, r10	2	
mov	1, r13	2	st.w r10, \$_u	4	
st.w	r13, -4+0x4[sp]	4	ld.w \$_v, r14	4	
jbr	.L4	2	cmp r14, r10	2	
.L2:			jlt .L2	2	
ld.w	\$_t, r14	4	st.w r0, \$_u	4	
cmp	r0, r14	2	.L2:		
jne	.L4	2			
mov	1, r15	2			
st.w	r15, -4+0x4[sp]	4			
.L4:					
ld.w	-4+0x4[sp], r16	4			
cmp	r0, r16	2			
je	.L6	2			
ld.w	\$_u, r10	4			
add	1, r10	2			
ld.w	\$ v, r18	4			
cmp	r18, r10	2			
jlt	.L7	2			
st.w	r0, \$ u	4			
jbr	_ .L6	2			
.L7:					
st.w	r10, \$_u	4			
.L6:	_				
Total code size		70 bytes	Total code size	34 bytes	

	Before Modification		After Modification	
	Program	Size [bytes]	Program	Size [bytes]
ld.w	\$_s, r11	4	ld.w \$_s, r11	4
cmp	r0, r11	2	ld.w	4
jne	.L2	2	xor r12, r11	2
ld.w	\$_t, r12	4	jne .L2	2
cmp	r0, r12	2	ld.w \$_u, r10	4
je	.L4	2	add 1, r10	2
mov	1, r13	2	st.w r10, \$_u	4
st.w	r13, -4+0x4[sp]	4	ld.w \$_v, r14	4
jbr	.L4	2	cmp r14, r10	2
.L2:			jlt .L2	2
ld.w	\$_t, r14	4	st.w r0, \$_u	4
cmp	r0, r14	2	. L2:	
jne	.L4	2		
mov	1, r15	2		
st.w	r15, -4+0x4[sp]	4		
.L4:				
ld.w	-4+0x4[sp], r16	4		
cmp	r0, r16	2		
je	.L6	2		
ld.w	\$_u, r10	4		
add	1, r10	2		
ld.w	\$_v, r18	4		
cmp	r18, r10	2		
jlt	.L7	2		
st.w	r0, \$_u	4		
jbr	_ .L6	2		
.L7:				
st.w	r10, \$_u	4		
.L6:	_			
Total code size		70 bytes	Total code size	34 bytes

2.9 Transforming a for or while Loop Into a goto Loop

When a loop begins with a conditional decision expression such as a for or while loop, the CA850 generates the conditional decision expression twice as shown in the following image diagram.

This kind of loop transformation is performed by the front end (syntax analyzer), which is the initial phase of the compiler. The initial conditional decision is often eliminated by the subsequent optimization, but this kind of transformation is performed because it is advantageous from the viewpoint of increasing execution speed. However, when this description is not eliminated, the code size is increased.

[Image diagram]

Syntax	Image
<pre>for (statement-1; expression-2; statement-3) { loop-body }</pre>	<pre>statement-1; if (expression-2) { do { loop-body statement-3; } while (expression-2); }</pre>
<pre>while (expression-1) { loop-body }</pre>	<pre>if (expression-1) { do { loop-body } while (expression-1) ; }</pre>

Therefore, if the first conditional decision expression is not eliminated by the optimization, transforming to a loop formed by goto statements as follows enables the number of conditional decision expressions to be reduced to one. Write the coding as shown in the following image diagram.

[Image diagram]

for loop	<pre>statement-1; loop_bgn: if (! expression-2) goto loop_end; loop-body statement-3; goto loop_bgn;</pre>
	loop_end:
while loop	<pre>loop_bgn: if (! expression-1) goto loop_end; loop-body goto loop_bgn; loop_end:</pre>

A sample program is shown below.

Remark s and array[] are assumed to be external variables.

Before Modification	After Modification
int i;	int i;
<pre>for (i = 0; i < s; ++i) { array[i] = array[i+1]; }</pre>	<pre>i = 0; bgn_loop: if (i >= s) goto end_loop; array[i] = array[i+1]; ++i; goto bgn_loop; end_loop: ;</pre>

	Before Modification			After Modification			
	Program		Size [bytes]	3		Program	Size [bytes]
	ld.w	\$_s, r12	4		mov	r0, r10	2
	cmp	r0, r12	2	.L16:			
	jle	.L18	2		ld.w	\$_s, r11	4
	movea	<pre>\$_array, gp, r10</pre>	4		cmp	r11, r10	2
	mov	r0, r11	2		jge	.L20	2
.L16:					mov	r10, r12	2
	mov	r11, r13	2		shl	2, r12	2
	shl	2, r13	2		movea	\$_array, gp, r13	4
	add	r10, r13	2		add	r12, r13	2
	ld.w	4[r13], r14	4		ld.w	4[r13], r14	4
	st.w	r14, [r13]	4		st.w	r14, [r13]	4
	add	1, r11	2		add	1, r10	2
	ld.w	\$_s, r15	4		jbr	.L16	2
	cmp	r15, r11	2	.L20:			
	jlt	.L16	2				
.L18:							
Total code	otal code size 38 bytes			Total code	size		32 bytes

	Before Modification			After Modification			
	Program		Size [bytes]	Program		Size [bytes]	
	ld.w	\$_s, r12	4		mov	r0, r10	2
	cmp	r0, r12	2	.L16:			
	jle	.L18	2		ld.w	\$_s, r11	4
	movea	\$_array, gp, r10	4		cmp	r11, r10	2
	mov	r0, r11	2		jge	.L20	2
.L16:					mov	r10, r12	2
	mov	r11, r13	2		shl	2, r12	2
	shl	2, r13	2		movea	\$_array, gp, r13	4
	add	r10, r13	2		add	r12, r13	2
	ld.w	4[r13], r14	4		ld.w	4[r13], r14	4
	st.w	r14, [r13]	4		st.w	r14, [r13]	4
	add	1, r11	2		add	1, r10	2
	ld.w	\$ s, r15	4		jbr	.L16	2
	cmp	r15, r11	2	.L20:			
	jlt	.L16	2				
.L18:	-						
Total code	size		38 bytes	Total code	size		32 bytes

2.10 Unrolling a Loop

When the execution count is small and the loop body is small, the coding size of the unrolled loop may be smaller. In this case, the execution speed is also increased.

A sample program is shown below.

Remark array[] is assumed to be an external variable.

Before Modification	After Modification
int i;	int *p;
for (i = 0; i < 4; i++) {	p = array;
array[i] = 0;	*p = 0; *(p+1) = 0;
1	* (p+1) = 0;
	*(p+3) = 0;

[Output Assembly List for the V850]

	Before Modification			At	ter Modification	
	Program		Size [bytes]	Pro	gram	Size [bytes]
	movea	\$_array, gp, r10	4	st.w	r0, \$_array	4
	mov	r0, r11	2	st.w	r0, \$_array+4	4
.L15:				st.w	r0, \$_array+8	4
	mov	r11, r12	2	st.w	r0,	4
	shl	2, r12	2			
	add	r10, r12	2			
	st.w	r0, [r12]	4			
	add	1, r11	2			
	cmp	4, r11	2			
	jlt	.L15	2			
Total code	size		22 bytes	Total code size		16 bytes

Before Modification			After Modification			
	Program		Size [bytes]	Program		Size [bytes]
	movea	\$_array, gp, r10	4	st.w	r0, \$_array	4
	mov	r0, r11	2	st.w	r0, \$_array+4	4
.L15:				st.w	r0, \$_array+8	4
	mov	r11, r12	2	st.w	r0, \$_array+12	4
	shl	2, r12	2			
	add	r10, r12	2			
	st.w	r0, [r12]	4			
	add	1, r11	2			
	cmp	4, r11	2			
	jlt	.L15	2			
tal code	size		22 bytes	Total code size		16 byte

2.11 Shortening the Lifespan of a Variable

If there is an interval between the location where a value is assigned to a stack variable and the location where that value is actually referenced, a register is occupied during that interval and the opportunities for other variables being allocated to that register are reduced. Although the assignment of values is often moved to the end of the interval by the compiler optimization, this optimization may not be performed if the interval contains a function call.

In this case, changing the coding so that the assignment is performed immediately before the value is actually referenced will increase the register allocation opportunities of other variables, reduce memory accesses, and decrease the code size.

A sample program is shown below.

Remark s1, s2, s3, and array[] are assumed to be external variables.

Before Modification	After Modification
int i = 0, j = 0, k = 0, m = 0;	int i, j, k, m;
:	:
<pre>while ((k & 0xFF) != 0xFF) { k = s1; if (k & 0xFF00) { if (m != 1) { s2 += 2; } }</pre>	<pre>i = 0; j = 0; k = 0; m = 0; while ((k & 0xFF) != 0xFF) {</pre>
<pre>m = 1; array[15+i+j] = 0xFF; j++; }</pre>	k = s1; if (k & 0xFF00) { if (m != 1) { s2 += 2;
}	<pre>m = 1; array[15+i+j] = 0xFF; j++; } }</pre>

		Before Modification		After Modification			
		Program	Size [bytes]			Program	Size [bytes
	ld.w	\$_s1, r12	4		ld.w	\$_s1, r14	
	and	0xff00, r12	4		andi	0xff00, r14, r11	4
	mov	r0, r29	2		and	0xff, r14	4
	mov	r0, r28	2		mov	r0, r13	2
.L21:					mov	r0, r12	2
	ld.w	\$_s1, r11	4	.L22:			
	cmp	r0, r12	2		cmp	r0, r11	2
	je	.L27	2		je	.L28	2
	cmp	1, r29	2		cmp	1, r12	2
	je	.L27	2		je	.L28	2
	ld.w	\$_s2, r13	4		ld.w	\$_s2, r15	4
	add	2, r13	2		add	2, r15	2
	st.w	r13, \$_s2	4		st.w	r15, \$_s2	4
	mov	1, r29	2		mov	1, r12	2
	mov	r28, r14	2		mov	r13, r16	2
	shl	2, r14	2		shl	2, r16	2
	movea	\$_array, gp, r15	4		movea	<pre>\$_array, gp, r17</pre>	4
	add	r14, r15	2		add	r16, r17	2
	mov	255, r16	4		mov	255, r18	4
	st.w	r16, 60[r15]	4		st.w	r18, 60[r17]	4
	add	1, r28	2		add	1, r13	2
.L27:				.L28:			
	andi	0xff, r11, r17	4		addi	-255, r14, r0	4
	addi	-255, r17, r0	4		jne	.L22	2
	jne	.L21	2				
otal code	size		66 bytes	Total code	62 byte		

	Before Modification				After Modification			
	!	Program	Size [bytes]	Program			Size [bytes]	
	ld.w	\$_s1, r12	4		ld.w	\$_s1, r14	4	
	and	0xff00, r12	4		andi	0xff00, r14, r11	4	
	mov	r0, r29	2		zxb	r14	2	
	mov	r0, r28	2		mov	r0, r13	2	
.L21:					mov	r0, r12	2	
	ld.w	\$_s1, r11	4	.L22:				
	cmp	r0, r12	2		cmp	r0, r11	2	
	je	.L27	2		je	.L28	2	
	cmp	1, r29	2		cmp	1, r12	2	
	je	.L27	2		je	.L28	2	
	ld.w	\$_s2, r13	4		ld.w	\$_s2, r15	4	
	add	2, r13	2		add	2, r15	2	
	st.w	r13, \$_s2	4		st.w	r15, \$_s2	4	
	mov	1, r29	2		mov	1, r12	2	
	mov	r28, r14	2		mov	r13, r16	2	
	shl	2, r14	2		shl	2, r16	2	
	movea	\$_array, gp, r15	4		movea	\$_array, gp, r17	4	
	add	r14, r15	2		add	r16, r17	2	
	mov	255, r16	4		mov	255, r18	4	
	st.w	r16, 60[r15]	4		st.w	r18, 60[r17]	4	
	add	1, r28	2		add	1, r13	2	
.L27:				.L28:				
	zxb	r11	2		addi	-255, r14, r0	4	
	addi	-255, r11, r0	4		jne	.L22	2	
	jne	.L21	2					
Total code	Total code size			Total code	size		60 bytes	

2.12 Eliminating an Induction Variable

A variable that controls a loop is called an induction variable, and optimization that eliminates an induction variable by changing the loop control so that another variable is used is referred to as "eliminating an induction variable."

Although this optimization is also included in the CA850, since the conditions for which it is applied are limited, not all cases can be optimized.

Therefore, perform this optimization "manually" by altering the program as follows.

A sample program is shown below.

Remark x and *table are assumed to be external variables.

Before Modification	After Modification
int i;	const unsigned short *p;
<pre>for (i = 0; *(table + i) != NULL; ++i) { if ((*(table + i) & 0xFF) == x) { return(*(table + i) & 0xFF00); } }</pre>	<pre>for (p = table; *p != NULL; ++p) { if ((*p & 0xFF) == x) { return(*p & 0xFF00); } }</pre>

		Before Modification				After Modification	
		Program	Size [bytes]			Program	Size [bytes]
1	Ld.w	\$_table, r12	4		ld.w	\$_table, r12	4
1	ld.h	[r12], r13	4		ld.h	[r12], r13	4
а	and	0xffff, r13	4		and	0xffff, r13	4
j	je	.L1	2		je	.L1	2
1	ld.h	[r12], r10	4		ld.h	[r12], r10	4
π	nov	r0, r11	2		mov	r12, r11	2
.L2:				.L2:			
а	andi	0xff, r10, r14	4		andi	0xff, r10, r14	4
1	ld.h	\$_x, r15	4		ld.h	\$_x, r15	4
а	and	0xffff, r15	4		and	0xffff, r15	4
C	cmp	r14, r15	2		cmp	r14, r15	2
j	jne	.L5	2		jne	.L5	2
а	andi	0xff00, r10, r16	4		andi	0xff00, r10, r16	4
S	st.w	r16, -4+0x4[sp]	4		st.w	r16, -4+0x4[sp]	4
j	jbr	.L1	2		jbr	.L1	2
.L5:				.L5:			
а	add	1, r11	2		add	2, r11	2
π	nov	r11, r10	2		ld.h	[r11], r10	4
S	shl	1, r10	2		and	0xffff, r10	4
а	add	r12, r10	2		jne	.L2	2
1	ld.h	[r10], r10	4	.L1:			
а	and	0xffff, r10	4		ld.w	-4+0x4[sp], r10	4
j	jne	.L2	2				
.L1:							
1	Ld.w	-4+0x4[sp], r10	4				
Total code	e size		68 bytes	Total co	ode size		62 byte

		Before Modification				After Modification	
		Program	Size [bytes]			Program	Size [bytes]
ld	d.w	\$_table, r12	4		ld.w	\$_table, r12	4
ld	d.hu	[r12], r13	4		ld.hu	[r12], r13	4
cm	mp	r0, r13	2		cmp	r0, r13	2
je	Э	.L1	2		je	.L1	2
ld	d.hu	[r12], r10	4		ld.hu	[r12], r10	4
mo	ov	r0, r11	2		mov	r12, r11	2
.L2:				.L2:			
mc	vc	r10, r14	2		mov	r10, r14	2
ZX	хb	r14	2		zxb	r14	2
ld	d.hu	\$_x, r15	4		ld.hu	\$_x, r15	4
cm	np	r14, r15	2		cmp	r14, r15	2
jn	ne	.L5	2		jne	.L5	2
an	ndi	0xff00, r10, r16	4		andi	0xff00, r10, r16	4
st	t.w	r16, -4+0x4[sp]	4		st.w	r16, -4+0x4[sp]	4
jb	or	.L1	2		jbr	.L1	2
.L5:				.L5:			
ad	dd	1, r11	2		add	2, r11	2
mc	vc	r11, r10	2		ld.hu	[r11], r10	4
sh	nl	1, r10	2		cmp	r0, r10	2
ad	dd	r12, r10	2		jne	.L2	2
ld	d.hu	[r10], r10	4	.L1:			
cm	np	r0, r10	2		ld.w	-4+0x4[sp], r10	4
jn	ne	.L2	2				
.L1:							
ld	d.w	-4+0x4[sp], r10	4				
Total code	size		60 bytes	Total co	de size		54 bytes

2.13 Setting (unsigned) int Type for (unsigned) short or char Type Variables

According to the (unsigned) ANSI-C specifications, since (unsigned) short or (unsigned) char type variables are extended to int type or unsigned int type when calculations are performed, type conversion instructions often are generated for programs that use these kinds of variables (in particular, when these kinds of variables are allocated to registers).

Since this type conversion will be unnecessary if (unsigned) int type is set, the code size is reduced.

In particular, (unsigned) int type should be used as much as possible for stack variables, which are allocated to registers relatively often.

A sample program is shown below.

Remark array[] and *p are assumed to be external variables.

Before Modification	After Modification
unsigned char i;	int i;
<pre>for(i = 0; i < 4; i++) { array[2 + i] = *(p + i); }</pre>	<pre>for(i = 0; i < 4; i++) { array[2 + i] = *(p + i); }</pre>

		Before Modification				After Modification	
		Program	Size [bytes]			Program	Size [bytes]
	movea	\$_array, gp, r10	4		movea	\$_array, gp, r10	4
	mov	r0, r12	2		mov	r0, r12	2
.L2:				.L2:			
	mov	r12, r11	2		mov	r12, r11	2
	shl	2, r11	2		shl	2, r11	2
	mov	r10, r13	2		mov	r10, r13	2
	add	r11, r13	2		add	r11, r13	2
	ld.w	\$_p, r15	4		ld.w	\$_p, r15	4
	add	r11, r15	2		add	r11, r15	2
	ld.w	[r15], r15	4		ld.w	[r15], r15	4
	st.w	r15, 8[r13]	4		st.w	r15, 8[r13]	4
	add	1, r12	2		add	1, r12	2
	and	0xff, r12	4		cmp	4, r12	2
	cmp	4, r12	2		jlt	.L2	2
	jlt	.L2	2				
Total co	de size		38 bytes	Total co	ode size		34 bytes

		Before Modification				After Modification	
		Program	Size [bytes]			Program	Size [bytes]
	movea	\$_array, gp, r10	4	mo	ovea	\$_array, gp, r10	4
	mov	r0, r12	2	mo	ov	r0, r12	2
.L2:				.L2:			
	mov	r12, r11	2	mo	ov	r12, r11	2
	shl	2, r11	2	sl	hl	2, r11	2
	mov	r10, r13	2	mo	ov	r10, r13	2
	add	r11, r13	2	ac	dd	r11, r13	2
	ld.w	\$_p, r15	4	10	d.w	\$_p, r15	4
	add	r11, r15	2	ac	dd	r11, r15	2
	ld.w	[r15], r15	4	10	d.w	[r15], r15	4
	st.w	r15, 8[r13]	4	st	t.w	r15, 8[r13]	4
	add	1, r12	2	ac	dd	1, r12	2
	zxb	r12	2	cr	mp	4, r12	2
	cmp	4, r12	2	j]	lt	.L2	2
	jlt	.L2	2				
Total co	ode size	·	36 bytes	Total code	size		34 bytes

2.14 Consolidating in a Single Statement When an Assigned Value Is Referenced in the Statement Following the Assignment Statement

When an assigned value is referenced in the statement following an assignment statement, consolidating the two statements in a single statement by replacing the referencing location with the assignment statement may eliminate an extra register transfer and reduce the code size. However, in many cases, the code size will not change because the extra register transfer has been eliminated by the compiler optimization.

A sample program is shown below.

Remark s and t are assumed to be external variables.

Before Modification	After Modification
s;	if (s == 0) {
if (s == 0) {	t++;
t++;	}
}	

In this example, the same code size will be produced by compiler optimization.

2.15 Eliminating the if~else Statement When the Branch Destinations of the if~else Statement Are return Statements That Return the Result of the Branch Condition

When each of the branch destinations of an if~else statement contains only a return statement and the corresponding return values are the branch condition results themselves, eliminate the if~else statement and make the program return the value of the branch condition expression.

A sample program is shown below.

Remark s1 and s2 are assumed to be external variables.

Before Modification	After Modification
<pre>if (s1 == s2) { return(1);</pre>	return (s1 == s2);
}	
return(0);	

[Output Assembly List for the V850]

	Before Modification		After Modification	
	Program	Size [bytes]	Program	Size [bytes]
ld.w	\$_s1, r11	4	ld.w \$_s1, r11	4
ld.w	\$_s2, r12	4	ld.w \$_s2, r12	4
cmp	r12, r11	2	cmp r12, r11	2
jne	.L2	2	setfe r10	4
mov	1, r10	2		
jbr	.L1	2		
.L2:				
mov	r0, r10	2		
.L1:				
Total code size		18 bytes	Total code size	14 bytes

	Before Modification		After Modification		
	Program	Size [bytes]	Program	Size [bytes]	
ld.w	\$_s1, r11	4	ld.w \$_s1, r11	4	
ld.w	\$_s2, r12	4	ld.w \$_s2, r12	4	
cmp	r12, r11	2	cmp r12, r11	2	
jne	.L2	2	setfe r10	4	
mov	1, r10	2			
jbr	.L1	2			
.L2:					
mov	r0, r10	2			
.L1:					
Total code size		18 bytes	Total code size	14 bytes	

2.16 Changing the Condition and Setting the Operand to 15 or –16 When One of the Operands of a Comparison Operation Is the Constant 16 or –17

When one of the operands of the comparison instruction cmp is not immediate data in the range –16 to 15, the assembler performs instruction expansion and the comparison instruction will become multiple instructions.

Therefore, if the operand value is held to the range from –16 to 15 by changing the condition, the expansion will be suppressed and the coding size will be reduced.

Specifically, change the value used for the comparison in the conditional expression of a statement such as a for or if statement.

A sample program is shown below.

Remark s is assumed to be an external variable.

Before Modification	After Modification
int i;	int i;
for (i = 0; i < 16; i++) { s++;	for (i = 0; i <= 15; i++) { s++;
}	}

In this example, the same code size will be produced by compiler optimization.

2.17 Initializing a Variable

When an auto variable is used in a function without being initialized, the code size is increased because that variable remains in memory without being allocated to a register.

In the example shown below, when the variable a does not correspond to either case of the switch statement, it is referenced by the return statement without being initialized. Actually, even if it always corresponds to one of the cases, it may be treated as if it is not initialized because its value is not known when the program is analyzed during register allocation. In this case, no register will be allocated during CA850 register allocation.

Therefore, if the variable is allocated to a register by adding an initialization, the code size can be reduced.

A sample program is shown below.

Remark a is assumed to be an auto variable.

Before Modification	After Modification
int a;	int a = 0;
<pre>switch(x) { case 0: a = 0; break;</pre>	<pre>switch(x) { case 0: a = 0; break;</pre>
case 1:	case 1:
<pre>a = 1; } return a;</pre>	<pre>a = 1; } return a;</pre>

	Before Modification		After Modification		
	Program	Size [bytes]		Program	Size [bytes]
cmp	r0, r6	2	cmp	r0, r6	2
je	.L4	2	mov	r0, r11	2
cmp	1, r6	2	je	.L3	2
je	.L5	2	cmp	1, r6	2
jbr	.L3	2	jne	.L3	2
.L4:			mov	1, r11	2
st.w	r0, -4+0x4[sp]	4	.L3:		
jbr	.L3	2	mov	r11, r10	2
.L5:					
mov	1, r11	2			
st.w	r11, -4+0x4[sp]	4			
.L3:					
ld.w	-4+0x4[sp], r10	4			
Total code size		26 bytes	Total code size		14 bytes

	Before Modification			After Modification	
	Program	Size [bytes]		Program	Size [bytes]
cmp	r0, r6	2	cmp	r0, r6	2
je	.L4	2	mov	r0, r11	2
cmp	1, r6	2	je	.L3	2
je	.L5	2	cmp	1, r6	2
jbr	.L3	2	jne	.L3	2
.L4:			mov	1, r11	2
st.v	w r0, -4+0x4[sp]	4	.L3:		
jbr	.L3	2	mov	r11, r10	2
.L5:					
mov	1, r11	2			
st.v	w r11, -4+0x4[sp]	4			
.L3:					
ld.v	w -4+0x4[sp], r10	4			
Total code siz	ze	26 bytes	Total code size		14 bytes

2.18 Declaring void for a Function Having No Return Value

Declare void for a function that has no return value.

The extra register transfer instruction will be eliminated.

A sample program is shown below.

Before Modification	After Modification
<pre>func(int a)</pre>	<pre>void func(int a)</pre>
{	{
switch(a) {	switch(a) {
case 0:	case 0:
s = 0;	s = 0;
break;	break;
case 1:	case 1:
s = 1;	s = 1;
}	}
}	}

[Output Assembly List for the V850]

	Before Modification		After Modification	
	Program	Size [bytes]	Program	Size [bytes]
cmp	r0, r6	2	cmp r0, r6	2
je	.L4	2	je .L4	2
cmp	1, r6	2	cmp 1, r6	2 2 2 2
je	.L5	2 2	je .L5	2
jbr	.L3	2	jbr .L3	2
.L4:			.L4:	
st.w	r0, \$_s	4	st.w r0, \$_s	4
jbr	.L3	2	jbr .L3	2
.L5:			.L5:	
mov	1, r11	2	mov 1, r10	2
st.w	r11, \$_s	4	st.w r10, \$_s	4
.L3:			.L3:	
mov	r0, r10	2	jmp [lp]	2
jmp	[lp]	2		
Total code size		26 bytes	Total code size	24 bytes

_	Before Modification			After Modification	·
	Program	Size		Program	Size
	_	[bytes]		-	[bytes]
cmp	r0, r6	2	cmp	r0, r6	2
je	.L4	2	je	.L4	2
cmp	1, r6	2 2 2	cmp	1, r6	2
je	.L5	2	je	.L5	2
jbr	.L3	2	jbr	.L3	2
.L4:			.L4:		
st.w	r0, \$_s	4	st.w	r0, \$_s	4
jbr	.L3	2	jbr	.L3	2
.L5:			.L5:		
mov	1, r11	2	mov	1, r10	2
st.w	r11, \$_s	4	st.w	r10, \$_s	4
.L3:			.L3:		
mov	r0, r10	2	jmp	[lp]	2
jmp	[lp]	2			
Total code size		26 bytes	Total code size		24 bytes

2.19 Consolidating Common case Processing in a switch Statement

If the identical case processing descriptions of a switch statement are consolidated, the code size may be reduced.

A sample program is shown below.

Remark x is assumed to be an external variable.

Before Modification	After Modification
switch(x) {	switch(x) {
case 0:	case 0:
<pre>dummy1();</pre>	case 1:
break;	case 2:
case 1:	dummy1();
<pre>dummy1();</pre>	break;
break;	case 3:
case 2:	case 4:
<pre>dummy1();</pre>	dummy2();
break;	break;
case 3:	default:
dummy2();	break;
break;	}
case 4:	
dummy2();	
break;	
default:	
break;	
}	

Before Modification		After Modification		
	Program	Size [bytes]	Program	Size [bytes]
ld.w	-4+.A2[sp], r10	4	ld.w -4+.A2[sp], r10	4
cmp	4, r10	2	cmp 4, r10	2
jh	.L3	2	jh .L3	2
shl	1, r10	2	shl 1, r10	2
add	tp, r10	2	add tp, r10	2
ld.h	.L10[r10], r11	8	ld.h .L10[r10], r11	8
add	.L10, r11	10	add .L10, r11	10
add	tp, r11	2	add tp, r11	2
jmp	[r11]	2	jmp [r11]	2
.L10:			.L10:	
.hword	.L13L10	2	.hword .L13L10	2
.hword	.L16L10	2	.hword .L13L10	2
.hword	.L19L10	2	.hword .L13L10	2
.hword	.L22L10	2	.hword .L16L10	2
.hword	.L25L10	2	.hword .L16L10	2
.L13:			.L13:	
jarl	_dummy1, lp	4	jarl _dummy1, lp	4
jbr	L3	2	jbr .L3	2
.L16:			.L16:	
jarl	_dummy1, lp	4	jarl _dummy2, lp	4
jbr	L3	2	.L3:	
.L19:				
jarl	_dummy1, lp	4		
jbr	.L3	2		
.L22:				
jarl	_dummy2, lp	4		
jbr	.L3	2		
.L25:				
jarl	dummy2, lp	4		
.L3:				
Total code size		72 bytes	Total code size	54 bytes

Before Modification		After Modification	
Program	Size [bytes]	Program	Size [bytes]
ld.w -4+.A2[sp], r10	4	ld.w -4+.A2[sp], r10	4
cmp 4, r10	2	cmp 4, r10	2
jh .L3	2	jh .L3	2
switch r10	2	switch r10	2
.L10:		.L10:	
.shword .L13L10	2	.shword .L13L10	2
.shword .L16L10	2	.shword .L13L10	2
.shword .L19L10	2	.shword .L13L10	2
.shword .L22L10	2	.shword .L16L10	2
.shword .L25L10	2	.shword .L16L10	2
.L13:		.L13:	
jarl _dummy1, lp	4	jarl _dummy1, lp	4
jbr .L3	2	jbr .L3	2
.L16:		.L16:	
jarl _dummy1, lp	4	jarl _dummy2, lp	4
jbr .L3	2	L3:	
.L19:			
jarl _dummy1, lp	4		
jbr .L3	2		
.L22:			
jarl _dummy2, lp	4		
jbr .L3	2		
.L25:			
jarl dummy2, lp	4		
.L3:			
Total code size	48 bytes	Total code size	30 bytes

2.20 Consolidating return Statements Having Identical Values

If descriptions of return statements that return the same value are consolidated, the code size may be reduced.

A sample program is shown below.

Remark s, t, and u are assumed to be external variables.

Before Modification	After Modification
<pre>if(s == 1)return 0xff;</pre>	if((s == 1) (t == 1) (u == 1))return 0xff;
<pre>if(t == 1)return 0xff;</pre>	return 0x0;
<pre>if(u == 1)return 0xff;</pre>	
return 0x0;	

	Before Modification		After Modification	
	Program	Size [bytes]	Program	Size [bytes]
ld.w	\$_s, r11	4	ld.w \$_s, r11	4
cmp	1, r11	2	cmp 1, r11	2
jne	.L2	2	je .L3	2
mov	255, r10	4	ld.w	4
jbr	.L1	2	cmp 1, r12	2
.L2:			je .L3	2
ld.w	\$_t, r12	4	ld.w \$_u, r13	4
cmp	1, r12	2	cmp 1, r13	2
jne	.L3	2	je .L3	2
mov	255, r10	4	mov r0, r10	2
jbr	.L1	2	jbr .L1	2
.L3:			.L3:	
ld.w	\$_u, r13	4	mov 255, r10	4
cmp	1, r13	2	.L1:	
jne	.L4	2		
mov	255, r10	4		
jbr	.L1	2		
.L4:				
mov	r0, r10	2		
.L1:				
Total code size		44 bytes	Total code size	32 bytes

	Before Modification			After Modification	
	Program	Size [bytes]		Program	Size [bytes]
ld.w	s_s, r11	4	ld.w	\$_s, r11	4
cmp	1, r11	2	cmp	1, r11	2
jne	.L2	2	je	.L3	2
mov	255, r10	4	ld.w	\$_t, r12	4
jbr	.L1	2	cmp	1, r12	2
.L2:			je	.L3	2
ld.w	, \$_t, r12	4	ld.w	\$_u, r13	4
cmp	1, r12	2	cmp	1, r13	2
jne	.L3	2	je	.L3	2
mov	255, r10	4	mov	r0, r10	2
jbr	.L1	2	jbr	.L1	2
.L3:			.L3:		
ld.w	, \$_u, r13	4	mov	255, r10	4
cmp	1, r13	2	.L1:		
jne	.L4	2			
mov	255, r10	4			
jbr	.L1	2			
.L4:					
mov	r0, r10	2			
.L1:					
Total code siz	re	44 bytes	Total code size		32 bytes

2.21 Making an Expanded Inline Function into a static Function

When a function for which inline expansion is specified is not referenced from another file, making it into a static function will cause the code for the function itself to be eliminated by the optimization. As a result, the code size may be reduced. However, when the program^{Note} contains an assembly language description, the code for the function itself is output because this optimization is not performed. In this case, make the function that contains the assembly language description a separate file.

Note This refers to the program (including the include file) that is to be compiled.

Example For a static function and a normal function

Note that the function name has been changed for the sake of explanation in the Before Modification and After Modification coding.

Before Modification	After Modification
#pragma inline funclsub	#pragma inline func2sub
int s,t;	int s,t;
<pre>void funclsub()</pre>	static void func2sub()
{	{
int tmp;	int tmp;
tmp = s;	tmp = s;
s = t;	s = t;
t = tmp;	t = tmp;
}	}
<pre>void func1()</pre>	void func2()
{	{
if(s == 1){	if(s == 1){
func1sub();	func2sub();
}	}
}	}

Although the code for the function func1sub is output, no code is output for the function func2sub.

Example For a function containing an assembly language description and a function not containing an assembly language description

Note that the function name has been changed for the sake of explanation in the Before Modification and After Modification coding.

Before Modification	After Modification
#pragma inline func3sub	#pragma inline func4sub
int s,t;	int s,t;
static void func3sub()	static void func4sub()
{	{
int tmp;	int tmp;
<pre>tmp = s; s = t; t = tmp; }</pre>	<pre>tmp = s; s = t; t = tmp; }</pre>
void func3()	void func4()
{	{
if(s == 1){	if(s == 1){
func3sub();	func4sub();
}	}
}	}
<pre>void dummy(void) { asm("nop"); }</pre>	

Since the program shown in the Before Modification coding contains an assembly language description, the code for the function func3sub is output. However, since the program shown in the After Modification coding does not contain an assembly language description, no code is output for the function func4sub.

CHAPTER 3 INCREASING EXECUTION SPEED

This chapter describes examples in which the effective speed of processing that is essentially entrusted to the compiler can be increased when a human operator alters the description.

3.1 Using a Pointer for Consecutive Accesses to an Array

Use a pointer for consecutive accesses to an array within a loop. If a pointer is not used, processing for obtaining the actual address from the array subscript may be output each time.

A sample program is shown below.

Remark sum and array[] are assumed to be external variables.

Before Modification	After Modification				
int i;	int i,*p;				
<pre>sum = 0; for(i = 0 ; i < 10 ; i++){ sum += array[i]; }</pre>	<pre>sum = 0; p = &array[0]; for(i = 0 ; i < 10 ; i++){ sum += *p++; }</pre>				

	Before Modification				After Modification			
	Program				Program			
	st.w	r0, \$_sum	4		movea	\$_array, gp, r12	4	
	movea	<pre>\$_array, gp, r10</pre>	4		st.w	r0, \$_sum	4	
	mov	r0, r12	2		mov	r0, r13	2	
	mov	r0, r11	2		mov	r0, r11	2	
.L16:				.L17:				
	mov	r11, r14	2		mov	r12, r14	2	
	shl	2, r14	2		add	4, r12	2	
	add	r10, r14	2		ld.w	[r14], r14	4	
	ld.w	[r14], r14	4		add	r14, r13	2	
	add	r14, r12	2		add	1, r11	2	
	add	1, r11	2		st.w	r13, \$_sum	4	
	st.w	r12, \$_sum	4		cmp	10, r11	2	
	cmp	10, r11	2		jlt	.L17	2	
	jlt	.L16	2					
Total code	Total code size 34		34 bytes	Total code	size		32 bytes	

	Before Modification				After Modification			
	Program			Program			Size [bytes]	
	st.w	r0, \$_sum	4		movea	\$_array, gp, r12	4	
	movea	<pre>\$_array, gp, r10</pre>	4		st.w	r0, \$_sum	4	
	mov	r0, r12	2		mov	r0, r13	2	
	mov	r0, r11	2		mov	r0, r11	2	
.L16:				.L17:				
	mov	r11, r14	2		mov	r12, r14	2	
	shl	2, r14	2		add	4, r12	2	
	add	r10, r14	2		ld.w	[r14], r14	4	
	ld.w	[r14], r14	4		add	r14, r13	2	
	add	r14, r12	2		add	1, r11	2	
	add	1, r11	2		st.w	r13, \$_sum	4	
	st.w	r12, \$_sum	4		cmp	10, r11	2	
	cmp	10, r11	2		jlt	.L17	2	
	jlt	.L16	2					
Total code	Total code size 34		34 bytes	Total code	size		32 bytes	

3.2 Replacing an Access to an External Variable with an Access to a Temporary Variable

Avoid using an external variable as much as possible within a loop.

Since an address calculation or memory access (load/store instruction) may be output every time, replace the access to an external variable with an access to a temporary variable.

A sample program is shown below.

Remark sum and array[] are assumed to be external variables.

Before Modification	After Modification
int i;	int i;
int *p;	int *p;
	int tmp;
sum = 0;	
<pre>p = &array[0];</pre>	tmp = 0;
for(i = 0; i < 10; i++){	<pre>p = &array[0];</pre>
sum += *p++;	for(i = 0; i < 10; i++){
}	tmp += *p++;
	}
	sum = tmp;

	Before Modification				After Modification			
	Program			Program			Size [bytes]	
	movea	\$_array, gp, r12	4		movea	\$_array, gp, r12	4	
	st.w	r0, \$_sum	4		mov	r0, r11	2	
	mov	r0, r13	2		mov	r0, r13	2	
	mov	r0, r11	2	.L18:				
.L17:					mov	r12, r14	2	
	mov	r12, r14	2		add	4, r12	2	
	add	4, r12	2		ld.w	[r14], r14	4	
	ld.w	[r14], r14	4		add	r14, r13	2	
	add	r14, r13	2		add	1, r11	2	
	add	1, r11	2		cmp	10, r11	2	
	st.w	r13, \$ sum	4		jlt	.L18	2	
	cmp	10, r11	2		st.w	r13, \$_sum	4	
	jlt	.L17	2					
Total code	Total code size		32 bytes	Total code	size		28 bytes	

	Before Modification				After Modification			
		Program	Size [bytes]	Program			Size [bytes]	
	movea	\$_array, gp, r12	4		movea	\$_array, gp, r12	4	
	st.w	r0, \$_sum	4		mov	r0, r11	2	
	mov	r0, r13	2		mov	r0, r13	2	
	mov	r0, r11	2	.L18:				
.L17:					mov	r12, r14	2	
	mov	r12, r14	2		add	4, r12	2	
	add	4, r12	2		ld.w	[r14], r14	4	
	ld.w	[r14], r14	4		add	r14, r13	2	
	add	r14, r13	2		add	1, r11	2	
	add	1, r11	2		cmp	10, r11	2	
	st.w	r13, \$ sum	4		jlt	.L18	2	
	cmp	10, r11	2		st.w	r13, \$_sum	4	
	jlt	.L17	2					
Total code	Total code size		32 bytes	Total code	size		28 bytes	

3.3 Not Using a Variable Expression for a Loop Ending Condition

Avoid using a variable expression as much as possible for a loop ending condition. Use a temporary variable instead.

If a variable expression is used, a calculation for the ending expression comparison may be output every time.

A sample program is shown below.

Remark array[][] is assumed to be an external variable.

Before Modification	After Modification			
int i;	int i;			
int nSize;	int nSize;			
int mSize;	int mSize;			
int *p;	int *p;			
	int s;			
p = &array[0][0];				
for(i = 0; i < nSize * mSize; i++){	p = &array[0][0];			
*p++ = 0;	s = nSize * mSize;			
}	for(i = 0; i < s; i++){			
	*p++ = 0;			
	}			

In this example, the same code size will be produced by compiler optimization.

3.4 Using a Comparison with Zero for the Loop Ending Condition

If a comparison with zero expression is used for the loop ending condition, the calculation of the ending condition each time through the loop may become faster. Also, the number of registers that are used may be reduced.

A sample program is shown below.

Remark array[][] is assumed to be an external variable.

Before Modification	After Modification			
int i;	int i;			
int nSize;	int nSize;			
int mSize;	int mSize;			
int *p;	int *p;			
int s;				
	p = &array[0][0];			
p = &array[0][0];	for(i = nSize * mSize; i > 0; i){			
s = nSize * mSize;	*p++ = 0;			
for(i = 0; i < s; i++){	}			
*p++ = 0;				
}				

[Output Assembly List for the V850]

		Before Modification		After Modification			
	Program			Size Program [bytes]			Size [bytes]
	ld.w	-4+0x8[sp], r7	4		ld.w	-4+0x8[sp], r7	4
	ld.w	-8+0x8[sp], r6	4		ld.w	-8+0x8[sp], r6	4
	jarl	mul, lp	4		jarl	mul, lp	4
	cmp	r0, r6	2		cmp	r0, r6	2
	mov	r6, r13	2		mov	r6, r10	2
	jle	.L24	2		jle	.L23	2
	movea	\$_array, gp, r11	4		movea	\$_array, gp, r11	4
	mov	r0, r12	2		mov	r10, r12	2
.L22:				.L21:			
	mov	r11, r10	2		mov	r11, r10	2
	add	4, r11	2		add	4, r11	2
	st.w	r0, [r10]	4		st.w	r0, [r10]	4
	add	1, r12	2		add	4294967295, r12	2
	cmp	r12, r13	2		cmp		2
	jgt	.L22	2		jgt		2
.L24:				.L23:			
Total code	size		38 bytes	Total code	size		38 bytes

Remark In this example, the total code size does not change.

	Before Modification				After Modification			
	Program			Program			Size [bytes]	
	ld.w	-4+0x8[sp], r13	4		ld.w	-4+0x8[sp], r10	4	
	ld.w	-8+0x8[sp], r14	4		ld.w	-8+0x8[sp], r13	4	
	mul	r14, r13, r0	4		mul	r13, r10, r0	4	
	cmp	r0, r13	2		cmp	r0, r10	2	
	jle	.L24	2		jle	.L23	2	
	movea	\$_array, gp, r11	4		movea	<pre>\$_array, gp, r11</pre>	4	
	mov	r0, r12	2		mov	r10, r12	2	
.L22:				.L21:				
	mov	r11, r10	2		mov	r11, r10	2	
	add	4, r11	2		add	4, r11	2	
	st.w	r0, [r10]	4		st.w	r0, [r10]	4	
	add	1, r12	2		add	4294967295, r12	2	
	cmp	r12, r13	2		cmp	r0, r12	2	
	jgt	.L22	2		jgt	.L21	2	
.L24:				.L23:				
Total code	Total code size		36 bytes	Total code	size		36 bytes	

Remark In this example, the total code size does not change.

3.5 Unrolling a Loop

If the looping frequency is reduced, the overhead for the branching instructions due to the loop will be reduced.

A sample program is shown below.

Remark array[] is assumed to be an external variable.

```
Before Modification
                                                                      After Modification
int i;
                                                      int i;
int *p;
                                                      int *p;
p = array;
                                                      p = array;
for(i = N; i > 0; i--){
                                                      for(i = N >> 2; i > 0; i--){ /* N/4 */}
  *p++ = 0;
                                                        *p++ = 0;
                                                        *p++ = 0;
                                                        *p++ = 0;
                                                        *p++ = 0;
                                                      for(i = N & 3; i > 0; i--) \{ /* \text{ N mod 4 *} / \}
                                                        *p++ = 0;
```

	Before Modification				After Modification			
	I	Program	Size [bytes]	Program			Size [bytes]	
	movea	\$_array, gp, r12	4		mov	2, r6	2	
	mov	10, r11	2		movea	\$_array, gp, r10	4	
.L16:					mov	r6, r11	2	
	mov	r12, r10	2	.L16:				
	add	4, r12	2		addi	4, r10, r13	4	
	st.w	r0, [r10]	4		st.w	r0, [r10]	4	
	add	4294967295, r11	2		addi	4, r13, r14	4	
	cmp	r0, r11	2		st.w	r0, [r13]	4	
	jgt	.L16	2		addi	4, r14, r7	4	
					st.w	r0, [r14]	4	
					addi	4, r7, r10	4	
					st.w	r0, [r7]	4	
					add	4294967295, r11	2	
					cmp	r0, r11	2	
					<u>jg</u> t	.L16	2	
					mov	r6, r12	2	
				.L23:				
					mov	r10, r11	2	
					add	4, r10	2	
					st.w	r0, [r11]	4	
					add	4294967295, r12	2	
					cmp	r0, r12	2	
					jgt	.L23	2	
Total code	size		20 bytes	Total code	size		62 bytes	

		Before Modification		After Modification			
Program		Size [bytes]	Program			Size [bytes]	
	movea	\$_array, gp, r12	4		mov	2, r6	2
	mov	10, r11	2		movea	<pre>\$_array, gp, r10</pre>	4
.L16:					mov	r6, r11	2
	mov	r12, r10	2	.L16:			
	add	4, r12	2		addi	4, r10, r13	4
	st.w	r0, [r10]	4		st.w	r0, [r10]	4
	add	4294967295, r11	2		addi	4, r13, r14	4
	cmp	r0, r11	2		st.w	r0, [r13]	4
	jgt	.L16	2		addi	4, r14, r7	4
					st.w	r0, [r14]	4
					addi	4, r7, r10	4
					st.w	r0, [r7]	4
					add	4294967295, r11	2
					cmp	r0, r11	2
					jgt	.L16	2
					mov	r6, r12	2
				.L23:			
					mov	r10, r11	2
					add	4, r10	2
					st.w	r0, [r11]	4
					add	4294967295, r12	2
					cmp	r0, r12	2
					jgt	.L23	2
Total code	size		20 bytes	Total code	size		62 bytes

3.6 Optimizing a Pointer

A sample program is shown below.

Remark array[] is assumed to be an external variable and N is assumed to be 10.

```
Before Modification
                                                                    After Modification
int i;
                                                    int i;
int *p;
                                                    int *p;
p = array;
                                                    p = array;
for(i = N >> 2; i > 0; i--) \{ /* N/4 */
                                                    for(i = N >> 2; i > 0; i--) { /* N/4 */
                                                      *(p + 0) = 0;
  *p++ = 0;
                                                      *(p + 1) = 0;
  *p++ = 0;
                                                      *(p + 2) = 0;
  *p++ = 0;
                                                      *(p + 3) = 0;
                                                      p += 4;
for(i = N \& 3; i > 0; i--){ /* N mod 4 */
                                                    for(i = N & 3; i > 0; i--) \{ /* \text{ N mod } 4 */ \}
  *p++ = 0;
                                                       *p++ = 0;
```

		Before Modification		After Modification				
		Program	Size [bytes]	Program			Size [bytes]	
	mov	2, r6	2		mov	2, r10	2	
	movea	<pre>\$_array, gp, r10</pre>	4		movea	<pre>\$_array, gp, r12</pre>	4	
	mov	r6, r11	2		mov	r10, r11	2	
.L16:				.L16:				
	addi	4, r10, r13	4		st.w	r0, [r12]	4	
	st.w	r0, [r10]	4		st.w	r0, 4[r12]	4	
	addi	4, r13, r14	4		st.w	r0, 8[r12]	4	
	st.w	r0, [r13]	4		st.w	r0, 12[r12]	4	
	addi	4, r14, r7	4		add	16, r12	4	
	st.w	r0, [r14]	4		add	4294967295, r11	2	
	addi	4, r7, r10	4		cmp	r0, r11	2	
	st.w	r0, [r7]	4		jgt	.L16	2	
	add	4294967295, r11	2		mov	r10, r11	2	
	cmp	r0, r11	2	.L23:				
	jgt	.L16	2		mov	r12, r10	2	
	mov	r6, r12	2		add	4, r12	2	
.L23:					st.w	r0, [r10]	4	
	mov	r10, r11	2		add	4294967295, r11	2	
	add	4, r10	2		cmp	r0, r11	2	
	st.w	r0, [r11]	4		jgt	.L23	2	
	add	4294967295, r12	2					
	cmp	r0, r12	2					
	jgt	.L23	2					
Total code	size		62 bytes	Total code	size		50 bytes	

		Before Modification		After Modification			
	Program			Program			Size [bytes]
	mov	2, r6	2		mov	2, r10	2
	movea	\$_array, gp, r10	4		movea	<pre>\$_array, gp, r12</pre>	4
	mov	r6, r11	2		mov	r10, r11	2
.L16:				.L16:			
	addi	4, r10, r13	4		st.w	r0, [r12]	4
	st.w	r0, [r10]	4		st.w	r0, 4[r12]	4
	addi	4, r13, r14	4		st.w	r0, 8[r12]	4
	st.w	r0, [r13]	4		st.w	r0, 12[r12]	4
	addi	4, r14, r7	4		add	16, r12	4
	st.w	r0, [r14]	4		add	4294967295, r11	2
	addi	4, r7, r10	4		cmp	r0, r11	2
	st.w	r0, [r7]	4		jgt	.L16	2
	add	4294967295, r11	2		mov	r10, r11	2
	cmp	r0, r11	2	.L23:			
	jgt	.L16	2		mov	r12, r10	2
	mov	r6, r12	2		add	4, r12	2
.L23:					st.w	r0, [r10]	4
	mov	r10, r11	2		add	4294967295, r11	2
	add	4, r10	2		cmp	r0, r11	2
	st.w	r0, [r11]	4		jgt	.L23	2
	add	4294967295, r12	2				
	cmp	r0, r12	2				
	jgt	.L23	2				
Total code	size		62 bytes	Total code	size		50 bytes

3.7 Outputting a setf Instruction to Output 0 or 1 According to the Result of a Conditional Comparison

Branch instructions are output by the if~else statement. As a result, the pipeline, which is the greatest feature of a RISC CPU, is disrupted. Coding that is not an if~else statement can be written so that a setf instruction (that outputs 0 or 1 according to the result of a conditional comparison) is output for assigning a variable.

A sample program is shown below.

Remark s and flag are assumed to be external variables.

It is also assumed that flag = 1 if the variable s is greater than 100, and flag = 0 if the variable s is less than or equal to 100.

Before Modification	After Modification
if(s > 100){	flag = (s > 100);
flag = 1;	
}	
else{	
flag = 0;	
}	

[Output Assembly List for the V850]

	Before Modification		After Modification		
	Program	Size [bytes]	Program	Size [bytes]	
ld.	i -100, r10, r0	4 4 2	ld.w \$_s, r10 cmp 100, r10	4 6 4	
jle mov st.	1, r11	2 4	setfgt r11 st.w r11, \$_flag	4	
jbr .L2:		2			
st. .L3:	w r0, \$_flag	4			
Total code si	ze	22 bytes	Total code size	18 bytes	

	Before Modification		After Modification		
Program		Size [bytes]	Program	Size [bytes]	
ld.w addi jle mov st.w jbr	-100, r10, r0 .L2 1, r11 r11, \$_flag .L3	4 4 2 2 4 2	ld.w \$_s, r10 cmp 100, r10 setfgt r11 st.w r11, \$_flag	4 6 4 4	
st.w .L3:		4 22 bytes	Total code size	18 bytes	

3.8 Using at Most Four Arguments

A function is usually called with arguments.

The values of the first four arguments are copied to registers r6, r7, r8, and r9 to call the function. The called side performs its processing by using the values of these registers as the first four arguments. The fifth and subsequent arguments are loaded in the stack when the function is called.

The called side fetches the fifth and subsequent arguments from the stack. If the number of arguments is increased from four to five, the overhead when the function is called increases because memory access is added to processing that had been accomplished by using only registers. Therefore, remember to use at most four arguments when designing the passing of arguments.

3.9 Using at Most 10 Local Variables (auto Variables) and Only 6 or 7 If Possible

Local variables (auto variables) are allocated to registers.

The V800 compiler (in 32-register mode) can use a total of 20 registers for variables within a single function since 10 registers are used as work registers and 10 registers for register variables.

If the processing within a single function is time consuming, many local variables should be used. All 20 registers can even be used for local variables.

If a function does not take very much time, it should use only the 10 work registers.

If the time overhead for pushing/popping registers that are used for register variables is taken into account, the use of register variables should be avoided. The compiler freely decides whether or not to use register variables.

Holding the number of local variables to approximately 6 or 7 is recommended.

The remaining 3 or 4 registers are used as real work registers.

3.10 Rearranging an Expression in Advance

When an expression is written, the compiler will normally expand the expression. However, the number of calculations can be reduced depending on how the expression is written.

A sample program is shown below.

Remark x and y are assumed to be external variables.

Before Modification (7 calculations)	After Modification (6 calculations)		
y = 7*x*x*x+5*x*x+x;	y = x*(7*x*x+5*x+1);		

	Before Modification			After Modification		
	Program	Size [bytes]		Program	Size [bytes]	
ld.w	\$_x, r7	4	ld.w	\$_x, r6	4	
mov	r7, r6	2	mov	r6, r13	2	
mov	r7, r11	2	mov	r6, r7	2	
jarl	mul, lp	4	jarl	mul, lp	4	
mov	r6, r14	2	mov	r6, r11	2	
mov	r6, r7	2	shl	3, r6	2	
mov	r11, r6	2	sub	r11, r6	2	
jarl	mul, lp	4	mov	r13, r7	2	
mov	r6, r13	2	shl	2, r13	2	
shl	3, r13	2	add	r7, r13	2	
sub	r6, r13	2	add	r13, r6	2	
mov	r14, r10	2	add	1, r6	2	
shl	2, r14	2	jarl	mul, lp	4	
add	r10, r14	2	st.w	r6, \$_y	4	
add	r14, r13	2				
add	r11, r13	2				
st.w	r13, \$_y	4				
Total code size		42 bytes	Total code size		36 bytes	

[Output Assembly List for the V850E]

	Before Modification			After Modification		
	Program	Size [bytes]		Program	Size [bytes]	
ld.w	\$_x, r11	4	ld.w	\$_x, r12	4	
mov	r11, r13	2	mov	r12, r11	2	
mul	r11, r13, r0	4	mul	r12, r11, r0	4	
mov	r13, r12	2	mul	7, r11, r0	4	
mul	r11, r12, r0	4	mov	r12, r10	2	
mul	7, r12, r0	4	mul	5, r12, r0	4	
mul	5, r13, r0	4	add	r12, r11	2	
add	r13, r12	2	add	1, r11	2	
add	r11, r12	2	mul	r10, r11, r0	4	
st.w	r12, \$_y	4	st.w	r11, \$_y	4	
Total code size		32 bytes	Total code size		32 bytes	

Remark In the V850E, the same code size will be produced by compiler optimization.

3.11 Replacing a Multiplication or Division Involving a Power of 2 with a Shift Operation

If a formula contains a multiplication or division involving a power of 2 (2, 4, 8, 16, 32, ...), the execution speed can be increased by replacing it with a shift instruction.

Although this normally should be left to the compiler, a calculation that is a power of 2 should explicitly be replaced with a shift instruction when coding.

Note, however, that this method will create a problem if negative values are handled. Only use it when positive values are handled. If unsigned type is used for handling positive values, no problem will occur.

A sample program is shown below.

Remark s, t, and u are assumed to be external variables (unsigned int).

Before Modification	After Modification		
s = s / 2;	s = s >> 1;		
t = t * 8;	t = t << 3;		
u = u * 64;	u = u << 6;		

In this example, the same code size will be produced by compiler optimization.

CHAPTER 4 DEFINING VARIABLES

This chapter summarizes precautions that should be kept in mind when defining variables.

4.1 Data Alignment

Data in memory must be aligned according to the device architecture. Therefore, the compiler places variables so that they are aligned properly (by inserting padding areas without changing the order).

The basic alignment conditions are that char-type data is at a one-byte boundary, short-type data is at a two-byte boundary, and int-type data is at a four-byte boundary.

When defining variables, arrange them so that they are in decreasing data length.

A sample program is shown below.

Before Modification	After Modification				
struct{	struct{				
char data1;	long data2;				
long data2;	long data4;				
short data3;	short data3;				
long data4;	char data1;				
}data;	}data;				

The memory allocation is shown below.

Before Modification (16 bytes)

After Modification	(12	bytes)	١
--------------------	-----	--------	---

Hig	her	
dat	a4	
	dat	a3
dat	a2	
		data1
Lov	ver	

Higher					
	data1	data3			
	data4				
data2					
	Lower				

Caution The shaded areas indicate padding areas.

4.2 volatile Specification

When coding external I/Os such as ports or external variables used for interrupt servicing, include a volatile specification. If volatile is not specified, an unexpected operation may occur due to the C compiler optimization.

A sample program is shown below.

Remark a is assumed to be an external variable.

```
Before Modification
#define PORT1
                 *((unsigned char *)0x100000) /* Address 0x100000 8 bits */
#define PORT2
                 *((unsigned short *)0x100004)
                                                  /* Address 0x100000 16 bits */
#define PORT3
                 *((unsigned int *)0x100008)
                                                  /* Address 0x100000 32 bits */
struct bitf {
                                                    /* Bit field */
 unsigned char bit00:1;
 unsigned char bit01:1;
 unsigned char bit02:1;
 unsigned char bit03:1;
 unsigned char bit04:1;
 unsigned char bit05:1;
 unsigned char bit06:1;
 unsigned char bit07:1;
#define PORTb
               ((struct bitf *)0x100000)->bit00 /* Address 0x100000 0th bit */
void func()
                       /* Write to PORT1 */
        PORT1 = 0xFF;
                        /* Read from PORT1 */
        a = PORT1;
                        /* Set PORTb */
        PORTb = 1;
```

```
After Modification
#define PORT1
                 *((volatile unsigned char *)0x100000)
                                                            /* Address 0x100000 8 bits */
#define PORT2
                 *((volatile unsigned short *)0x100004)
                                                         /* Address 0x100000 16 bits */
#define PORT3
                 *((volatile unsigned int *)0x100008)
                                                           /* Address 0x100000 32 bits */
                                                            /* Bit field */
struct bitf {
 unsigned char bit00:1;
 unsigned char bit01:1;
 unsigned char bit02:1;
 unsigned char bit03:1;
 unsigned char bit04:1;
 unsigned char bit05:1;
 unsigned char bit06:1;
 unsigned char bit07:1;
#define PORTb ((struct bitf *)0x100000)->bit00
                                                          /* Address 0x100000 0th bit */
void func()
 PORT1 = 0xFF; /* Write to PORT1 */
                /* Read from PORT1 */
 a = PORT1;
                 /* Set PORTb */
 PORTb = 1;
```

Before Modification			After Modification		
	Program		Program	Size [bytes]	
mov mov st.b st.b set1	1048576, r10 255, r11 r11, [r10] r11, \$_a 0, [r10]	4 4 4 4	mov 1048576, r10 mov 255, r11 st.b r11, [r10] ld.b [r10], r12 st.b r12, \$_a set1 0, [r10]	4 4 4 4	
Total code size		20 bytes	Total code size	24 bytes	

Before Modification				After Modification	
Program		Size [bytes]	Program		Size [bytes]
mov	1048576, r10	4	mov	1048576, r10	4
mov	255, r12	4	mov	255, r11	4
st.b	r12, [r10]	4	st.b	r11, [r10]	4
sxb	r12	2	ld.b	[r10], r12	4
st.b	r12, \$_a	4	st.b	r12, \$_a	4
set1	0, [r10]	4	set1	0, [r10]	4
Total code size		22 bytes	Total code size	_	24 bytes

4.3 Read-Only Variables

Specify the const variable as the variable for read only.

The variable definition will be allocated to ROM, reducing the size of the RAM variable area.

4.4 Reducing Alignment Between Files

When variables are defined in multiple files, an alignment hole may be generated between files during linking. The alignment hole between files can be avoided by gathering variable definitions in 1 file.

A sample program is shown below.

Before Modification	After Modification
sub1_1.c	sub2.c
long data1 = 0;	long data1 = 0;
char data2 = 0;	char data2 = 0;
	char data3 = 0;
sub1_2.c	char data4 = 0;
char data3 = 0;	
char data4 = 0;	

[Link map]

		E	Before Modification	
.sdata		0x00ffe00c	0x0000000a	
	.sdata	0x00ffe00c	0x0000005	sub1_1.o
		0x00ffe011	0x0000003	*(align-hole)*
	.sdata	0x00ffe014	0x0000002	sub1 2.0

After Modification					
.sdata	_	0x00ffe00c	0x0000007		
	.sdata	0x00ffe00c	0x00000007	sub2.o	

4.5 Consolidating Flags

Consolidate flags that have a capacity within a few bits in bit field.

A sample program is shown below.

Before Modification	After Modification
<pre>unsigned char flag1; unsigned char flag2; unsigned char flag3; flag1 = 1; flag2 = 1;</pre>	<pre>struct bitf { unsigned char flag1:1; unsigned char flag2:1; unsigned char flag3:1; } flags;</pre>
flag3 = 1;	<pre>flags.flag1 = 1; flags.flag2 = 1; flags.flag3 = 1;</pre>

The memory allocation is shown below.

Before Modification

After Modification

flag3				
flag2		flag3	flag2	flag1
flag1				

4.6 Reducing Nesting Levels of Functions

Reducing the nesting of functions by changing the algorithm may reduce the usage rate of the stack. Since the operation may not be affected much by inline expansion, the algorithm must be changed.

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